Beyond Sturmian: a characterization of octagon cutting sequences

Corinna Ulcigrai

(based on joint work with John Smillie Cornell University)

ICTP, Trieste, 23 July 2018

▲□▶ ▲□▶ ▲三▶ ▲三▶ 三三 のへで

Consider a regular polygon, for simplicity with 2n sides. As an example, in the talk we will consider a *regular octagon*. Glue opposite sides. Label pairs of sides by $\{A, B, C, D\}$.

Let φ_t^{θ} be the *linear flow* in direction θ : trajectories which do not hit singularities are straight lines in direction θ .

Definition (Cutting sequence)

The *cutting sequence* in $\{A, B, C, D\}^{\mathbb{Z}}$ that codes a bi-infinite linear trajectory of φ_t^{θ} consists of the sequence of labels of the sides hit by the trajectory.

・ロト ・四ト ・ヨト ・ヨト

- 2

Example

Consider a regular polygon, for simplicity with 2n sides. As an example, in the talk we will consider a *regular octagon*. Glue opposite sides. Label pairs of sides by $\{A, B, C, D\}$.

Let φ_t^{θ} be the *linear flow* in direction θ : trajectories which do not hit singularities are straight lines in direction θ .

Definition (Cutting sequence)

The *cutting sequence* in $\{A, B, C, D\}^{\mathbb{Z}}$ that codes a bi-infinite linear trajectory of φ_t^{θ} consists of the sequence of labels of the sides hit by the trajectory.

Example

Consider a regular polygon, for simplicity with 2n sides. As an example, in the talk we will consider a *regular octagon*. Glue opposite sides. Label pairs of sides by $\{A, B, C, D\}$.

Let φ_t^{θ} be the *linear flow* in direction θ : trajectories which do not hit singularities are straight lines in direction θ .

Definition (Cutting sequence)

The *cutting sequence* in $\{A, B, C, D\}^{\mathbb{Z}}$ that codes a bi-infinite linear trajectory of φ_t^{θ} consists of the sequence of labels of the sides hit by the trajectory.

Example

Consider a regular polygon, for simplicity with 2n sides. As an example, in the talk we will consider a *regular octagon*. Glue opposite sides. Label pairs of sides by $\{A, B, C, D\}$.

Let φ_t^{θ} be the *linear flow* in direction θ : trajectories which do not hit singularities are straight lines in direction θ .

Definition (Cutting sequence)

The *cutting sequence* in $\{A, B, C, D\}^{\mathbb{Z}}$ that codes a bi-infinite linear trajectory of φ_t^{θ} consists of the sequence of labels of the sides hit by the trajectory.

Example

Consider a regular polygon, for simplicity with 2n sides. As an example, in the talk we will consider a *regular octagon*. Glue opposite sides. Label pairs of sides by $\{A, B, C, D\}$.

Let φ_t^{θ} be the *linear flow* in direction θ : trajectories which do not hit singularities are straight lines in direction θ .

Definition (Cutting sequence)

The *cutting sequence* in $\{A, B, C, D\}^{\mathbb{Z}}$ that codes a bi-infinite linear trajectory of φ_t^{θ} consists of the sequence of labels of the sides hit by the trajectory.

Example

Consider a regular polygon, for simplicity with 2n sides.

As an example, in the talk we will consider a *regular octagon*.

Glue opposite sides.

Label pairs of sides by $\{A, B, C, D\}$.

Let φ_t^{θ} be the *linear flow* in direction θ : trajectories which do not hit singularities are straight lines in direction θ .

Definition (Cutting sequence)

The *cutting sequence* in $\{A, B, C, D\}^{\mathbb{Z}}$ that codes a bi-infinite linear trajectory of φ_t^{θ} consists of the sequence of labels of the sides hit by the trajectory.

Example

The cutting sequence of the trajectory in the example is:



<ロト <四ト <注ト <注ト = 三

Consider a regular polygon, for simplicity with 2n sides.

As an example, in the talk we will consider a *regular octagon*.

Glue opposite sides.

Label pairs of sides by $\{A, B, C, D\}$.

Let φ_t^{θ} be the *linear flow* in direction θ : trajectories which do not hit singularities are straight lines in direction θ .

Definition (Cutting sequence)

The cutting sequence in $\{A, B, C, D\}^{\mathbb{Z}}$ that codes a bi-infinite linear trajectory of φ_t^{θ} consists of the sequence of labels of the sides hit by the trajectory.

Example

The cutting sequence of the trajectory in the example is:



Consider a regular polygon, for simplicity with 2n sides.

As an example, in the talk we will consider a *regular octagon*.

Glue opposite sides.

Label pairs of sides by $\{A, B, C, D\}$.

Let φ_t^{θ} be the *linear flow* in direction θ : trajectories which do not hit singularities are straight lines in direction θ .

Definition (Cutting sequence)

The cutting sequence in $\{A, B, C, D\}^{\mathbb{Z}}$ that codes a bi-infinite linear trajectory of φ_t^{θ} consists of the sequence of labels of the sides hit by the trajectory.

Example

The cutting sequence of the trajectory in the example is:



< □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > □ □

Consider a regular polygon, for simplicity with 2n sides. As an example, in the talk we will consider a *regular octagon*. Glue opposite sides. Label pairs of sides by $\{A, B, C, D\}$.

Let φ_t^{θ} be the *linear flow* in direction θ : trajectories which do not hit singularities are straight lines in direction θ .

Definition (Cutting sequence)

The *cutting sequence* in $\{A, B, C, D\}^{\mathbb{Z}}$ that codes a bi-infinite linear trajectory of φ_t^{θ} consists of the sequence of labels of the sides hit by the trajectory.

Example

Consider a regular polygon, for simplicity with 2n sides. As an example, in the talk we will consider a *regular octagon*. Glue opposite sides. Label pairs of sides by $\{A, B, C, D\}$.

Let φ_t^{θ} be the *linear flow* in direction θ : trajectories which do not hit singularities are straight lines in direction θ .

Definition (Cutting sequence)

The cutting sequence in $\{A, B, C, D\}^{\mathbb{Z}}$ that codes a bi-infinite linear trajectory of φ_t^{θ} consists of the sequence of labels of the sides hit by the trajectory.

Example

Consider a regular polygon, for simplicity with 2n sides. As an example, in the talk we will consider a *regular octagon*. Glue opposite sides. Label pairs of sides by $\{A, B, C, D\}$.

Let φ_t^{θ} be the *linear flow* in direction θ : trajectories which do not hit singularities are straight lines in direction θ .

Definition (Cutting sequence)

The cutting sequence in $\{A, B, C, D\}^{\mathbb{Z}}$ that codes a bi-infinite linear trajectory of φ_t^{θ} consists of the sequence of labels of the sides hit by the trajectory.

Example

Consider a regular polygon, for simplicity with 2n sides. As an example, in the talk we will consider a *regular octagon*. Glue opposite sides. Label pairs of sides by $\{A, B, C, D\}$.

Let φ_t^{θ} be the *linear flow* in direction θ : trajectories which do not hit singularities are straight lines in direction θ .

Definition (Cutting sequence)

The cutting sequence in $\{A, B, C, D\}^{\mathbb{Z}}$ that codes a bi-infinite linear trajectory of φ_t^{θ} consists of the sequence of labels of the sides hit by the trajectory.

Example

Consider a regular polygon, for simplicity with 2n sides.

As an example, in the talk we will consider a *regular octagon*.

Glue opposite sides.

Label pairs of sides by $\{A, B, C, D\}$.

Let φ_t^{θ} be the *linear flow* in direction θ : trajectories which do not hit singularities are straight lines in direction θ .

Definition (Cutting sequence)

The cutting sequence in $\{A, B, C, D\}^{\mathbb{Z}}$ that codes a bi-infinite linear trajectory of φ_t^{θ} consists of the sequence of labels of the sides hit by the trajectory.

Example



Consider a regular polygon, for simplicity with 2n sides. As an example, in the talk we will consider a *regular octagon*. Glue opposite sides. Label pairs of sides by $\{A, B, C, D\}$.

Let φ_t^{θ} be the *linear flow* in direction θ : trajectories which do not hit singularities are straight lines in direction θ .

Definition (Cutting sequence)

The cutting sequence in $\{A, B, C, D\}^{\mathbb{Z}}$ that codes a bi-infinite linear trajectory of φ_t^{θ} consists of the sequence of labels of the sides hit by the trajectory.

Example

Consider a regular polygon, for simplicity with 2n sides. As an example, in the talk we will consider a *regular octagon*. Glue opposite sides. Label pairs of sides by $\{A, B, C, D\}$.

Let φ_t^{θ} be the *linear flow* in direction θ : trajectories which do not hit singularities are straight lines in direction θ .

