# Beyond Sturmian: a <br> characterization of octagon cutting sequences 

## Corinna Ulcigrai

(based on joint work with John Smillie Cornell University)

ICTP, Trieste, 23 July 2018

## Linear trajectories and cutting sequences

Consider a regular polygon, for simplicity with $2 n$ sides.
As an example, in the talk we will consider a regular octagon.
Glue opposite sides.
Label pairs of sides by $\{A, B, C, D\}$.
Let $\varphi_{t}^{\theta}$ be the linear flow in direction $\theta$ :
trajectories which do not hit singularities
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## A classical case: Sturmian sequences

Consider the special case in which the polygon is a square.


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Sturmian sequences are characterized by having the smallest possible complexity among non-periodic sequences.
(Let $P_{w}(n)$ the number of words of lenght $n$ which appear in the
sequence $w: P_{w}(n)=n$ iff $w$ is periodic. Sturmian sequences satisfy
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In this case the cutting sequence correspond to the sequence of horizontal (letter A) and vertical (letter B) sides crossed by a line in direction $\theta$ in a square grid..

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Cutting sequences give a symbolic coding of the following systems:

- Translation surfaces
- Poligonal Billiards
- Interval exchange transformations


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Glueing opposite sides one gets a surface of genus 2 , with a flat metric with a singularity (it's a translation surface); $\varphi_{t}^{\theta}$ is the geodesic flow with respect to the flat metric;

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For example: A billiard with angles $\frac{\pi}{2}, \frac{\pi}{8}$ and $\frac{3 \pi}{8}$
with elastic reflections at sides) by a procedure called unfolding is
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Problem:
Describe explicitely the symbolic sequences which are cutting sequences of trajectories.

In particular, answer the following questions:
D1) Which sequences in $\{A, B, C, D\}^{\mathbb{Z}}$ are cutting sequences?
D2) Given a cutting sequence, can one recover the direction of the trajectory?
Given a finite piece of a cutting sequence, can one recover a sector of possible directions?

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    - Caracterization of Sturmian sequences;
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    > Connection with Continued Fractions;
    * Sketch of proof for the square;
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joint work with Iohn Smillie (Cormell University)
- Formulation of results in the case of the octagon;
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## The octagon: isometries and sectors

Let $D_{8}$ be the isometries group of the octagon.
The letters $\{A, B, C, D\}$ are invariant with respect to a central symmetry. The other
elements induce permutations of
$\{A, B, C, D\}$ for example:
$A \mapsto C, C \mapsto A, B \mapsto B, D \mapsto D$
Let us assume that the direction of the trajectory is $\theta \in\left[0, \frac{\pi}{8}\right]$

 we can consider $\theta \in \Sigma_{0}:=\left[0, \frac{\pi}{8}\right]$

The other sectors of angle $\pi / 8$ are in order $\Sigma_{1}, \Sigma_{2}, \ldots, \Sigma_{7}$.

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$A \mapsto C, C \mapsto A, B \mapsto B, D \mapsto D$.
Let us assume that the direction of the trajectory is $\theta \in\left[0, \frac{\pi}{8}\right]$.
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## The octagon: possible transitions

Permuting the letters we obtain the diagrams corresponding to the other sectors:



## Admissible sequences

Definition
A sequence $w \in\{A, B\}^{\mathbb{Z}}$ is admissible if it gives an infinite path on one of the following diagrams:

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Lemma
An octagon cutting sequence is admissible.

## Derived sequences

Definition
A letter in $\{A, B, C, D\}$ is sandwitched if it is preceeded and followed by the same letter.

Example
In D B B C B A A D the letter C is sandwitched between to Bs.
Definition (Derived sequence)
If $w$ is an octagon cutting sequence, the derived sequence $w^{\prime}$ is obtained erasing all letters which are NOT sandwitched.

Example

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Example
If $w=\ldots$ D A D B C C B C C B DADBCBDBDBCBD...,

Definition (Derivable sequences)
A sequence $w \in\{A, B, C, D\}^{\mathbb{Z}}$ is derivable if it is admissible and its
derivative is still admissible. The sequence $w$ is infinitely derivable if each
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$\begin{aligned} \text { If } w & =\ldots \text { D } \frac{A}{A} \text { D B C C B C C B D A D B C B D B D B C B D } \ldots,\end{aligned}$
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## Necessary condition and sequence of sectors

Theorem
An octagon cutting sequence is infinitly derivable.
The converse is not true, but we can describe exactly the condition which one needs to add.

Definition
Let $w$ be infinitly derivable and let $w^{(n)}$ be the $n^{t h}$ derived sequence. The sequence $\left\{s_{k}\right\}_{k \in \mathbb{N}} \in\{0,1, \ldots, 7\}^{\mathbb{N}}$ is a sequence of sectors for $w$ if for each $k, w^{(k)}$ gives a path on the diagram $\mathscr{D}_{s_{k}}$.

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w=C \underline{C} C \underline{B C C B} \underline{D} B
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w=\mathrm{C} \underline{C} C \underline{B} C C B \underline{D} B \\
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s_{1}=4
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## Octagon Continued Fractions

Let $F:[0, \pi] \rightarrow[\pi / 8, \pi]$ the following map, that we call Octagon Farey map:

Definition
The octagon continued fraction expansion of $\theta$ is

$$
\theta=\left[s_{0}, s_{1}, s_{2}, \ldots, s_{k}, \ldots\right]_{O} \quad \text { iff } \quad\{\theta\}=\cap_{k} F_{s_{0}}^{-1} F_{s_{1}}^{-1} \ldots F_{s_{k}}^{-1}[0, \pi] .
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## Direction Recognitionn

Let $w$ be an octagon cutting sequence.
Lemma
If $w$ is not a periodic sequence, the sequence of sectors $\left\{s_{k}\right\}_{k \in \mathbb{N}}$ is univoquely determined. In particular, each derivative $w^{(k)}$ is admissible in an unique $\mathscr{D}_{s_{k}}$.

