

SHOT-NOISE SPATIAL BIRTH-and-DEATH PROCESSES

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Structure of the Lecture

- Background and Motivation
- Methodology
- **Part 1**
Birth and Death Wireless Queues
Joint work with A. Sankararaman and S. Foss
- **Part 2**
Wireless Queues under Multihop Routing and Motion
Joint work with S. Rybko, S. Shlosman and S. Vladimirov

Motivations in Wireless Networks

■ Lack of understanding and analysis of

Space-time interactions

- Static spatial setting well understood: Stochastic Geometry [FB, Blaszczyszyn 01]
- Churn partly taken into account in flow-based queuing [Bonald, Proutiere 06], [Shakkottai, De Veciana 07]

■ Contents of this lecture:

First models with such dynamics in stochastic geometry

Methodology

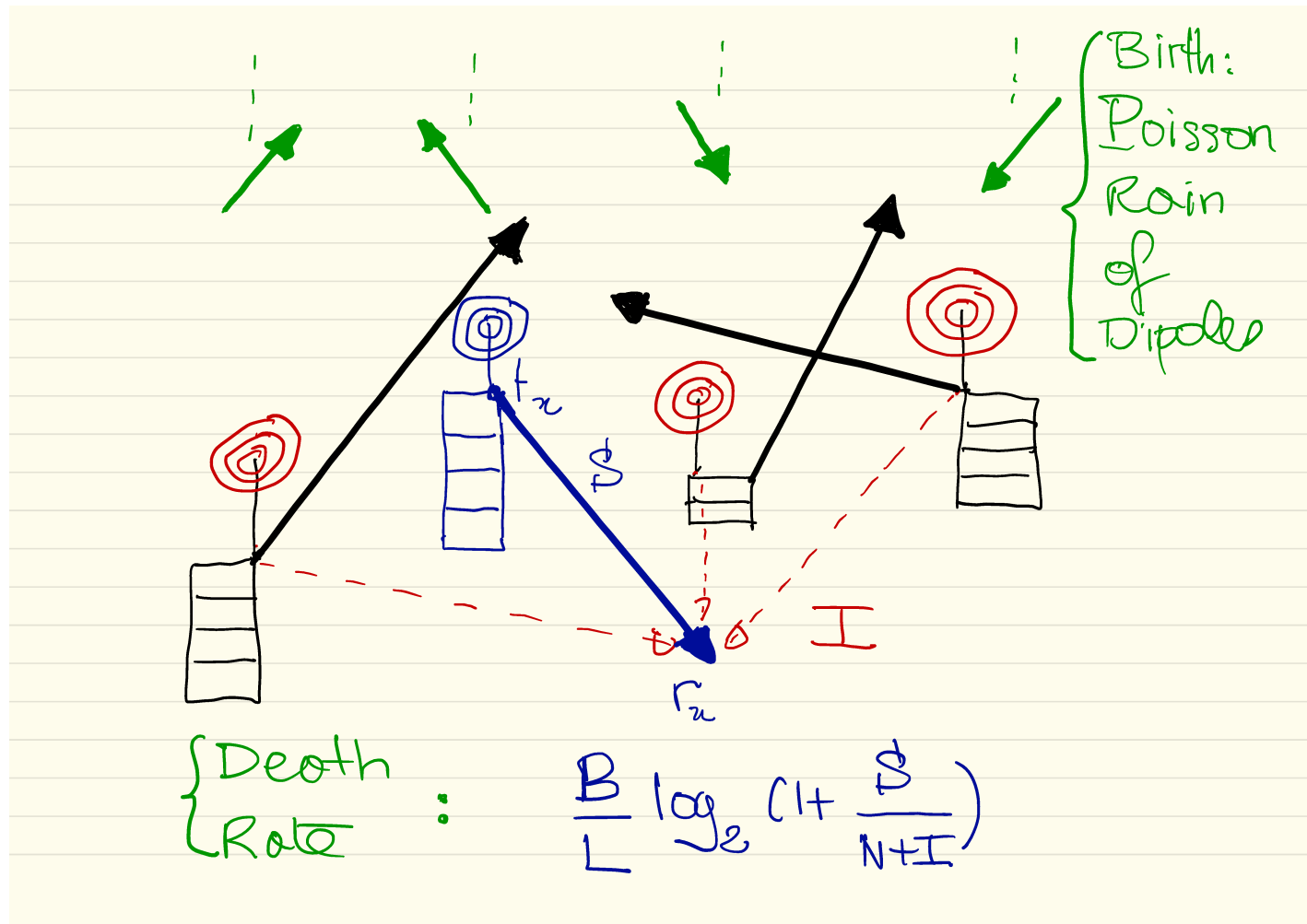
Everything Should Be Made as Simple as Possible,
But Not Simpler
A. E.

Setting

- **Infrastructureless Wireless Network:**
Ad-hoc Networks, D2D Networks, IoT
- **Markov Models:**
Poisson, Exponential
- **Mathematical tools:**
Stochastic Geometry, Fluid, Mean-Field

I: Wireless Birth and Death

- Problem Statement
- Summary of Results
- Proof Overview



Stochastic Network Model

- $S = [-Q, Q] \times [-Q, Q]$: torus where the wireless links live
- **Links**: (Tx-Rx pairs)
- **Links**: **arrive** as a PPP on $\mathbb{R} \times S$ with intensity λ :
Prob. of a point arriving in space dx and time dt : $\lambda dx dt$
- Each Tx has an **i.i.d. exponential file size**
of mean **L** bits to transmit to its Rx
- A point **exits** after the Tx finishes transmitting its file
- Φ_t : set of locations of links present at time t :

$$\Phi_t = \{\mathbf{x}_1, \dots, \mathbf{x}_{N_t}\}, \quad \mathbf{x}_i \in S$$

Interference and Service Rate

- Interference seen at point x due to configuration Φ

$$I(\mathbf{x}, \Phi) = \sum_{\mathbf{x}_i \in \Phi \neq \mathbf{x}} l(\|\mathbf{x} - \mathbf{x}_i\|)$$

- Distance on the torus
- $l(\cdot): \mathbb{R}^+ \rightarrow \mathbb{R}^+$: path loss function

- The **speed of file transfer** by link at $\mathbf{x}$ in configuration Φ is

$$R(\mathbf{x}, \Phi) = B \log_2 \left(1 + \frac{1}{N + I(\mathbf{x}, \Phi)} \right)$$

- **B, N** Positive constants

B& D Master Equation

- A point born at $\mathbf{x}_p$ and time b_p with file-size L_p dies at time

$$d_p = \inf \left\{ t > b_p : \int_{u=b_p}^t R(\mathbf{x}_p, \Phi_u) du \geq L_p \right\}$$