Definition (Cutting sequence)

The cutting sequence in $\{A, B, C, D\}^{\mathbb{Z}}$ that codes a bi-infinite linear trajectory of φ_t^{θ} consists of the sequence of labels of the sides hit by the trajectory.

Example

Consider the special case in which the polygon is a square.



Minima complessita'

Sturmian sequences are characterized by having the smallest possible *complexity* among non-periodic sequences.

(Let $P_w(n)$ the number of words of lenght *n* which appear in the sequence *w*: $P_w(n) = n$ iff *w* is periodic. Sturmian sequences satisfy $P_w(n) = n + 1$.)

Consider the special case in which the polygon is a square.



In this case the cutting sequence correspond to the sequence of horizontal (letter A) and vertical (letter B) sides crossed by a line in direction θ in a square grid.

▲□▶ ▲□▶ ▲□▶ ▲□▶ □ のQで

Consider the special case in which the polygon is a square.



Square cutting sequences are *Sturmian sequences*. They were studied since Hedlund e Morse.

Minimal complexity

Sturmian sequences are characterized by having the smallest possible *complexity* among non-periodic sequences.

(Let $P_w(n)$ the number of words of lenght *n* which appear in the sequence *w*: $P_w(n) = n$ iff *w* is periodic. Sturmian sequences satisfy $P_w(n) = n + 1$.)

Consider the special case in which the polygon is a square.



Square cutting sequences are *Sturmian sequences*. They were studied since Hedlund e Morse.

Minimal complexity

Sturmian sequences are characterized by having the smallest possible *complexity* among non-periodic sequences.

(Let $P_w(n)$ the number of words of lenght *n* which appear in the sequence *w*: $P_w(n) = n$ iff *w* is periodic. Sturmian sequences satisfy $P_w(n) = n + 1$.)

Cutting sequences give a symbolic coding of the following systems:

▲□▶ ▲□▶ ▲ □▶ ▲ □▶ □ のへぐ

- Translation surfaces
- Poligonal Billiards
- Interval exchange transformations

Cutting sequences give a symbolic coding of the following systems:

Translation surfaces

Glueing opposite sides one gets a surface of genus 2, with a flat metric with a singularity (it's a translation surface); φ_t^{θ} is the geodesic flow with respect to the flat metric;

- Poligonal Billiards
- Interval exchange transformations

Cutting sequences give a symbolic coding of the following systems:

- Translation surfaces
- Poligonal Billiards

For example: A billiard with angles $\frac{\pi}{2}$, $\frac{\pi}{8}$ and $\frac{3\pi}{8}$ (motion of a particle with elastic reflections at sides) by a procedure called *unfolding* is equivalent to the flow φ_t^{θ} in the octagon.

▲□▶ ▲□▶ ▲□▶ ▲□▶ ▲□ ● ● ●

Cutting sequences give a symbolic coding of the following systems:

- Translation surfaces
- Poligonal Billiards

For example: A billiard with angles $\frac{\pi}{2}$, $\frac{\pi}{8}$ and $\frac{3\pi}{8}$ (motion of a particle with elastic reflections at sides) by a procedure called *unfolding* is equivalent to the flow φ_t^{θ} in the octagon.

Cutting sequences give a symbolic coding of the following systems:

- Translation surfaces
- Poligonal Billiards

For example: A billiard with angles $\frac{\pi}{2}$, $\frac{\pi}{8}$ and $\frac{3\pi}{8}$ (motion of a particle with elastic reflections at sides) by a procedure called *unfolding* is equivalent to the flow φ_t^{θ} in the octagon.

Cutting sequences give a symbolic coding of the following systems:

- Translation surfaces
- Poligonal Billiards

For example: A billiard with angles $\frac{\pi}{2}$, $\frac{\pi}{8}$ and $\frac{3\pi}{8}$ (motion of a particle with elastic reflections at sides) by a procedure called *unfolding* is equivalent to the flow φ_t^{θ} in the octagon.

Cutting sequences give a symbolic coding of the following systems:

- Translation surfaces
- Poligonal Billiards

For example: A billiard with angles $\frac{\pi}{2}$, $\frac{\pi}{8}$ and $\frac{3\pi}{8}$ (motion of a particle with elastic reflections at sides) by a procedure called *unfolding* is equivalent to the flow φ_t^{θ} in the octagon.

Cutting sequences give a symbolic coding of the following systems:

- Translation surfaces
- Poligonal Billiards

For example: A billiard with angles $\frac{\pi}{2}$, $\frac{\pi}{8}$ and $\frac{3\pi}{8}$ (motion of a particle with elastic reflections at sides) by a procedure called *unfolding* is equivalent to the flow φ_t^{θ} in the octagon.

Cutting sequences give a symbolic coding of the following systems:

- Translation surfaces
- Poligonal Billiards

For example: A billiard with angles $\frac{\pi}{2}$, $\frac{\pi}{8}$ and $\frac{3\pi}{8}$ (motion of a particle with elastic reflections at sides) by a procedure called *unfolding* is equivalent to the flow φ_t^{θ} in the octagon.

Cutting sequences give a symbolic coding of the following systems:

- Translation surfaces
- Poligonal Billiards

For example: A billiard with angles $\frac{\pi}{2}$, $\frac{\pi}{8}$ and $\frac{3\pi}{8}$ (motion of a particle with elastic reflections at sides) by a procedure called *unfolding* is equivalent to the flow φ_t^{θ} in the octagon.

Cutting sequences give a symbolic coding of the following systems:

- Translation surfaces
- Poligonal Billiards

For example: A billiard with angles $\frac{\pi}{2}$, $\frac{\pi}{8}$ and $\frac{3\pi}{8}$ (motion of a particle with elastic reflections at sides) by a procedure called *unfolding* is equivalent to the flow φ_t^{θ} in the octagon.

Cutting sequences give a symbolic coding of the following systems:

- Translation surfaces
- Poligonal Billiards

For example: A billiard with angles $\frac{\pi}{2}$, $\frac{\pi}{8}$ and $\frac{3\pi}{8}$ (motion of a particle with elastic reflections at sides) by a procedure called *unfolding* is equivalent to the flow φ_t^{θ} in the octagon.

Cutting sequences give a symbolic coding of the following systems:

- Translation surfaces
- Poligonal Billiards

For example: A billiard with angles $\frac{\pi}{2}$, $\frac{\pi}{8}$ and $\frac{3\pi}{8}$ (motion of a particle with elastic reflections at sides) by a procedure called *unfolding* is equivalent to the flow φ_t^{θ} in the octagon.

Cutting sequences give a symbolic coding of the following systems:

- Translation surfaces
- Poligonal Billiards
- Interval exchange transformations

The Poincaré first return map on a section is an interval exchange transformation (IET). As θ changes, one has a one-paramter family which is not generic.

Cutting sequences give a symbolic coding of the following systems:

- Translation surfaces
- Poligonal Billiards
- Interval exchange transformations

The Poincaré first return map on a section is an interval exchange transformation (IET). As θ changes, one has a one-paramter family which is not generic.

Cutting sequences give a symbolic coding of the following systems:

- Translation surfaces
- Poligonal Billiards
- Interval exchange transformations

The Poincaré first return map on a section is an interval exchange transformation (IET). As θ changes, one has a one-paramter family which is not generic.
Motivation

Cutting sequences give a symbolic coding of the following systems:

- Translation surfaces
- Poligonal Billiards
- Interval exchange transformations

Remark

All these systems have entropy zero: cutting sequences have linear complexity.

Remark

Translation surfaces and IETs which come from regular polygons are not *generic*: techniques that are used for the generic setting do not apply here.

▲□▶ ▲□▶ ▲□▶ ▲□▶ ▲□ ● ● ●

Motivation

Cutting sequences give a symbolic coding of the following systems:

- Translation surfaces
- Poligonal Billiards
- Interval exchange transformations

Remark

All these systems have entropy zero: cutting sequences have linear complexity.

Remark

Translation surfaces and IETs which come from regular polygons are not *generic*: techniques that are used for the generic setting do not apply here.

▲□▶ ▲□▶ ▲□▶ ▲□▶ ▲□ ● ● ●

Motivation

Cutting sequences give a symbolic coding of the following systems:

- Translation surfaces
- Poligonal Billiards
- Interval exchange transformations

Remark

All these systems have entropy zero: cutting sequences have linear complexity.

Remark

Translation surfaces and IETs which come from regular polygons are not *generic*: techniques that are used for the generic setting do not apply here.

▲□▶ ▲□▶ ▲□▶ ▲□▶ ▲□ ● ● ●

Problem:

Describe explicitely the symbolic sequences which are cutting sequences of trajectories.