Theorem
If $w$ is not periodic, there is a unique sequence of sectors $\left\{s_{k}\right\}_{k \in \mathbb{N}}$ for $w$ and the direction of the trajectories with cutting sequence $w$ is given by

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## Ideas from proofs

As in the case of the square, the theorems follow if we prove that:
Theorem
If $w$ is an octagon cutting sequence, also the derived sequence $w^{\prime}$ is an octagon cutting sequence.

Furthermore, if $w$ is the cutting sequence of a trajectory in direction $\theta$, the derived sequence $w^{\prime}$ is a cutting sequence of a trajectory in direction $\theta^{\prime}=F(\theta)$, where $F$ is the octagon Farey map.
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i.e. one uses renormalization in the space of affine deformations of the octagon.
Regular polygons are rich of affine automorphism.

## Affine automorphism and Veech group

Let $S_{O}$ be the surface obtained glueing opposite sides of the octagon by translations: $S_{O}$ is an example of a translation surface, i.e. it has an atlas whose changes of coordinates are of the form $z \mapsto z+c$.

> Definition
> An automorphism $\Psi: S \mapsto S$ of a translation surface $S$ is an affine automorphism if it is affine in each chart and $D \Psi(z)$ is independent on $z \in S$. Let
> $\operatorname{Aff}(S)=\{\Psi, \quad \Psi$ affine automorphism $\}$

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Example (1)
Se $\psi \in D_{8}$
is an isometry of $O$, clearly one has
$\psi \in \operatorname{Aff}(O)$.

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Example (2)
The matrix $\left(\begin{array}{cc}1 & 2(1+\sqrt{2}) \\ 0 & 1\end{array}\right) \in V(O)$.

## The Veech group of the octagon

The Veech group $V(S)$ is the group of linear parts of $A f f(S)$ :

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V(S)=\{D \Psi, \quad \Psi \in \operatorname{Aff}(S)\} \subset S L(2, \mathbb{R})^{ \pm}
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Ex $1 \quad V\left(\mathbb{T}^{2}\right)^{+}=<\left(\begin{array}{ll}1 & 1 \\ 0 & 1\end{array}\right), \quad\left(\begin{array}{ll}1 & 0 \\ 1 & 1\end{array}\right)>=S L(2, \mathbb{Z})$.


If $S$ is a translation surface glued from a regular polygon, $V(S)$ is a lattice in $S L(2, \mathbb{R})^{ \pm}$(Veech)
The surfaces $S$ for which $V(S)$ is a lattice are actively researched in Teichmüller dynamics.

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## Derivation and renormalization in the octagon

Let $w$ be the cutting sequence of a trajectory in direction $\theta \in \Sigma_{0}$. Let
$O^{\prime}=\left(\begin{array}{cc}-1 & 2(1+\sqrt{2}) \\ 0 & 1\end{array}\right) O$.
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## Renormalization, modular surface and continued fractions

The space of lattices is $\mathbb{H} / S L(2, \mathbb{Z})$. (moduli space of tori with a flat metric)
Farey Tessellation: $V(Q)=<\left(\begin{array}{ll}1 & 1 \\ 0 & 1\end{array}\right),\left(\begin{array}{ll}1 & 0 \\ 1 & 1\end{array}\right)>\mathscr{T}(0,1, \infty)$

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Renormalization and dynamics on the Teichmüller disk $S L(2, \mathbb{R}) \cdot O=\{A \cdot O, A \in S L(2, \mathbb{R})\}$ affine deformations of the octagon

(Teichmueller disk)

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## Generation of cutting sequences

Let us define operators $g_{i}^{j}$ that invert derivation.
The operator $g_{i}^{j}$ interpolates a sequence $w$ admissible in $\mathscr{D}_{j}$ producing a sequence ammissible in $\mathscr{D}_{i}$ and such that $\left(g_{i}^{j}(w)\right)^{\prime}=w$.
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Let $w$ be ammissible in $\mathscr{D}_{0}$. The sequence $g_{0}^{2} w$ us obtained interpolating the letters corresponding to vertices in Figure with the words on the arrows:

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Example $w \quad=\ldots$ D C A B B B A C $\ldots$

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$$
\begin{array}{ll}
\mathrm{w} & =\ldots \\
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\text { D B C BD A } & \text { DBCC B B } \ldots
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## The interpolation operators




The other operators are obtained from these ones by permuting the letters.

## Characterization of the closure of cutting sequences

Lemma
If $w$ is a cutting sequence and $\left\{s_{n}\right\}_{n \in \mathbb{N}}$ a sequence of sectors, we have

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w \in \cap_{n} g_{s_{1}}^{s_{0}} g_{s_{2}}^{s_{1}} \ldots g_{s_{n}}^{s_{n-1}}\{A, B, C, D\}^{\mathbb{Z}}
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This condition, together with infinite derivability, is necessary and sufficient to characterize the closure of octagon cutting sequences (in $\{A, B, C, D\}^{Z}$ ):

Theorem (Smillie-U)
A sequence $w \in\{A, B, C, D\}^{2}$ belongs to the closure of cutting sequences in the octagon iff there exists a squence $\left\{s_{n}\right\}_{n \in \mathbb{N}} \in\{0$, such that
$w \in \cap_{n} g_{s_{1}}^{s_{0}} g_{s_{2}}^{s_{1}} \ldots g_{s_{n}}^{s_{n-1}}\{A, B, C, D\}^{\mathbb{Z}}$.

An analogous theorem holds for every regular polyon.

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