- **Spatial Birth-Death Process**

- Arrivals from the Poisson Rain
- Departures happen at file transfer completion

Properties of the Dynamics

- The statistical assumptions imply that Φ_t is a Markov Process on the set of simple counting measures on $\mathbf{S}$
- **Euclidean extension** of the flow-level models of [Bonald, Proutiere 06], [Shakkottai, De Veciana 07]

Questions

- Existence and uniqueness of the stationary regimes of Φ_t
- Characterization of the stationary regime(s) if existence

Main Stability Results

$$a := \int_{x \in S} l(\|x\|) dx$$

■ Theorem

- If $\lambda > \frac{B}{\ln(2)La}$, then Φ_t admits no stationary regime.
- If $\lambda < \frac{B}{\ln(2)La}$, and $r \rightarrow l(r)$ bounded and monotone, then Φ_t admits a unique stationary regime

■ Necessary condition by Palm calculus

■ Sufficient condition by fluid limit

■ Corollary

For the path-loss model $l(r) = r^{-\alpha}$, $\alpha \geq 2$, for all $\lambda > 0$, and all mean file sizes, the process Φ_t admits no stationary-regime

Main Qualitative Result

- Φ stationary point-process on S with Palm distribution $\mathbb{P}^0$
- **Clustering**
 Φ is clustered if for all bounded, positive, non-increasing functions $f(\cdot) : \mathbb{R}^+ \rightarrow \mathbb{R}^+$, the shot noise;

$$\mathbf{F}(\mathbf{x}, \Phi) := \sum_{\mathbf{y} \in \Phi \setminus \{\mathbf{x}\}} f(\|\mathbf{y} - \mathbf{x}\|)$$

satisfies

$$\mathbb{E}^0[\mathbf{F}(\mathbf{0}, \Phi)] \geq \mathbb{E}[\mathbf{F}(\mathbf{0}, \Phi)]$$

Main Qualitative Result (*continued*)

■ **Theorem**

The steady-state point process, when it exists, is **clustered**

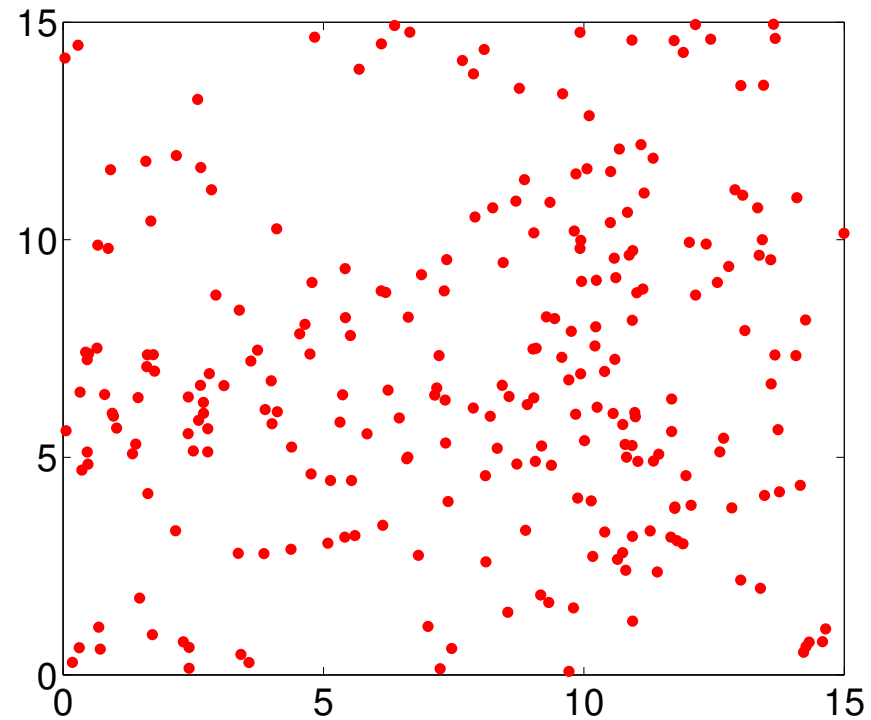
■ Follows from **Palm calculus + the FKG inequality**

■ **Interpretation of the result**

The steady-state interference measured at a uniformly randomly chosen point of is larger in mean than that at an uniformly random location of space.

■ **Key Observation**

- Dynamics Shapes Geometry
- Geometry Shapes Dynamics



A sample of Φ when $\lambda = 0.99$ and $l(\mathbf{r}) = (\mathbf{r} + \mathbf{1})^{-4}$.

First Quantitative Results

- Mean-field approximations for the intensity of the steady-state process
 1. Poisson heuristic β_f - derived by neglecting clustering and assuming Poisson
 2. Second-order heuristic β_s based on a second-order cavity approximation of the dynamics

Poisson Heuristic

Exact Rate Conservation Law:

$$\lambda \mathbf{L} = \beta \mathbb{E}_{\Phi}^0 \left[\log_2 \left(1 + \frac{1}{\mathbf{N} + \mathbf{I}(\mathbf{0})} \right) \right].$$

Poisson Heur.: Largest solution to the fixed point equation:

$$\lambda \mathbf{L} = \frac{\beta_f}{\ln(2)} \int_{\mathbf{z}=0}^{\infty} \frac{e^{-\mathbf{N}\mathbf{z}}(1 - e^{-\mathbf{z}})}{\mathbf{z}} e^{-\beta_f \int_{\mathbf{x} \in \mathbf{S}} (1 - e^{-\mathbf{z}l(\|\mathbf{x}\|)}) d\mathbf{x}} d\mathbf{z}$$

Ignores the Palm effect and uses that if $\mathbf{X}, \mathbf{Y}$ are non-negative and independent,

$$\mathbb{E} \left[\ln \left(1 + \frac{\mathbf{X}}{\mathbf{Y} + \mathbf{a}} \right) \right] = \int_{\mathbf{z}=0}^{\infty} \frac{e^{-\mathbf{a}\mathbf{z}}}{\mathbf{z}} (1 - \mathbb{E}[e^{-\mathbf{z}\mathbf{X}}]) \mathbb{E}[e^{-\mathbf{z}\mathbf{Y}}] d\mathbf{z}.$$

Second Order Heuristic

The intensity β_s is given by

$$\beta_s = \frac{\lambda L}{B \log_2 \left(1 + \frac{1}{N+I_s} \right)}$$

where I_s is the smallest solution of the fixed-point equation

$$I_s = \lambda L \int_{\mathbf{x} \in S} \frac{I(||\mathbf{x}||)}{B \log_2 \left(1 + \frac{1}{N+I_s+I(||\mathbf{x}||)} \right)} d\mathbf{x}$$

Second Order Heuristic (*continued*)

- **Rationale** based on $\rho_2(\mathbf{x}, \mathbf{y})$: second moment measure of Φ
- **Rate Conservation for ρ_2** : when considering I_s as a constant