In particular, answer the following questions:

- D1) Which sequences in $\{A, B, C, D\}^{\mathbb{Z}}$ are cutting sequences?
- D2) Given a cutting sequence, can one recover the direction of the trajectory?Given a *finite* piece of a cutting sequence, can one recover a sector of possible directions?
- D3) Given a direction or more in general a sector of directions, can one produce all cutting sequences of trajectories in that direction or sector?

Problem:

Describe explicitely the symbolic sequences which are cutting sequences of trajectories.

In particular, answer the following questions:

- D1) Which sequences in $\{A, B, C, D\}^{\mathbb{Z}}$ are cutting sequences?
- D2) Given a cutting sequence, can one recover the direction of the trajectory? Given a *finite* piece of a cutting sequence can one recover a sect

Given a *finite* piece of a cutting sequence, can one recover a sector of possible directions?

<ロト <回ト < 注ト < 注ト = 注

D3) Given a direction or more in general a sector of directions, can one produce all cutting sequences of trajectories in that direction or sector?

Problem:

Describe explicitely the symbolic sequences which are cutting sequences of trajectories.

In particular, answer the following questions:

- D1) Which sequences in $\{A, B, C, D\}^{\mathbb{Z}}$ are cutting sequences?
- D2) Given a cutting sequence, can one recover the direction of the trajectory?Given a *finite* piece of a cutting sequence,can one recover a sector of possible directions?
- D3) Given a direction or more in general a sector of directions, can one produce all cutting sequences of trajectories in that direction or sector?

Problem:

Describe explicitely the symbolic sequences which are cutting sequences of trajectories.

In particular, answer the following questions:

D1) Which sequences in $\{A, B, C, D\}^{\mathbb{Z}}$ are cutting sequences?

D2) Given a cutting sequence, can one recover the direction of the trajectory?

Given a *finite* piece of a cutting sequence, can one recover a sector of possible directions?

《曰》 《聞》 《臣》 《臣》 三臣

D3) Given a direction or more in general a sector of directions, can one produce all cutting sequences of trajectories in that direction or sector?

Problem:

Describe explicitely the symbolic sequences which are cutting sequences of trajectories.

In particular, answer the following questions:

- D1) Which sequences in $\{A, B, C, D\}^{\mathbb{Z}}$ are cutting sequences?
- D2) Given a cutting sequence, can one recover the direction of the trajectory?Given a *finite* piece of a cutting sequence, can one recover a sector of possible directions?
- D3) Given a direction or more in general a sector of directions, can one produce all cutting sequences of trajectories in that direction or sector?

《曰》 《聞》 《臣》 《臣》 三臣 …

Problem:

Describe explicitely the symbolic sequences which are cutting sequences of trajectories.

In particular, answer the following questions:

- D1) Which sequences in $\{A, B, C, D\}^{\mathbb{Z}}$ are cutting sequences?
- D2) Given a cutting sequence, can one recover the direction of the trajectory?Given a *finite* piece of a cutting sequence, can one recover a sector of possible directions?
- D3) Given a direction or more in general a sector of directions, can one produce all cutting sequences of trajectories in that direction or sector?

▶ The classical case: the square (Sturmian sequences)

- Caracterization of Sturmian sequences; (revisiting Caroline Series work)
- Connection with Continued Fractions;
- Sketch of proof for the square;

Regular polygons with 2n lati:

joint work with John Smillie (Cornell University)

- Formulation of results in the case of the octagon;
- Continued Fractions for regular polygons;

Renormalization, Fuchsian groups and Teichmüller flow;

イロト イヨト イヨト イヨト

2

▶ The classical case: the square (Sturmian sequences)

- Caracterization of Sturmian sequences; (revisiting Caroline Series work)
- Connection with Continued Fractions;
- Sketch of proof for the square;

Regular polygons with 2n lati:

joint work with John Smillie (Cornell University)

- Formulation of results in the case of the octagon;
- Continued Fractions for regular polygons;

Renormalization, Fuchsian groups and Teichmüller flow;

《曰》 《圖》 《臣》 《臣》

- 2

▶ The classical case: the square (Sturmian sequences)

- Caracterization of Sturmian sequences;
 - (revisiting Caroline Series work)
- Connection with Continued Fractions;
- Sketch of proof for the square;

Regular polygons with 2n lati:

joint work with John Smillie (Cornell University)

- Formulation of results in the case of the octagon;
- Continued Fractions for regular polygons;

Renormalization, Fuchsian groups and Teichmüller flow;

・ロト ・四ト ・ヨト ・ヨト

- 2

► The classical case: the square (Sturmian sequences)

- Caracterization of Sturmian sequences; (revisiting Caroline Series work)
- Connection with Continued Fractions;
- Sketch of proof for the square;

Regular polygons with 2n lati:

joint work with John Smillie (Cornell University)

- Formulation of results in the case of the octagon;
- Continued Fractions for regular polygons;

Renormalization, Fuchsian groups and Teichmüller flow;

・ロト ・四ト ・ヨト ・ヨト

- 2

▶ The classical case: the square (Sturmian sequences)

- Caracterization of Sturmian sequences; (revisiting Caroline Series work)
- Connection with Continued Fractions;
- Sketch of proof for the square;

Regular polygons with 2n lati:

joint work with John Smillie (Cornell University)

- Formulation of results in the case of the octagon;
- Continued Fractions for regular polygons;

Renormalization, Fuchsian groups and Teichmüller flow;

<ロト <四ト <至ト <至ト = 至

▶ The classical case: the square (Sturmian sequences)

- Caracterization of Sturmian sequences; (revisiting Caroline Series work)
- Connection with Continued Fractions;
- Sketch of proof for the square;

Regular polygons with 2n lati:

joint work with John Smillie (Cornell University)

- Formulation of results in the case of the octagon;
- Continued Fractions for regular polygons;

Renormalization, Fuchsian groups and Teichmüller flow;

▶ The classical case: the square (Sturmian sequences)

- Caracterization of Sturmian sequences; (revisiting Caroline Series work)
- Connection with Continued Fractions;
- Sketch of proof for the square;

Regular polygons with 2n lati:

joint work with John Smillie (Cornell University)

- Formulation of results in the case of the octagon;
- Continued Fractions for regular polygons;

Renormalization, Fuchsian groups and Teichmüller flow;

▶ The classical case: the square (Sturmian sequences)

- Caracterization of Sturmian sequences; (revisiting Caroline Series work)
- Connection with Continued Fractions;
- Sketch of proof for the square;

Regular polygons with 2n lati:

joint work with John Smillie (Cornell University)

- Formulation of results in the case of the octagon;
- Continued Fractions for regular polygons;

Renormalization, Fuchsian groups and Teichmüller flow;

▶ The classical case: the square (Sturmian sequences)

- Caracterization of Sturmian sequences; (revisiting Caroline Series work)
- Connection with Continued Fractions;
- Sketch of proof for the square;

Regular polygons with 2n lati:

joint work with John Smillie (Cornell University)

- Formulation of results in the case of the octagon;
- Continued Fractions for regular polygons;

Renormalization, Fuchsian groups and Teichmüller flow;

▶ The classical case: the square (Sturmian sequences)

- Caracterization of Sturmian sequences; (revisiting Caroline Series work)
- Connection with Continued Fractions;
- Sketch of proof for the square;

Regular polygons with 2n lati:

joint work with John Smillie (Cornell University)

- Formulation of results in the case of the octagon;
- Continued Fractions for regular polygons;

Renormalization, Fuchsian groups and Teichmüller flow;

《曰》 《聞》 《臣》 《臣》 三臣

▶ The classical case: the square (Sturmian sequences)

- Caracterization of Sturmian sequences; (revisiting Caroline Series work)
- Connection with Continued Fractions;
- Sketch of proof for the square;

Regular polygons with 2n lati:

joint work with John Smillie (Cornell University)

- Formulation of results in the case of the octagon;
- Continued Fractions for regular polygons;

Renormalization, Fuchsian groups and Teichmüller flow;

▲□▶ ▲□▶ ▲□▶ ▲□▶ = ● ● ●

Let D_8 be the isometries group of the octagon.

The letters $\{A, B, C, D\}$ are invariant with respect to a central symmetry. The other elements induce permutations of $\{A, B, C, D\}$ for example: $A \mapsto C, C \mapsto A, B \mapsto B, D \mapsto D$.

Let us assume that the direction of the trajectory is $\theta \in [0, \frac{\pi}{8}]$. A fundamental domain for D_8 is $\Sigma_0 := [0, \frac{\pi}{8}]$. Thus, acting by an element of D_8 , up to a permutation of the letters $\{A, B, C, D\}$, we can consider $\theta \in \Sigma_0 := [0, \frac{\pi}{8}]$.

Let D_8 be the isometries group of the octagon.