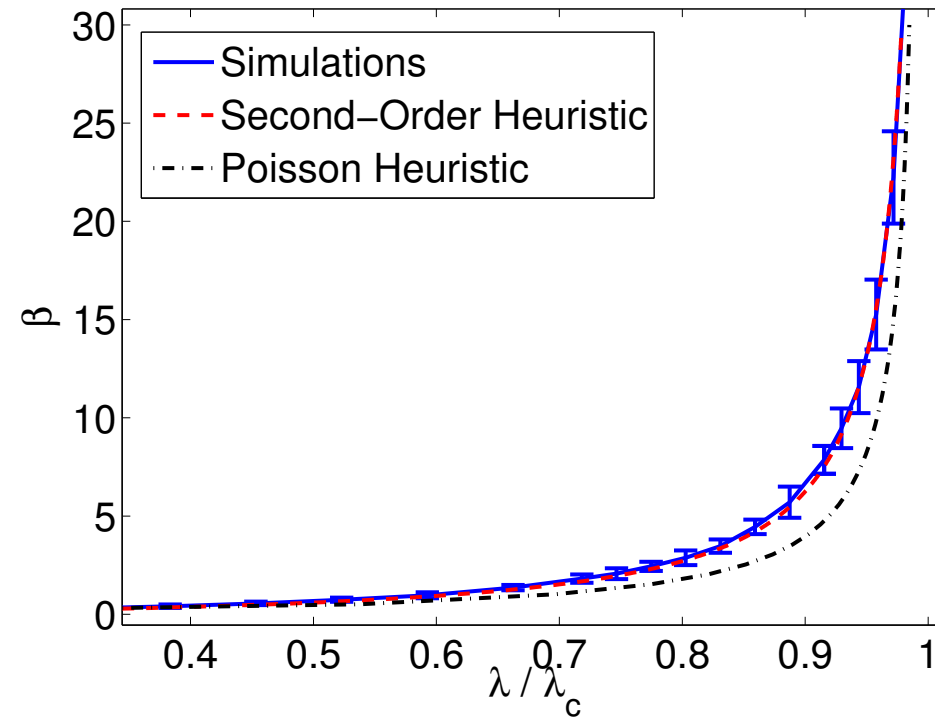
$$\rho_2(\mathbf{x}, \mathbf{y}) \frac{1}{L} B \log_2 \left(1 + \frac{1}{N + I_s + l(\|\mathbf{x} - \mathbf{y}\|)} \right) = \lambda \beta_s$$

- From the definition of second moment measure,

$$I_s = \int_{\mathbf{x} \in S} l(\|\mathbf{x}\|) \frac{\rho_2(\mathbf{0}, \mathbf{x})}{\beta_s} d\mathbf{x}$$

which gives the fixed point equation for I_s

- The formula for β_s follows from **Rate Conservation for $\rho_1 = \beta_s$**



95% confidence interval when $l(r) = (r + 1)^{-4}$

Tightness Results & Extensions

- The Poisson heuristic is **tight** in heavy and light traffic
- **Recent Extensions** obtained with **S. Foss**:
 - **Exact expression** for the intensity β of Φ in the **Low SINR** regime when replacing the death rate by

$$\frac{B}{\ln(2)L} \frac{S}{N + I(\mathbf{x}, \Phi)}$$
 - **Scalability result**: extension to dynamics on $\mathbb{R}^2$ using **Coupling from the Past** techniques.
- **Future**: introduction of **scheduling** or **multi-user IT**

Summary

- A new basic representation of **space time interactions** in wireless networks
- A **generative model for clustering** as assumed in 3GPP simulation standards
- A new **dynamic notion of capacity** involving both queuing and IT
- **First analytical results** in the Low SINR case and good **heuristics** in general

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Motion of Wireless Queues under Multihop Routing
Joint work with S. Rybko, S. Shlosman and S. Vladimirov

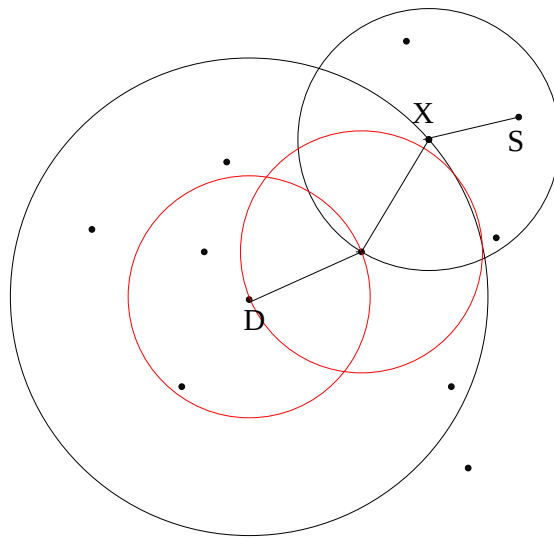
II: Motion and Multi-Hop Routing

- Problem Statement
- Summary of Results

Problem Statement

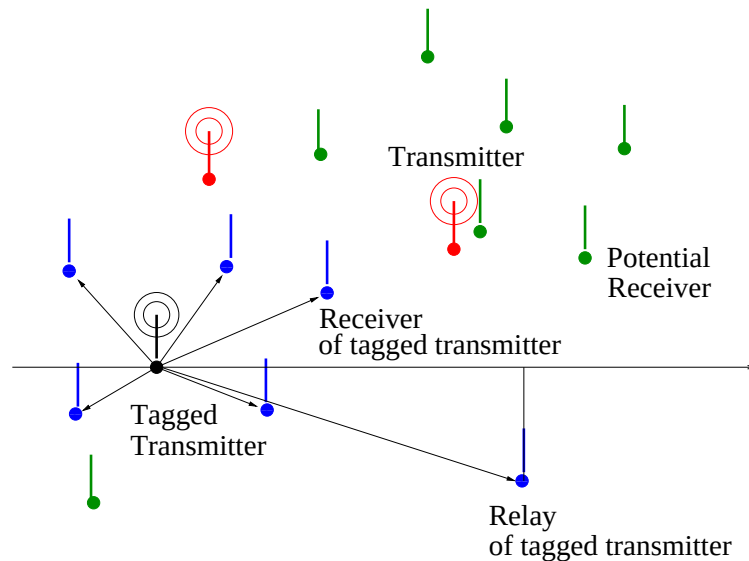
- **Setting:** Grossglauser & Tse 02 scaling law problem
 - Multihop relaying
 - Opportunistic geographic routing
 - Motion of nodes
- New **SG+QT** view of the problem

Example of Geographic Routing



- **Nearest Neighbor Geographic Routing**
The next hop on the route from S to D is the nearest among the nodes which are closer from D than X .
- On a Poisson P.P., a.s.
 - No ties
 - Converges in finite number of steps

Wireless Geographic Routing



- Each node uses **Aloha** to split the Poisson p.p. into transmitters and potential receivers
- **Potential relays of a transmitter:** receivers with a large enough SINR
- **Geographic Routing:** next hop := potential relay nearest to destination