The letters $\{A, B, C, D\}$ are invariant with respect to a central symmetry. The other

elements induce permutations of $\{A, B, C, D\}$ for example: $A \mapsto C, C \mapsto A, B \mapsto B, D \mapsto D$.

Let us assume that the direction of the trajectory is $\theta \in [0, \frac{\pi}{8}]$. A fundamental domain for D_8 is $\Sigma_0 := [0, \frac{\pi}{8}]$. Thus, acting by an element of D_8 , up to a permutation of the letters $\{A, B, C, D\}$, we can consider $\theta \in \Sigma_0 := [0, \frac{\pi}{8}]$.

Let D_8 be the isometries group of the octagon. The letters $\{A, B, C, D\}$ are invariant with respect to a central symmetry. The other elements induce permutations of $\{A, B, C, D\}$ for example: $A \mapsto C, C \mapsto A, B \mapsto B, D \mapsto D$.

Let us assume that the direction of the trajectory is $\theta \in [0, \frac{\pi}{8}]$. A fundamental domain for D_8 is $\Sigma_0 := [0, \frac{\pi}{8}]$. Thus, acting by an element of D_8 , up to a permutation of the letters $\{A, B, C, D\}$, we can consider $\theta \in \Sigma_0 := [0, \frac{\pi}{8}]$.

A D N A 目 N A E N A E N A B N A C N

Let D_8 be the isometries group of the octagon.

The letters $\{A, B, C, D\}$ are invariant with respect to a central symmetry. The other elements induce permutations of $\{A, B, C, D\}$ for example:

 $A\mapsto C, C\mapsto A, B\mapsto B, D\mapsto D$.

Let us assume that the direction of the trajectory is $\theta \in [0, \frac{\pi}{8}]$. A fundamental domain for D_8 is $\Sigma_0 := [0, \frac{\pi}{8}]$. Thus, acting by an element of D_8 , up to a permutation of the letters $\{A, B, C, D\}$, we can consider $\theta \in \Sigma_0 := [0, \frac{\pi}{8}]$.

A D N A 目 N A E N A E N A B N A C N

Let D_8 be the isometries group of the octagon.

The letters $\{A, B, C, D\}$ are invariant with respect to a central symmetry. The other elements induce permutations of $\{A, B, C, D\}$ for example: $A \mapsto C, C \mapsto A, B \mapsto B, D \mapsto D$.

Let us assume that the direction of the trajectory is $\theta \in [0, \frac{\pi}{8}]$. A fundamental domain for D_8 is $\Sigma_0 := [0, \frac{\pi}{8}]$. Thus, acting by an element of D_8 , up to a permutation of the letters $\{A, B, C, D\}$, we can consider $\theta \in \Sigma_0 := [0, \frac{\pi}{8}]$.

A D N A 目 N A E N A E N A B N A C N

Let D_8 be the isometries group of the octagon.

The letters $\{A, B, C, D\}$ are invariant with respect to a central symmetry. The other elements induce permutations of $\{A, B, C, D\}$ for example: $A \mapsto C, C \mapsto A, B \mapsto B, D \mapsto D$.

Let us assume that the direction of the trajectory is $\theta \in [0, \frac{\pi}{8}]$. A fundamental domain for D_8 is $\Sigma_0 := [0, \frac{\pi}{8}]$. Thus, acting by an element of D_8 , up to a permutation of the letters $\{A, B, C, D\}$, we can consider $\theta \in \Sigma_0 := [0, \frac{\pi}{8}]$.

A D N A 目 N A E N A E N A B N A C N

Let D_8 be the isometries group of the octagon.

The letters $\{A, B, C, D\}$ are invariant with respect to a central symmetry. The other elements induce permutations of $\{A, B, C, D\}$ for example: $A \mapsto C, C \mapsto A, B \mapsto B, D \mapsto D$.

Let us assume that the direction of the trajectory is $\theta \in [0, \frac{\pi}{8}]$. A fundamental domain for D_8 is $\Sigma_0 := [0, \frac{\pi}{8}]$. Thus, acting by an element of D_8 , up to a permutation of the letters $\{A, B, C, D\}$, we can consider $\theta \in \Sigma_0 := [0, \frac{\pi}{8}]$.

A D N A 目 N A E N A E N A B N A C N

Let $\theta \in \Sigma_0 := \left[0, \frac{\pi}{8}\right]$.

The transitions (pairs of consecutive letters) which can appear are:

Each octagon cutting sequence of a trajectory in direction $\theta \in \Sigma_0$ determines a path in the diagram in Figure.

▲□▶ ▲□▶ ▲ 三▶ ▲ 三▶ 三 のへぐ

Let
$$\theta \in \Sigma_0 := \left[0, \frac{\pi}{8}\right]$$
.

The transitions (pairs of consecutive letters) which can appear are:



Each octagon cutting sequence of a trajectory in direction $\theta \in \Sigma_0$ determines a path in the diagram in Figure.

・ロト ・ 国 ト ・ ヨ ト ・ ヨ ト

э

Let $\theta \in \Sigma_0 := \left[0, \frac{\pi}{8}\right]$.

The transitions (pairs of consecutive letters) which can appear are:



Each octagon cutting sequence of a trajectory in direction $\theta \in \Sigma_0$ determines a path in the diagram in Figure.

Let
$$\theta \in \Sigma_0 := \left[0, \frac{\pi}{8}\right]$$
.

The transitions (pairs of consecutive letters) which can appear are:



Each octagon cutting sequence of a trajectory in direction $\theta \in \Sigma_0$ determines a path in the diagram in Figure.

イロト イボト イヨト イヨト 三日

Let
$$\theta \in \Sigma_0 := \left[0, \frac{\pi}{8}\right]$$
.

The transitions (pairs of consecutive letters) which can appear are:



Each octagon cutting sequence of a trajectory in direction $\theta \in \Sigma_0$ determines a path in the diagram in Figure.

イロト 不得 トイヨト イヨト

3

Let $\theta \in \Sigma_0 := \left[0, \frac{\pi}{8}\right]$.

The transitions (pairs of consecutive letters) which can appear are:

R

・ロト ・ 国 ト ・ ヨ ト ・ ヨ ト

э

Each octagon cutting sequence of a trajectory in direction $\theta \in \Sigma_0$ determines a path in the diagram in Figure.

The octagon: possible transitions

Permuting the letters we obtain the diagrams corresponding to the other sectors:











▲□▶ ▲□▶ ▲□▶ ▲□▶ □ のQで

Admissible sequences

Definition

A sequence $w \in \{A, B\}^{\mathbb{Z}}$ is *admissible* if it gives an infinite path on one of the following diagrams:

・ロト ・ 同ト ・ ヨト ・ ヨト

ж



Lemma

An octagon cutting sequence is admissible.

Admissible sequences

Definition

A sequence $w \in \{A, B\}^{\mathbb{Z}}$ is *admissible* if it gives an infinite path on one of the following diagrams:



・ロト ・ 同ト ・ ヨト ・ ヨト

ж

Lemma

An octagon cutting sequence is admissible.
Definition

A letter in $\{A, B, C, D\}$ is *sandwitched* if it is preceded and followed by the same letter.

Example

In D B B C B A A D the letter C is *sandwitched* between to Bs.

Definition (Derived sequence)

If w is an octagon cutting sequence, the derived sequence w' is obtained erasing all letters which are *NOT sandwitched*.

Example

 $\begin{array}{ll} \mathsf{If} \ \mathsf{w} \ = \ \ldots \ \mathsf{D} \ \underline{\mathsf{A}} \ \mathsf{D} \ \mathsf{B} \ \mathsf{C} \ \mathsf{B} \ \mathsf{C} \ \mathsf{B} \ \mathsf{D} \ \mathsf{A} \ \mathsf{D} \ \mathsf{B} \ \mathsf{C} \ \mathsf{B} \ \mathsf{D} \ \mathsf{B} \ \mathsf{D} \ \mathsf{B} \ \mathsf{C} \ \mathsf{B} \ \mathsf{D} \ \mathsf{D} \ \mathsf{B} \ \mathsf{C} \ \mathsf{B} \ \mathsf{D} \ \mathsf{D} \ \mathsf{B} \ \mathsf{C} \ \mathsf{B} \ \mathsf{D} \ \mathsf{D} \ \mathsf{B} \ \mathsf{C} \ \mathsf{B} \ \mathsf{D} \ \mathsf{D} \ \mathsf{B} \ \mathsf{C} \ \mathsf{B} \ \mathsf{D} \ \mathsf{D} \ \mathsf{B} \ \mathsf{C} \ \mathsf{B} \ \mathsf{D} \ \mathsf{D} \ \mathsf{B} \ \mathsf{C} \ \mathsf{B} \ \mathsf{C} \ \mathsf{B} \ \mathsf{D} \ \mathsf{B} \ \mathsf{C} \ \mathsf{$

Definition (Derivable sequences)

A sequence $w \in \{A, B, C, D\}^{\mathbb{Z}}$ is *derivable* if it is admissible and its derivative is still admissible. The sequence w is *infinitely derivable* if each of its derivatives is derivable.