Traffic and Relaying

- **DTN** type assumptions:
 - Wireless nodes **move**
 - Each moving node generates **packets**
 - Each generated packet has a **destination**, e.g. another node at finite distance
 - **Multihop relaying**: each node transmits
 - * its own packets
 - * those of other nodes on their route to destination
 - **Contention**: on each node, packets are queued **FIFO** and are served as above (SINR condition)

Reduction (as per Methodology :)

- Wireless nodes move randomly on a grid or a graph G (e.g. $\mathbb{Z}$ or $\mathbb{Z}/K\mathbb{Z}$, $\mathbb{Z}^2$, d -regular graph)
- **Traffic:**
 - Each moving node generates packets at rate λ
 - Each generated packet has a destination (e.g. a point of the grid, vertex of the graph)
- **Contention:** on each node packets are queued **FIFO**

Reduction (as per Methodology :) (*continued*)

- **Motion:** neighbor nodes swap their positions with rate β
- **Wireless:** communication to neighbors only
- **Opportunistic multihop routing:**
upon service at a node, a packet:
 - is routed to the neighbor the **closest to its destination**
 - leaves the system if the destination is distance 1 or 0

Finite Network Markovization

■ Assumptions

- Poisson arrivals with intensity λ
- exponential service times with mean 1
- finite connected graph with K nodes

■ Markov representation with discrete non compact state:

- Permutation on $[1, \dots, K]$
(locations of wireless nodes)
- Ordered queue at each node
(finite ordered list of destinations)

Finite Network Instability Result

- Maximal degree in the graph: d
- Motion rate β
- **Theorem**

For all $\beta > 0$, this Markov process is transient when

$$K > K^* = \frac{d+1}{\lambda}$$

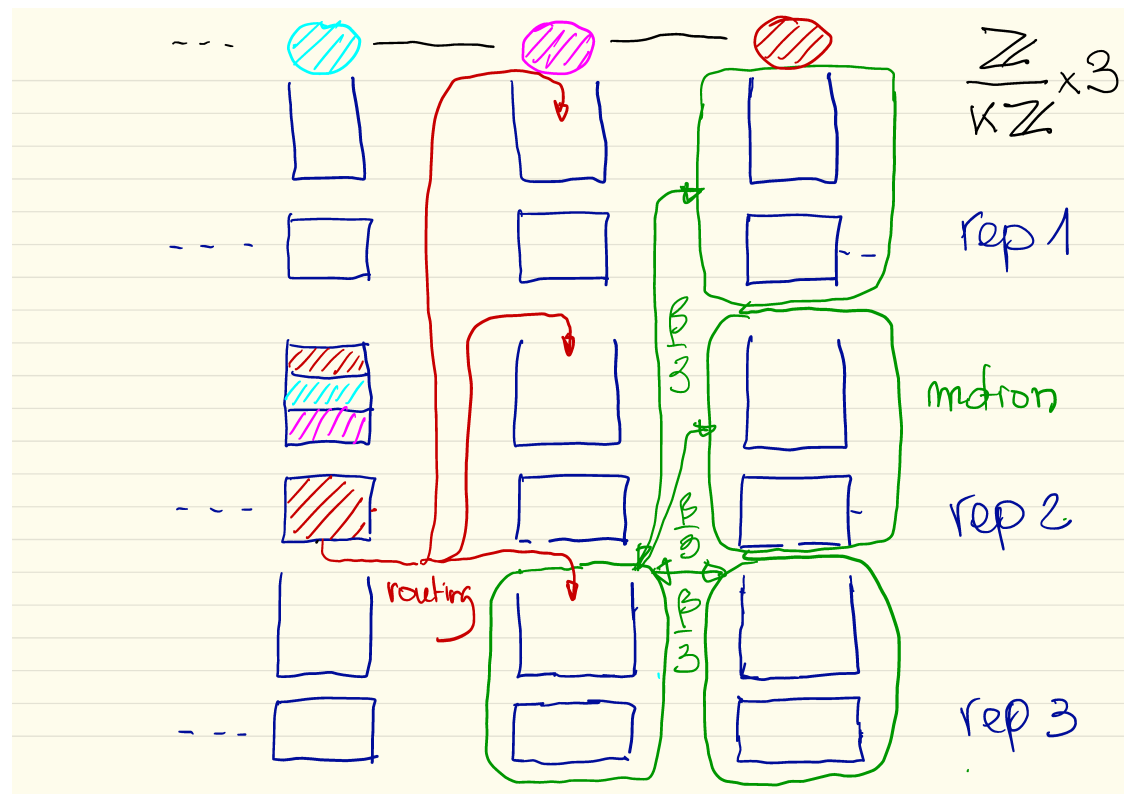
- For all load factors, a finite network is unstable when the diameter of the network is large enough!

■ Sketch of Proof in the d -regular case

- nodes mix to the uniform distribution on $[1, \dots, K]$
- the proportion of time the server harboring a packet is a neighbor of its destination vertex is order $\frac{d+1}{K}$.
- if $\lambda > \frac{d+1}{K}$, drift inside orthant is positive on all coordinates
- supermartingale argument

Replicated Version

- **N-Replica version** of the network on G
- New graph G_N with
 - $V_N = V \times \{1, \dots, N\}$
 - edge between (u, i) and (v, j) if (u, v) edge in G
- **Routing**
 - destination is any replica of the destination vertex in G
 - routing to one of the N replicas closest to destination at random
- **Swaping**: between a replica and any neighboring replica at random



The 3-Replica Version of $\mathbb{Z}/K\mathbb{Z}$

Instability of the N -Replica Version of a Finite Network

- Same instability result for G_N as for G for fixed N
- **Example** The N -replica version of the network on $\frac{\mathbb{Z}}{K\mathbb{Z}}$ is unstable under the same condition as the network on $\frac{\mathbb{Z}}{K\mathbb{Z}}$ namely if $K > 3/\lambda$

Mean-Field Version

- **Mean-Field version** of the network on G :
weak limit of the latter when N tends to infinity.
- **Notation**: network on G^∞
- **Existence**: tightness arguments

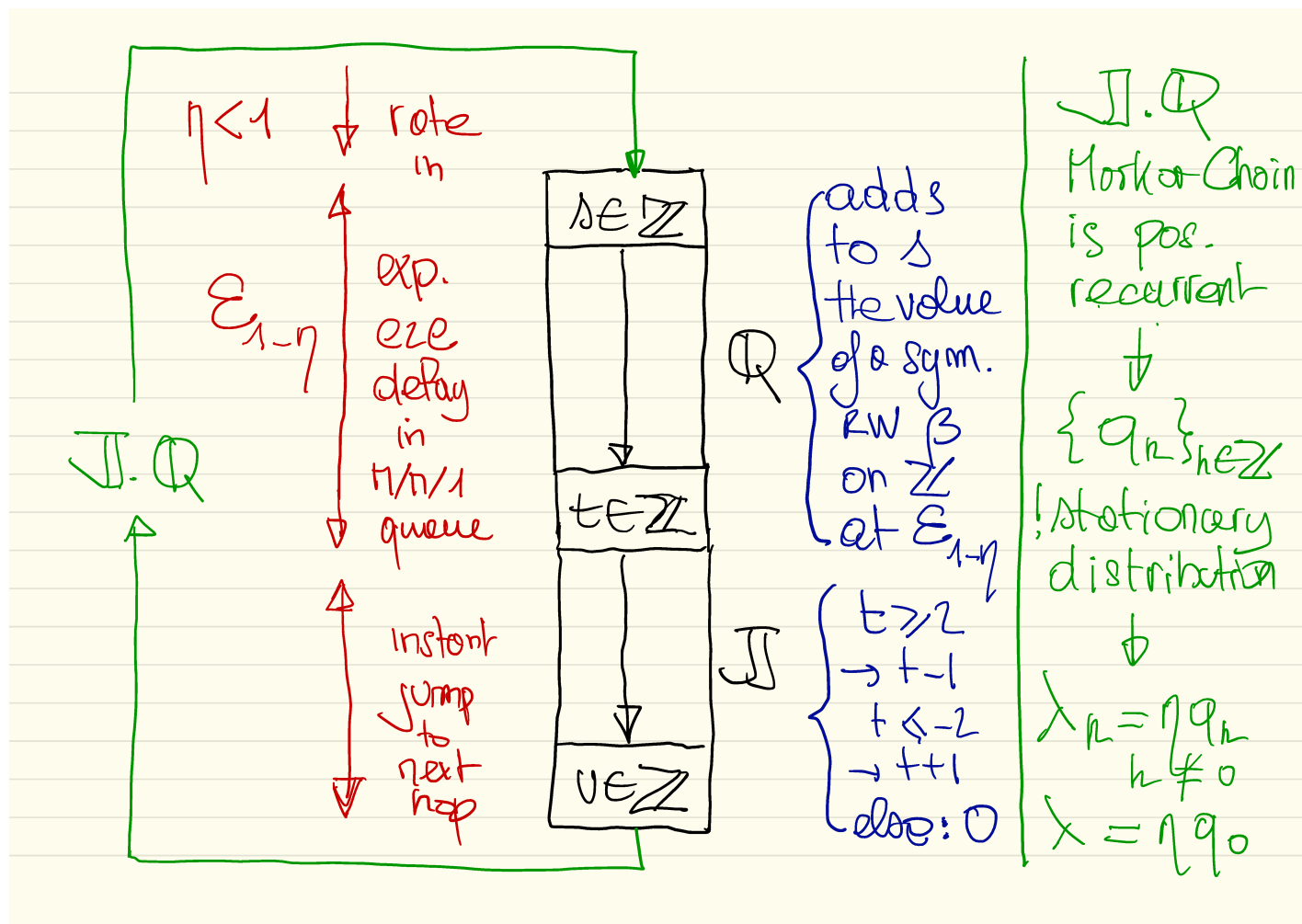
Mean Field Networks on $\mathbb{Z}$

- **Non Linear Markov Process** roughly, dynamical system on probability measures μ on queue states

$$\mu(\mathbf{q}) = \mu(\mathbf{n}_1, \dots, \mathbf{n}_l), \quad \mathbf{n}_k : \text{relative location of dest}(\mathbf{c}_k)$$

- **Functional equation** for fixed points $\mu(\mathbf{q})$ of this dynamical system

- **Theorem** For all $0 < \eta < 1$, there exists
 - a unique $0 < \lambda = \lambda_\eta < 1$
 - a unique probability distribution $\mu = \mu_\eta$ on the queue state such that, for the exogenous arrival λ ,
 - μ is solution of the functional equation and
 - the total rate in a node under μ is η
- **Sketch of Proof**
 - Special case where the destination is the vertex of birth



Existence of Multiple Solutions

- **Theorem** For the mean-field version of the network on $\mathbb{Z}$, there exists a λ_* such that for all $\lambda < \lambda_*$, there are at least two different values $\eta = \eta_-(\lambda)$ and $\eta = \eta_+(\lambda)$ s.t.

- $\lambda(\eta) = \lambda$

- $\eta_-(\lambda) \rightarrow 0$ as $\lambda \rightarrow 0$

- $\eta_+(\lambda) \rightarrow 1$ as $\lambda \rightarrow 0$

- **Sketch of Proof**

- When η tends to 0, $\lambda(\eta) = \lambda q_0$ tends to 0 by M/M/1

- When η tends to 1, $\lambda(\eta) = \lambda q_0$ tends to 0 by M/M/1 as well

Illustration

- MC of the **distance to destination**:

$$\mathbf{p}(\mathbf{n}, \mathbf{n} + \mathbf{1}) = \beta, \quad \mathbf{p}(\mathbf{n}, \mathbf{n} - \mathbf{1}) = \beta + \gamma,$$

with $\gamma = 1 - \eta$,

$$\mathbf{p}(\mathbf{1}, \mathbf{2}) = \beta, \quad \mathbf{p}(\mathbf{1}, \mathbf{0}) = \beta, \quad \mathbf{p}(\mathbf{1}, *) = \gamma$$

and

$$\mathbf{p}(\mathbf{0}, \mathbf{1}) = 2\beta, \quad \mathbf{p}(\mathbf{0}, *) = \gamma,$$

where $*$ is absorbing

Illustration (continued)

- Mean # nodes visited by a customer till absorption:

$$\mathbb{E}[\mathbf{N}] = 1 + \frac{2\beta^2}{\gamma(3\beta + \gamma)}$$

- Flow equation:

$$\eta = 1 - \gamma = \lambda \left(1 + \frac{2\beta^2}{\gamma(3\beta + \gamma)} \right)$$

When β is large, there are two roots

$$\eta^- = \lambda + \frac{1 - \lambda - \sqrt{(1 - \lambda)^2 - \frac{8}{3}\lambda\beta}}{2}$$

$$\eta^+ = \lambda + \frac{1 - \lambda + \sqrt{(1 - \lambda)^2 - \frac{8}{3}\lambda\beta}}{2}$$

with

$$0 < \eta^- < \eta^+ < 1, \quad \nu^- \rightarrow 0, \quad \nu^+ \rightarrow 1, \quad \text{when } \lambda \rightarrow 0$$

Generalization

- The multiple fixed point result can be extended to the Cayley graph of any group G s.t.
 - G has a finite generating set
 - G is infinite
 - the arrival rate, swap rate, swap rule, destination rule are G -invariant
- **Example:** the network of $(\mathbb{Z})^\infty$ has at least to fixed points stationary regimes for every $\lambda < \lambda_*$ with $\lambda_* > 0$.