Definition

A letter in $\{A, B, C, D\}$ is *sandwitched* if it is preceded and followed by the same letter.

Example

In D B B C B A A D the letter C is *sandwitched* between to Bs.

Definition (Derived sequence)

If w is an octagon cutting sequence, the derived sequence w' is obtained erasing all letters which are *NOT* sandwitched.

Example

If w = ... D A D B C C B C C B D A D B C B D B D B C B D ..., w' = ... A

Definition (Derivable sequences)

Definition

A letter in $\{A, B, C, D\}$ is *sandwitched* if it is preceded and followed by the same letter.

Example

In D B B C B A A D the letter C is *sandwitched* between to Bs.

Definition (Derived sequence)

If w is an octagon cutting sequence, the derived sequence w' is obtained erasing all letters which are *NOT* sandwitched.

Example

If w = \dots D A D B C C B C C B D A D B C B D B D B C B D \dots , w' = \dots A

Definition (Derivable sequences)

Definition

A letter in $\{A, B, C, D\}$ is *sandwitched* if it is preceded and followed by the same letter.

Example

In D B B C B A A D the letter C is *sandwitched* between to Bs.

Definition (Derived sequence)

If w is an octagon cutting sequence, the derived sequence w' is obtained erasing all letters which are *NOT* sandwitched.

Example

```
\begin{array}{lll} \text{If } w \ = \ \ldots \ D \ \underline{A} \ D \ B \ C \ \underline{C} \ \underline{B} \ \underline{C} \ C \ B \ D \ A \ D \ B \ C \ B \ D \ B \ D \ B \ C \ B \ D \ \ldots, \\ w' \ = \ \ldots \ A \ B \end{array}
```

Definition (Derivable sequences)

Definition

A letter in $\{A, B, C, D\}$ is *sandwitched* if it is preceded and followed by the same letter.

Example

In D B B C B A A D the letter C is *sandwitched* between to Bs.

Definition (Derived sequence)

If w is an octagon cutting sequence, the derived sequence w' is obtained erasing all letters which are *NOT* sandwitched.

Example

 $\begin{array}{cccc} \mathsf{If} \ \mathsf{w} \ = \ \ldots \ \mathsf{D} \ \underline{\mathsf{A}} \ \mathsf{D} \ \mathsf{B} \ \mathsf{C} \ \mathsf{C} \ \underline{\mathsf{B}} \ \mathsf{C} \ \mathsf{C} \ \mathsf{B} \ \mathsf{D} \ \underline{\mathsf{A}} \ \mathsf{D} \ \mathsf{B} \ \mathsf{C} \ \mathsf{B} \ \mathsf{C} \ \mathsf{C} \ \mathsf{B} \ \mathsf{C} \ \mathsf{$

Definition (Derivable sequences)

Definition

A letter in $\{A, B, C, D\}$ is *sandwitched* if it is preceded and followed by the same letter.

Example

In D B B C B A A D the letter C is *sandwitched* between to Bs.

Definition (Derived sequence)

If w is an octagon cutting sequence, the derived sequence w' is obtained erasing all letters which are *NOT* sandwitched.

Example

Definition (Derivable sequences)

Definition

A letter in $\{A, B, C, D\}$ is sandwitched if it is preceded and followed by the same letter.

Example

In D B B C B A A D the letter C is *sandwitched* between to Bs.

Definition (Derived sequence)

If w is an octagon cutting sequence, the derived sequence w' is obtained erasing all letters which are *NOT* sandwitched.

Example

Definition (Derivable sequences)

Theorem An octagon cutting sequence is infinitly derivable.

The converse is not true, but we can describe exactly the condition which one needs to add.

Definition

Let w be infinitly derivable and let $w^{(n)}$ be the n^{th} derived sequence. The sequence $\{s_k\}_{k\in\mathbb{N}} \in \{0, 1, ..., 7\}^{\mathbb{N}}$ is a sequence of sectors for w if for each k, $w^{(k)}$ gives a path on the diagram \mathscr{D}_{s_k} .

$$w = C \subseteq C \subseteq B \subseteq C \subseteq B \supseteq E$$
$$s_0 = 0$$
$$w' = C \subseteq B \subseteq B \subseteq D$$
$$s_1 = 4$$

Theorem

An octagon cutting sequence is infinitly derivable.

The converse is not true, but we can describe exactly the condition which one needs to add.

Definition

Let w be infinitly derivable and let $w^{(n)}$ be the n^{th} derived sequence. The sequence $\{s_k\}_{k\in\mathbb{N}} \in \{0, 1, ..., 7\}^{\mathbb{N}}$ is a sequence of sectors for w if for each k, $w^{(k)}$ gives a path on the diagram \mathscr{D}_{s_k} .

$$w = C \subseteq C \subseteq B \subseteq C \subseteq B \subseteq C$$
$$s_0 = 0$$
$$w' = C \subseteq B \subseteq D \subseteq D$$
$$s_1 = 4$$

Theorem

An octagon cutting sequence is infinitly derivable.

The converse is not true, but we can describe exactly the condition which one needs to add.

Definition

Let w be infinitly derivable and let $w^{(n)}$ be the n^{th} derived sequence. The sequence $\{s_k\}_{k\in\mathbb{N}} \in \{0, 1, \dots, 7\}^{\mathbb{N}}$ is a sequence of sectors for w if for each k, $w^{(k)}$ gives a path on the diagram \mathscr{D}_{s_k} .

$$w = C \underline{C} C \underline{B} C C B \underline{D} B$$
$$s_0 = 0$$
$$w' = C B D B D$$
$$s_1 = 4$$

Theorem

An octagon cutting sequence is infinitly derivable.

The converse is not true, but we can describe exactly the condition which one needs to add.

Definition

Let w be infinitly derivable and let $w^{(n)}$ be the n^{th} derived sequence. The sequence $\{s_k\}_{k\in\mathbb{N}} \in \{0, 1, \dots, 7\}^{\mathbb{N}}$ is a sequence of sectors for w if for each k, $w^{(k)}$ gives a path on the diagram \mathscr{D}_{s_k} .

$$w= C \underline{C} C \underline{B} C C B \underline{D} B$$
$$s_0 = 0$$
$$w'= C B D B D$$
$$s_1 = 4$$

Theorem

An octagon cutting sequence is infinitly derivable.

The converse is not true, but we can describe exactly the condition which one needs to add.

Definition

Let w be infinitly derivable and let $w^{(n)}$ be the n^{th} derived sequence. The sequence $\{s_k\}_{k\in\mathbb{N}} \in \{0, 1, \dots, 7\}^{\mathbb{N}}$ is a sequence of sectors for w if for each k, $w^{(k)}$ gives a path on the diagram \mathscr{D}_{s_k} .

Example

$$w = C \subseteq C \subseteq B \subseteq C \subset B \supseteq B$$
$$s_0 = 0$$
$$w' = C \subseteq B \subseteq D \subseteq D$$
$$s_1 = 4$$

・ロト・西ト・西ト・西ト・日・ 今日・

Theorem

An octagon cutting sequence is infinitly derivable.

The converse is not true, but we can describe exactly the condition which one needs to add.

Definition

Let w be infinitly derivable and let $w^{(n)}$ be the n^{th} derived sequence. The sequence $\{s_k\}_{k\in\mathbb{N}} \in \{0, 1, \dots, 7\}^{\mathbb{N}}$ is a sequence of sectors for w if for each k, $w^{(k)}$ gives a path on the diagram \mathscr{D}_{s_k} .

Example

$$w = C \subseteq C \subseteq B \subseteq C \subseteq B \supseteq B$$
$$s_0 = 0$$
$$w' = C \subseteq B \subseteq D \subseteq D$$
$$s_1 = 4$$

Theorem

An octagon cutting sequence is infinitly derivable.

The converse is not true, but we can describe exactly the condition which one needs to add.

Definition

Let w be infinitly derivable and let $w^{(n)}$ be the n^{th} derived sequence. The sequence $\{s_k\}_{k\in\mathbb{N}} \in \{0, 1, \dots, 7\}^{\mathbb{N}}$ is a sequence of sectors for w if for each k, $w^{(k)}$ gives a path on the diagram \mathscr{D}_{s_k} .