Meta-Stability

- Finitely many replicas–Infinitely many replicas stability difference.
- No contradiction with the fact that, for $N < \infty$, the network of $(\mathbb{Z})^N$ has no stationary regimes
the time - replica diagram does not commute here!

Current and Future Extensions

- **Wireless primitives:** (SINR, MAC) to assess transmission between nearby nodes
- **Motion primitives:** Brownian, random waypoint
- **Queuing primitives:** beyond Poisson arrivals and exponential service times
- **Scheduling:** beyond FIFO: needed to see why motion increases capacity

Summary

- **Way to go** but first step of a mean-field representation of **mutihop routing** in wireless networks beyond scaling laws
- **Meta-stability** is encouraging news
- When fully interconnected with SG, new **analytical handle** for optimization in this class of problems.

References

■ Part 1

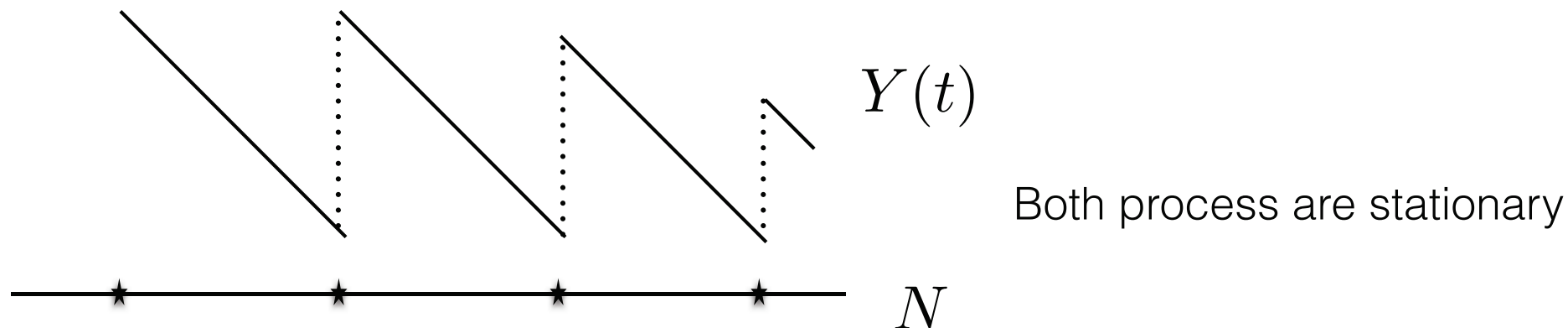
- with [A. Sankararaman](#),
Spatial birth-death wireless networks,
[ArXiv 1604.07884](#), *IEEE Tr. IT*, 63(6) 2017
- with [S. Foss](#) and [A. Sankararaman](#),
Infinite spatial birth-death wireless networks,
in preparation

■ Part 2

- with [S. Rybko](#), [S. Shlosman](#) and [S. Vladimirov](#),
Metastability of Queuing Networks with Mobile Servers,
[ArXiv 1704.02521](#), submitted

Proof Sketch - Necessary Condition

Assume stability and write down 'Rate Conservation Equations'.
Then find a contradiction.



$$\text{If } Y(t) = Y(0) + \int_{s=0}^t D(s)ds + \int_{s=0}^t (Y(s) - Y(s^-))N(ds)$$

$$\text{Implies } \mathbb{E}[D(0)] + \lambda_N \mathbb{E}_N^0[Y(0) - Y(0^-)] = 0$$

Proof Sketch - Necessary Condition

$\phi_t(\mathbf{S})$

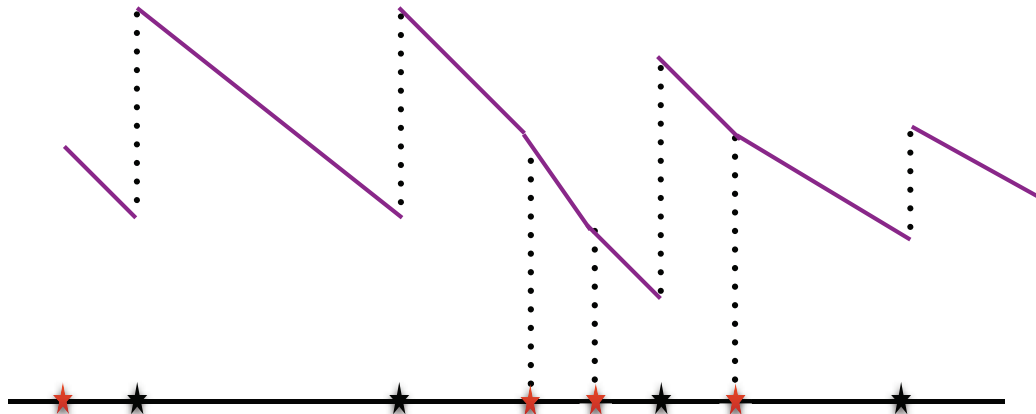


Red - Epochs of Death
Black - Epochs of Arrivals



RCL implies $\lambda|S| = \lambda_d$

Total bits left in the network i.e. remaining 'workload'



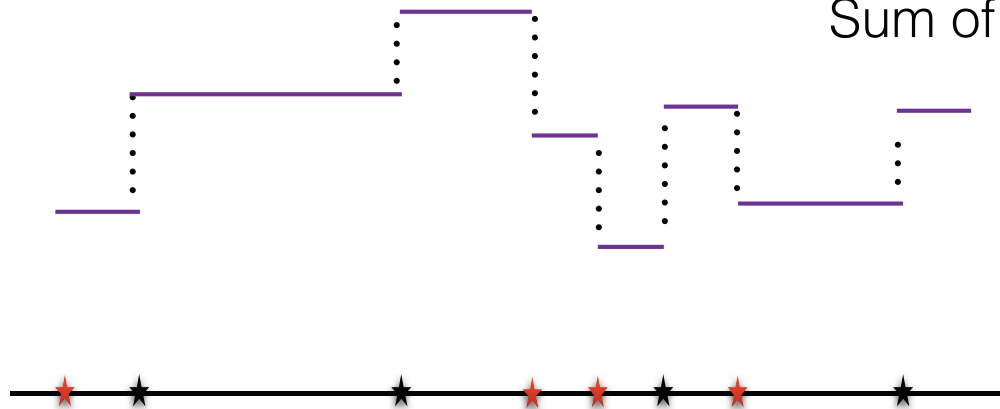
RCL implies

$$\lambda|S|L = \mathbb{E} \left[\sum_{x \in \phi_0} R(x, \phi_0) \right]$$

Proof Sketch - Necessary Condition

Sum of interference seen at all points

$$Y(t) = \sum_{x \in \phi_t} I(x, \phi_t)$$



Red - Epochs of Death
Black - Epochs of Arrivals

RCL implies

$$\lambda |S| \mathbb{E}_B^0[\mathcal{I}] = \lambda_d \mathbb{E}_D^0[\mathcal{D}]$$

PASTA and RCL for $\phi_t(\mathbf{S})$

$$\mathbb{E}[\mathcal{I}] = \mathbb{E}_D^0[\mathcal{D}]$$

Campbell's Theorem

$$\mathbb{E}[\mathcal{I}] = 2 \frac{\mathbb{E}[\phi_0(\mathbf{S})]}{|S|} \int_{x \in \mathbf{S}} l(\|x\|) dx$$

Handle $\mathbb{E}_D^0[\mathcal{D}]$ through Papangelou's Stochastic Intensity formula

Proof Sketch - Necessary Condition

We have
$$\mathbb{E}[\mathcal{I}] = \mathbb{E}_D^0[\mathcal{D}] = 2 \frac{\mathbb{E}[\phi_0(\mathbf{S})]}{|S|} \int_{x \in \mathbf{S}} l(||x||) dx$$

The Death Point process admits as stochastic intensity - $\mathbf{R}_t = \sum_{x \in \phi_t} R(x, \phi_t)$
with respect to the filtration $\mathcal{F}_t = \sigma(\phi_s : s \leq t)$

Papangelou's theorem implies
$$\frac{d\mathbb{P}_D^0}{d\mathbb{P}} \Big|_{\mathcal{F}_{0-}} = \frac{\mathbf{R}_0}{\mathbb{E}[\mathbf{R}_0]}$$

This gives

$$\mathbb{E}_D^0[\mathcal{D}] = \mathbb{E} \left[\frac{\mathbf{R}_0}{\mathbb{E}[\mathbf{R}_0]} \mathcal{D} \right] = 2\mathbb{E} \left[\frac{\mathbf{R}_0}{\mathbb{E}[\mathbf{R}_0]} \sum_{x \in \phi_0} \frac{R(x, \phi_0)}{\mathbf{R}_0} I(x, \phi_0) \right]$$

Proof Sketch - Necessary Condition

We have $\mathbb{E}[\mathcal{I}] = \mathbb{E}_D^0[\mathcal{D}] = 2 \frac{\mathbb{E}[\phi_0(\mathbf{S})]}{|S|} \int_{x \in \mathbf{S}} l(||x||) dx$

and

$$\mathbb{E}_D^0[\mathcal{D}] = \mathbb{E} \left[\frac{\mathbf{R}_0}{\mathbb{E}[\mathbf{R}_0]} \mathcal{D} \right] = 2 \mathbb{E} \left[\frac{\mathbf{R}_0}{\mathbb{E}[\mathbf{R}_0]} \sum_{x \in \phi_0} \frac{R(x, \phi_0)}{\mathbf{R}_0} I(x, \phi_0) \right]$$

Algebra -

$$2 \frac{\mathbb{E}[\phi_0(\mathbf{S})]}{|S|} \int_{x \in \mathbf{S}} l(||x||) dx = 2 \frac{\mathbb{E}_{\phi_0}^0 [R(0, \phi_0) I(\phi_0)] \mathbb{E}[\phi_0(\mathbf{S})]}{\lambda L |S|}$$
$$\int_{x \in \mathbf{S}} l(||x||) dx = \frac{\mathbb{E}_{\phi_0}^0 [R(0, \phi_0) I(0, \phi_0)]}{\lambda L}$$

Noticing that $R(x, \phi) I(x, \phi) \leq C \log_2(e)$ yields the necessary condition !

“Clustering” in Steady State

Assuming there is a stationary regime,

$$\lambda|S|L = \mathbb{E} \left[\sum_{x \in \phi_0} R(x, \phi_0) \right] = \mathbb{E}_{\phi_0}^0 [R(0, \phi_0)] \beta |S|$$

From RCL arguments

$$\int_{x \in \mathbf{S}} l(\|x\|) dx = \frac{\mathbb{E}_{\phi_0}^0 [R(0, \phi_0) I(0, \phi_0)]}{\lambda L}$$

Negative Association yields $\mathbb{E}_{\phi_0}^0 [R(0, \phi_0) I(0, \phi_0)] \leq \mathbb{E}_{\phi_0}^0 [R(0, \phi_0)] \mathbb{E}_{\phi_0}^0 [I(0, \phi_0)]$

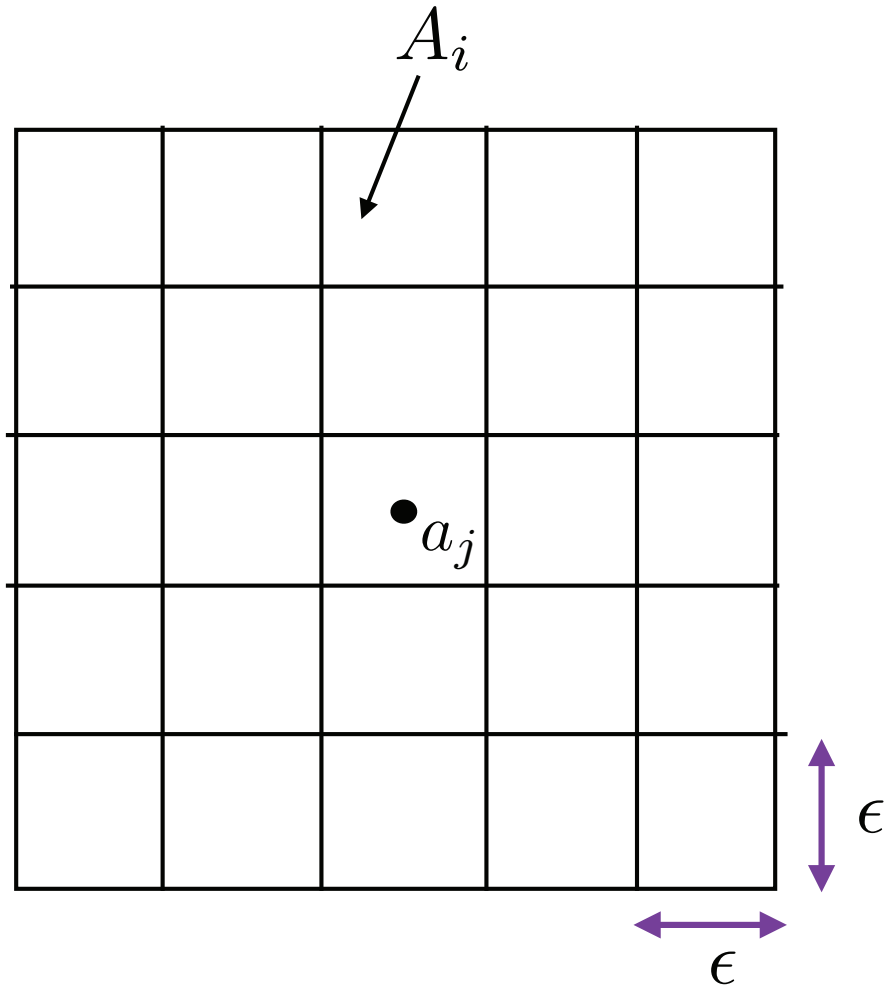
Putting the above together $\mathbb{E}[I(0, \phi_0)] \leq \mathbb{E}_{\phi_0}^0 [I(0, \phi_0)]$

Implies Clustering if path-loss is non-increasing and implies repulsion if path-loss is non-decreasing

Proof Sketch - Sufficient Condition

Consider an *approximate* birth-death process on the ϵ width lattice which is easy to study.