▲□▶ ▲冊▶ ▲臣▶ ▲臣▶ 三臣 - 釣�?

$$w = C \subseteq C \underset{s_0}{\underline{B}} \subset C \underset{D}{\underline{B}} \underset{D}{\underline{D}} \underset{B}{\underline{D}} \underset{B}{\underline{D}} \underset{C}{\underline{B}} \underset{C}{\underline{C}} \underset{B}{\underline{B}} \underset{D}{\underline{C}} \underset{A}{\underline{C}} \underset{C}{\underline{C}} \underset{B}{\underline{C}} \underset{D}{\underline{C}} \underset{A}{\underline{C}} \underset{A}{\underline{C}} \underset{B}{\underline{C}} \underset{A}{\underline{C}} \underset{A}$$

Octagon Continued Fractions

Let $F : [0, \pi] \to [\pi/8, \pi]$ the following map, that we call Octagon Farey map:

Definition

The octagon continued fraction expansion of θ is

 $\theta = [s_0, s_1, s_2, \dots, s_k, \dots]_O \quad \text{iff} \quad \{\theta\} = \cap_k F_{s_0}^{-1} F_{s_1}^{-1} \dots F_{s_k}^{-1} [0, \pi].$

In this case we have $F^k(\theta) \in \Sigma_{s_k}$ per tutti i k.

Octagon Continued Fractions

Let $F : [0, \pi] \to [\pi/8, \pi]$ the following map, that we call Octagon Farey map:

Definition

The octagon continued fraction expansion of θ is

 $\theta = [s_0, s_1, s_2, \dots, s_k, \dots]_O \quad \text{iff} \quad \{\theta\} = \cap_k F_{s_0}^{-1} F_{s_1}^{-1} \dots F_{s_k}^{-1} [0, \pi].$

▲□▶▲□▶▲≡▶▲≡▶ ≡ めぬぐ

In this case we have $F^k(\theta) \in \Sigma_{s_k}$ per tutti i k.

Direction Recognitionn

Let w be an octagon cutting sequence.

Lemma

If w is not a periodic sequence, the sequence of sectors $\{s_k\}_{k\in\mathbb{N}}$ is univoquely determined. In particular, each derivative $w^{(k)}$ is admissible in an unique \mathscr{D}_{s_k} .

Theorem

If w is not periodic, there is a unique sequence of sectors $\{s_k\}_{k\in\mathbb{N}}$ for w and the direction of the trajectories with cutting sequence w is given by

$$\theta = [s_0, s_1, s_2,]_O.$$

Direction Recognitionn

Let w be an octagon cutting sequence.

Lemma

If w is not a periodic sequence, the sequence of sectors $\{s_k\}_{k\in\mathbb{N}}$ is univoquely determined. In particular, each derivative $w^{(k)}$ is admissible in an unique \mathscr{D}_{s_k} .

Theorem

If w is not periodic, there is a unique sequence of sectors $\{s_k\}_{k\in\mathbb{N}}$ for w and the direction of the trajectories with cutting sequence w is given by

$$\boldsymbol{\theta} = [\mathbf{s}_0, \mathbf{s}_1, \mathbf{s}_2, \dots]_O.$$

As in the case of the square, the theorems follow if we prove that:

Theorem

If w is an octagon cutting sequence, also the derived sequence w' is an octagon cutting sequence.

Furthermore, if w is the cutting sequence of a trajectory in direction θ , the derived sequence w' is a cutting sequence of a trajectory in direction $\theta' = F(\theta)$, where F is the octagon Farey map.

To prove it, one uses an argument in the same spirit of:

i.e. one uses renormalization in the space of affine deformations of the octagon.

(日) (四) (日) (日) (日)

As in the case of the square, the theorems follow if we prove that:

Theorem

If w is an octagon cutting sequence, also the derived sequence w' is an octagon cutting sequence.

Furthermore, if w is the cutting sequence of a trajectory in direction θ , the derived sequence w' is a cutting sequence of a trajectory in direction $\theta' = F(\theta)$, where F is the octagon Farey map.

To prove it, one uses an argument in the same spirit of:

i.e. one uses renormalization in the space of affine deformations of the octagon.

As in the case of the square, the theorems follow if we prove that:

Theorem

If w is an octagon cutting sequence, also the derived sequence w' is an octagon cutting sequence.

Furthermore, if w is the cutting sequence of a trajectory in direction θ , the derived sequence w' is a cutting sequence of a trajectory in direction $\theta' = F(\theta)$, where F is the octagon Farey map.

To prove it, one uses an argument in the same spirit of:



i.e. one uses renormalization in the space of affine deformations of the octagon.

・ ロ ト ・ 西 ト ・ 日 ト ・ 日 ト

-

As in the case of the square, the theorems follow if we prove that:

Theorem

If w is an octagon cutting sequence, also the derived sequence w' is an octagon cutting sequence.

Furthermore, if w is the cutting sequence of a trajectory in direction θ , the derived sequence w' is a cutting sequence of a trajectory in direction $\theta' = F(\theta)$, where F is the octagon Farey map.

To prove it, one uses an argument in the same spirit of:



i.e. one uses renormalization in the space of affine deformations of the octagon.

Let S_O be the surface obtained glueing opposite sides of the octagon by translations: S_O is an example of a *translation surface*, i.e. it has an atlas whose changes of coordinates are of the form $z \mapsto z + c$.

Definition

An automorphism $\Psi : S \mapsto S$ of a translation surface S is an affine automorphism if it is affine in each chart and $D\Psi(z)$ is independent on $z \in S$. Let

 $Aff(S) = \{\Psi, \quad \Psi \text{ affine automorphism}\}.$

・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・

Let S_O be the surface obtained glueing opposite sides of the octagon by translations: S_O is an example of a *translation surface*, i.e. it has an atlas whose changes of coordinates are of the form $z \mapsto z + c$.

Definition

An automorphism $\Psi : S \mapsto S$ of a translation surface S is an affine automorphism if it is affine in each chart and $D\Psi(z)$ is independent on $z \in S$. Let

 $Aff(S) = \{\Psi, \quad \Psi \text{ affine automorphism}\}.$

・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・

Let S_O be the surface obtained glueing opposite sides of the octagon by translations: S_O is an example of a *translation surface*, i.e. it has an atlas whose changes of coordinates are of the form $z \mapsto z + c$.

Definition

An automorphism $\Psi : S \mapsto S$ of a translation surface S is an affine automorphism if it is affine in each chart and $D\Psi(z)$ is independent on $z \in S$. Let

 $Aff(S) = \{\Psi, \quad \Psi \text{ affine automorphism}\}.$

・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・

Example (1) Se $\Psi \in D_8$ is an isometry of O, clearly one has $\Psi \in Aff(O)$.

Let S_O be the surface obtained glueing opposite sides of the octagon by translations: S_O is an example of a *translation surface*, i.e. it has an atlas whose changes of coordinates are of the form $z \mapsto z + c$.

Definition

An automorphism $\Psi : S \mapsto S$ of a translation surface S is an affine automorphism if it is affine in each chart and $D\Psi(z)$ is independent on $z \in S$. Let

$$Aff(S) = \{\Psi, \quad \Psi \text{ affine automorphism}\}.$$

・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・

Example (2) The matrix $\begin{pmatrix} 1 & 2(1+\sqrt{2}) \\ 0 & 1 \end{pmatrix} \in V(O).$

The Veech group V(S) is the group of linear parts of Aff(S):

$$V(S) = \{ D\Psi, \quad \Psi \in Aff(S) \} \subset SL(2,\mathbb{R})^{\pm}.$$

$$\begin{array}{ll} \mathsf{Ex} \ 1 & V(\mathbb{T}^2)^+ = & < \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}, & \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} > & = SL(2,\mathbb{Z}). \\ \mathsf{Ex} \ 2 & \end{array}$$

$$V(\mathcal{O}) = \langle \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}, \quad \begin{pmatrix} rac{1}{\sqrt{2}} & rac{1}{\sqrt{2}} \\ -rac{1}{\sqrt{2}} & rac{1}{\sqrt{2}} \end{pmatrix}, \quad \begin{pmatrix} 1 & 2(1+\sqrt{2}) \\ 0 & 1 \end{pmatrix} >$$

If S is a translation surface glued from a regular polygon, V(S) is a lattice in $SL(2,\mathbb{R})^{\pm}$ (Veech)

The Veech group V(S) is the group of linear parts of Aff(S):

$$V(S) = \{ D\Psi, \quad \Psi \in Aff(S) \} \subset SL(2,\mathbb{R})^{\pm}.$$

$$\begin{array}{ll} \mathsf{Ex} \ 1 & V(\mathbb{T}^2)^+ = & < \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}, & \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} > & = SL(2,\mathbb{Z}). \\ \mathsf{Ex} \ 2 & \end{array}$$

$$V(O) = < \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}, \quad \begin{pmatrix} \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ -\frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \end{pmatrix}, \quad \begin{pmatrix} 1 & 2(1+\sqrt{2}) \\ 0 & 1 \end{pmatrix} >$$

If S is a translation surface glued from a regular polygon, V(S) is a lattice in $SL(2,\mathbb{R})^{\pm}$ (Veech)

The Veech group V(S) is the group of linear parts of Aff(S):

$$V(S) = \{ D\Psi, \quad \Psi \in Aff(S) \} \subset SL(2,\mathbb{R})^{\pm}.$$

$$\begin{array}{ll} \mathsf{Ex} \ 1 & V(\mathbb{T}^2)^+ = & < \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}, & \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} > & = SL(2,\mathbb{Z}). \\ \mathsf{Ex} \ 2 & \end{array}$$

$$V(O) = < \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}, \quad \begin{pmatrix} \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ -\frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \end{pmatrix}, \quad \begin{pmatrix} 1 & 2(1+\sqrt{2}) \\ 0 & 1 \end{pmatrix} >$$

If S is a translation surface glued from a regular polygon, V(S) is a lattice in $SL(2,\mathbb{R})^{\pm}$ (Veech)

The Veech group V(S) is the group of linear parts of Aff(S):

$$V(S) = \{ D\Psi, \quad \Psi \in Aff(S) \} \subset SL(2,\mathbb{R})^{\pm}.$$

$$\begin{array}{ll} \mathsf{Ex} \ 1 & V(\mathbb{T}^2)^+ = & < \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}, & \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} > & = SL(2,\mathbb{Z}). \\ \mathsf{Ex} \ 2 & \end{array}$$

$$V(O) = \langle egin{pmatrix} 1 & 0 \ 0 & -1 \end{pmatrix}, \quad egin{pmatrix} rac{1}{\sqrt{2}} & rac{1}{\sqrt{2}} \ -rac{1}{\sqrt{2}} & rac{1}{\sqrt{2}} \end{pmatrix}, \quad egin{pmatrix} 1 & 2(1+\sqrt{2}) \ 0 & 1 \end{pmatrix} >$$

If S is a translation surface glued from a regular polygon, V(S) is a lattice in $SL(2,\mathbb{R})^{\pm}$ (Veech)

The Veech group V(S) is the group of linear parts of Aff(S):

$$V(S) = \{ D\Psi, \quad \Psi \in Aff(S) \} \subset SL(2,\mathbb{R})^{\pm}.$$

$$\begin{array}{ll} \mathsf{Ex} \ 1 & V(\mathbb{T}^2)^+ = & < \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}, & \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} > & = SL(2,\mathbb{Z}). \\ \mathsf{Ex} \ 2 & \end{array}$$

$$V(O) = \langle egin{pmatrix} 1 & 0 \ 0 & -1 \end{pmatrix}, \quad egin{pmatrix} rac{1}{\sqrt{2}} & rac{1}{\sqrt{2}} \ -rac{1}{\sqrt{2}} & rac{1}{\sqrt{2}} \end{pmatrix}, \quad egin{pmatrix} 1 & 2(1+\sqrt{2}) \ 0 & 1 \end{pmatrix} >$$

If S is a translation surface glued from a regular polygon, V(S) is a lattice in $SL(2,\mathbb{R})^{\pm}$ (Veech)

The Veech group V(S) is the group of linear parts of Aff(S):

$$V(S) = \{ D\Psi, \quad \Psi \in Aff(S) \} \subset SL(2,\mathbb{R})^{\pm}.$$

$$\begin{array}{ll} \mathsf{Ex} \ 1 & V(\mathbb{T}^2)^+ = & < \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}, \quad \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} > & = SL(2,\mathbb{Z}). \\ \mathsf{Ex} \ 2 & \end{array}$$

▲□▶ ▲□▶ ▲□▶ ▲□▶ ▲□ ● ● ●

If S is a translation surface glued from a regular polygon, V(S) is a lattice in $SL(2,\mathbb{R})^{\pm}$ (Veech)

The Veech group V(S) is the group of linear parts of Aff(S):

$$V(S) = \{ D\Psi, \quad \Psi \in Aff(S) \} \subset SL(2,\mathbb{R})^{\pm}.$$

$$\begin{array}{ll} \mathsf{Ex} \ 1 & V(\mathbb{T}^2)^+ = & < \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}, & \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} > & = \mathit{SL}(2,\mathbb{Z}). \\ \mathsf{Ex} \ 2 & \end{array}$$

▲□▶ ▲□▶ ▲□▶ ▲□▶ ▲□ ● ● ●

If S is a translation surface glued from a regular polygon, V(S) is a lattice in $SL(2,\mathbb{R})^{\pm}$ (Veech)

Derivation and renormalization in the octagon

Let w be the cutting sequence of a trajectory in direction $\theta \in \Sigma_0$. Let $O' = \begin{pmatrix} -1 & 2(1 + \sqrt{2}) \\ 0 & 1 \end{pmatrix} O.$

Lemma

The derived sequence w' coincides with the cutting sequence of the same trajectory in direction θ with respect to the sides of O'.

Let us renormalize: $O' \mapsto O$ $\theta \mapsto \theta'$

Lemma

The derived sequence w' is an octagon cutting sequence in direction $\theta' = F_O(\theta)$.

Derivation and renormalization in the octagon

Let w be the cutting sequence of a trajectory in direction $\theta \in \Sigma_0$. Let $O' = \begin{pmatrix} -1 & 2(1 + \sqrt{2}) \\ 0 & 1 \end{pmatrix} O.$

Lemma

The derived sequence w' coincides with the cutting sequence of the same trajectory in direction θ with respect to the sides of O'.

▲□▶ ▲□▶ ▲□▶ ▲□▶ □ のQで



Lemma

The derived sequence w' is an octagon cutting sequence in direction $\theta' = F_O(\theta)$.

Derivation and renormalization in the octagon

Let w be the cutting sequence of a trajectory in direction $\theta \in \Sigma_0$. Let $O' = \begin{pmatrix} -1 & 2(1 + \sqrt{2}) \\ 0 & 1 \end{pmatrix} O.$

Lemma

The derived sequence w' coincides with the cutting sequence of the same trajectory in direction θ with respect to the sides of O'.

▲□▶ ▲□▶ ▲□▶ ▲□▶ □ のQで



Lemma

The derived sequence w' is an octagon cutting sequence in direction $\theta' = F_O(\theta)$.








The space of lattices is $\mathbb{H}/SL(2,\mathbb{Z})$. (moduli space of tori with a flat metric)

Farey Tessellation: $V(Q) = < \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}, \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} > \mathscr{T}(0, 1, \infty)$ $S_1 = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}$ $S_2 = \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix}$ $\frac{1}{0}$ Given θ , g_t^{θ} geodesics $\theta = \frac{1}{a_0 + \frac{1}{2}}$



affine deformations



(Teichmueller disk)



 $V(O) \cdot O$ centers of ideal octagons

tree of renormalization moves

Given θ , g_t^{θ} geodesics is approximated by a sequence of renormalization moves

the Farey map of the octagon is given by the action on $\partial \mathbb{D}$.

 $\frac{\text{affine deformations}}{\text{affini automorphisms}} = \frac{SL(2,\mathbb{R})}{V(O)}$

 $\frac{SL(2,\mathbb{R})}{V(O)}$ (Teichmueller disk)



 $V(O) \cdot O$ centers of ideal octagons

tree of renormalization moves

Given θ , g_t^{θ} geodesics is approximated by a sequence of renormalization moves

the Farey map of the octagon is given by the action on $\partial \mathbb{D}$.

 $\frac{\text{affine deformations}}{\text{affini automorphisms}} = \frac{SL(2,\mathbb{R})}{V(O)}$





$V(O) \cdot O$ centers of ideal octagons

tree of renormalization moves

Given θ , g_t^{θ} geodesics is approximated by a sequence of renormalization moves

the Farey map of the octagon is given by the action on $\partial \mathbb{D}$.

▲□▶ ▲□▶ ▲ 三▶ ▲ 三▶ ▲ 三 ● ● ●





$V(O) \cdot O$ centers of ideal octagons

tree of renormalization moves

Given θ , g_t^{θ} geodesics is approximated by a sequence of renormalization moves

the Farey map of the octagon is given by the action on $\partial \mathbb{D}$.

▲□▶ ▲□▶ ▲ 三▶ ▲ 三▶ 三 のへぐ

 $\frac{\text{affine deformations}}{\text{affini automorphisms}} = \frac{SL(2,\mathbb{R})}{V(O)} \quad \text{(Teichmueller disk)}$



 $V(O) \cdot O$ centers of ideal octagons

tree of renormalization moves

Given θ , g_t^{θ} geodesics

is approximated by a sequence of renormalization moves

the Farey map of the octagon is given by the action on $\partial \mathbb{D}$.

▲□▶ ▲□▶ ▲ 三▶ ▲ 三▶ ▲ 三 ● ● ●

 $\frac{\text{affine deformations}}{\text{affini automorphisms}} = \frac{SL(2,\mathbb{R})}{V(O)} \quad (\text{Teichmueller disk})$



 $V(O) \cdot O$ centers of ideal octagons

tree of renormalization moves

Given θ , g_t^{θ} geodesics is approximated by a sequence of renormalization moves

the Farey map of the octagon is given by the action on $\partial \mathbb{D}$.

▲□▶ ▲□▶ ▲ 三▶ ▲ 三▶ 三 のへぐ

 $\frac{\text{affine deformations}}{\text{affini automorphisms}} = \frac{SL(2,\mathbb{R})}{V(O)} \quad (\text{Teichmueller disk})$



 $V(O) \cdot O$ centers of ideal octagons

tree of renormalization moves

Given θ , g_t^{θ} geodesics is approximated by a sequence of renormalization moves

the Farey map of the octagon is given by the action on $\partial \mathbb{D}.$

▲□▶ ▲□▶ ▲ 三▶ ▲ 三▶ ▲ 三 ● ● ●

Let us define operators g_i^j that invert derivation.

The operator g_i^j interpolates a sequence w admissible in \mathcal{D}_j producing a sequence ammissible in \mathcal{D}_i and such that $(g_i^j(w))' = w$.

Definition

Let w be ammissible in \mathcal{D}_0 . The sequence $g_0^2 w$ us obtained interpolating the letters corresponding to vertices in Figure with the words on the arrows:

Example $w = \dots DCABBBAC\dots$ $g_0^3 w = \dots DBC\dots$

Let us define operators g_i^j that invert derivation.

The operator g_i^j interpolates a sequence w admissible in \mathcal{D}_j producing a sequence ammissible in \mathcal{D}_i and such that $(g_i^j(w))' = w$.

Definition

Let *w* be ammissible in \mathcal{D}_0 . The sequence $g_0^2 w$ us obtained interpolating the letters corresponding to vertices in Figure with the words on the arrows:



Example w = ... D C A B B B A C $g_0^3 w$ = ... D B C ...

Let us define operators g_i^j that invert derivation.

The operator g_i^j interpolates a sequence w admissible in \mathcal{D}_j producing a sequence ammissible in \mathcal{D}_i and such that $(g_i^j(w))' = w$.

Definition

Let *w* be ammissible in \mathcal{D}_0 . The sequence $g_0^2 w$ us obtained interpolating the letters corresponding to vertices in Figure with the words on the arrows:





Let us define operators g_i^j that invert derivation.

The operator g_i^j interpolates a sequence w admissible in \mathcal{D}_j producing a sequence ammissible in \mathcal{D}_i and such that $(g_i^j(w))' = w$.

Definition

Let w be ammissible in \mathcal{D}_0 . The sequence $g_0^2 w$ us obtained interpolating the letters corresponding to vertices in Figure with the words on the arrows:



Example $w = \dots D$ CABBBAC... $g_0^3 w = \dots D$ BC...

Let us define operators g_i^j that invert derivation.

The operator g_i^j interpolates a sequence w admissible in \mathcal{D}_j producing a sequence ammissible in \mathcal{D}_i and such that $(g_i^j(w))' = w$.

Definition

Let *w* be ammissible in \mathcal{D}_0 . The sequence $g_0^2 w$ us obtained interpolating the letters corresponding to vertices in Figure with the words on the arrows:



Example $w = \dots D C A B B B A C \dots$ $g_0^3 w = \dots D B C B D A \dots$

Let us define operators g_i^j that invert derivation.

The operator g_i^j interpolates a sequence w admissible in \mathcal{D}_j producing a sequence ammissible in \mathcal{D}_i and such that $(g_i^j(w))' = w$.

Definition

Let *w* be ammissible in \mathcal{D}_0 . The sequence $g_0^2 w$ us obtained interpolating the letters corresponding to vertices in Figure with the words on the arrows:



Example

w = ... D C A B B B A C ... $g_0^3 w$ = ... D B C BD A DBCC B ...

Let us define operators g_i^j that invert derivation.

The operator g_i^j interpolates a sequence w admissible in \mathcal{D}_j producing a sequence ammissible in \mathcal{D}_i and such that $(g_i^j(w))' = w$.

Definition

Let w be ammissible in \mathcal{D}_0 . The sequence $g_0^2 w$ us obtained interpolating the letters corresponding to vertices in Figure with the words on the arrows:



Example

w = ... D C A B B B A C ... $g_0^3 w$ = ... D B C BD A DBCC B CC B CC B CCBD A DB C ...

Let us define operators g_i^j that invert derivation.

The operator g_i^j interpolates a sequence w admissible in \mathcal{D}_j producing a sequence ammissible in \mathcal{D}_i and such that $(g_i^j(w))' = w$.

Definition

Let w be ammissible in \mathcal{D}_0 . The sequence $g_0^2 w$ us obtained interpolating the letters corresponding to vertices in Figure with the words on the arrows:



Example

 $w = \dots D C A B B B A C \dots$ $g_0^3 w = \dots D B C B D A D B C C B C C B C C B C C B D A D B C \dots$

The interpolation operators



The other operators are obtained from these ones by permuting the letters.

▲□▶ ▲□▶ ▲□▶ ▲□▶ □ のQで

Characterization of the closure of cutting sequences

Lemma

If w is a cutting sequence and $\{s_n\}_{n\in\mathbb{N}}$ a sequence of sectors, we have

$$w \in \cap_n g_{s_1}^{s_0} g_{s_2}^{s_1} \dots g_{s_n}^{s_{n-1}} \{A, B, C, D\}^{\mathbb{Z}}.$$

This condition, together with infinite derivability, is *necessary* and *sufficient* to characterize the *closure* of octagon cutting sequences (in $\{A, B, C, D\}^{\mathbb{Z}}$):

Theorem (Smillie-U)

A sequence $w \in \{A, B, C, D\}^{\mathbb{Z}}$ belongs to the closure of cutting sequences in the octagon iff there exists a squence $\{s_n\}_{n\in\mathbb{N}}\in\{0,\ldots,7\}^{\mathbb{N}}$ such that

$$w \in \bigcap_n g_{s_1}^{s_0} g_{s_2}^{s_1} \dots g_{s_n}^{s_{n-1}} \{A, B, C, D\}^{\mathbb{Z}}.$$

An analogous theorem holds for every regular polyon.

Characterization of the closure of cutting sequences

Lemma

If w is a cutting sequence and $\{s_n\}_{n\in\mathbb{N}}$ a sequence of sectors, we have

$$w \in \cap_n g_{s_1}^{s_0} g_{s_2}^{s_1} \dots g_{s_n}^{s_{n-1}} \{A, B, C, D\}^{\mathbb{Z}}.$$

This condition, together with infinite derivability, is *necessary* and *sufficient* to characterize the *closure* of octagon cutting sequences (in $\{A, B, C, D\}^{\mathbb{Z}}$):

Theorem (Smillie-U)

A sequence $w \in \{A, B, C, D\}^{\mathbb{Z}}$ belongs to the closure of cutting sequences in the octagon iff there exists a squence $\{s_n\}_{n\in\mathbb{N}}\in\{0,\ldots,7\}^{\mathbb{N}}$ such that

$$w \in \bigcap_n g_{s_1}^{s_0} g_{s_2}^{s_1} \dots g_{s_n}^{s_{n-1}} \{A, B, C, D\}^{\mathbb{Z}}.$$

An analogous theorem holds for every regular polyon.

Characterization of the closure of cutting sequences

Lemma

If w is a cutting sequence and $\{s_n\}_{n\in\mathbb{N}}$ a sequence of sectors, we have

$$w \in \cap_n g_{s_1}^{s_0} g_{s_2}^{s_1} \dots g_{s_n}^{s_{n-1}} \{A, B, C, D\}^{\mathbb{Z}}.$$

This condition, together with infinite derivability, is *necessary* and *sufficient* to characterize the *closure* of octagon cutting sequences (in $\{A, B, C, D\}^{\mathbb{Z}}$):

Theorem (Smillie-U)

A sequence $w \in \{A, B, C, D\}^{\mathbb{Z}}$ belongs to the closure of cutting sequences in the octagon iff there exists a squence $\{s_n\}_{n\in\mathbb{N}}\in\{0,\ldots,7\}^{\mathbb{N}}$ such that

$$w \in \cap_n g_{s_1}^{s_0} g_{s_2}^{s_1} \dots g_{s_n}^{s_{n-1}} \{A, B, C, D\}^{\mathbb{Z}}.$$

・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・

An analogous theorem holds for every regular polyon.