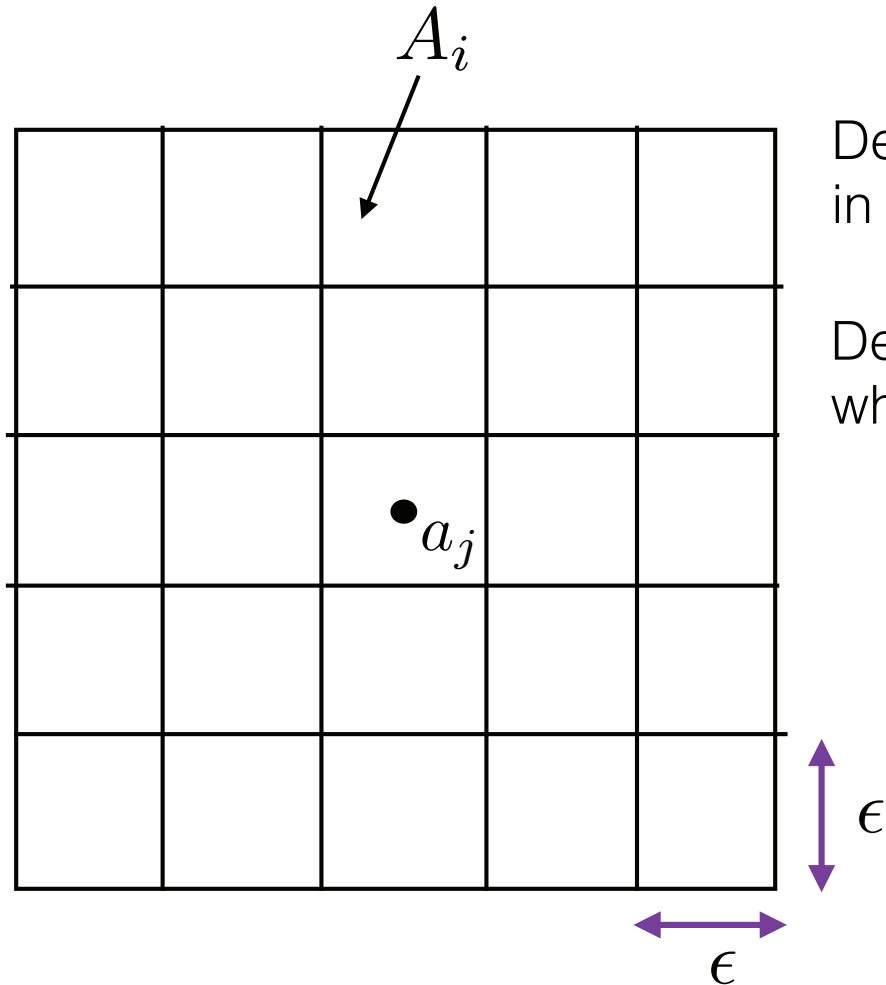
Tessellate into grids of side length at-most ϵ which results in N_ϵ grids.



Proof Sketch - Sufficient Condition

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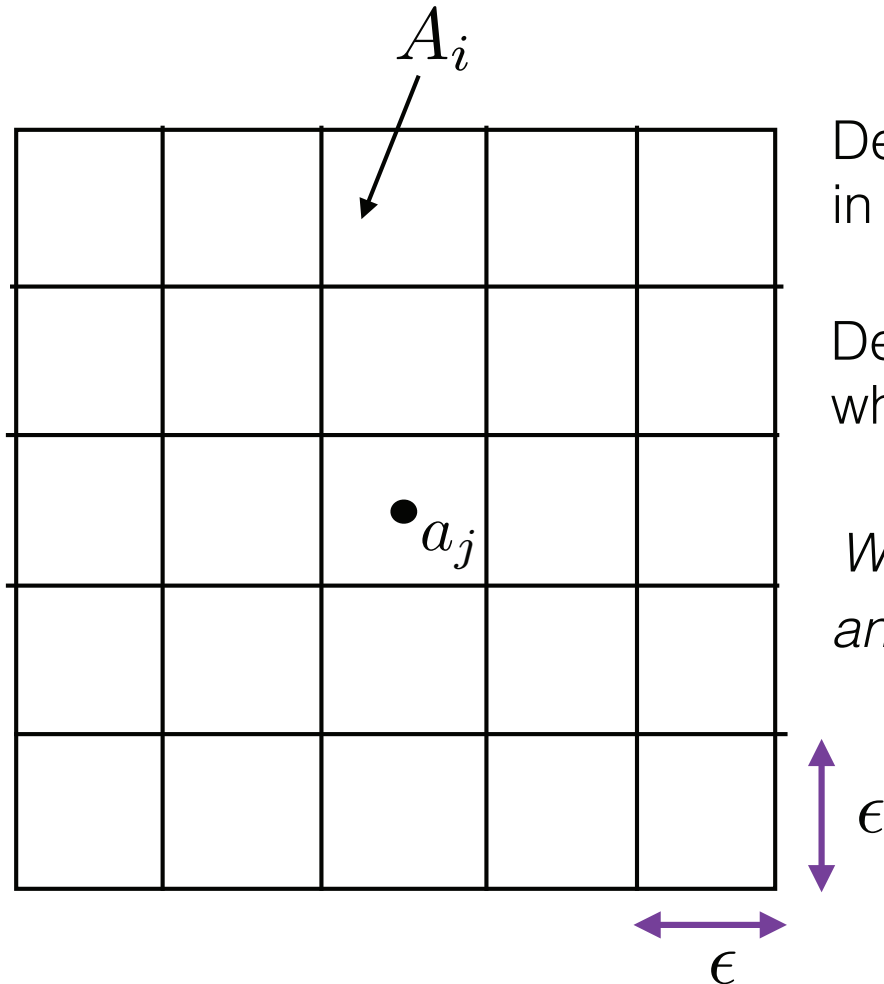
Denote by ϕ_t^ϵ as the configuration at time t in this ϵ approximate system

Denote by the vector $\mathbf{X}(t) \in \mathbb{N}^{N_\epsilon}$ where $X_i(t) = \phi_t^\epsilon(A_i)$.

Proof Sketch - Sufficient Condition

Consider an *approximate* birth-death process on the ϵ width lattice which is easy to study.

Tessellate into grids of side length at-most ϵ which results in N_ϵ grids.



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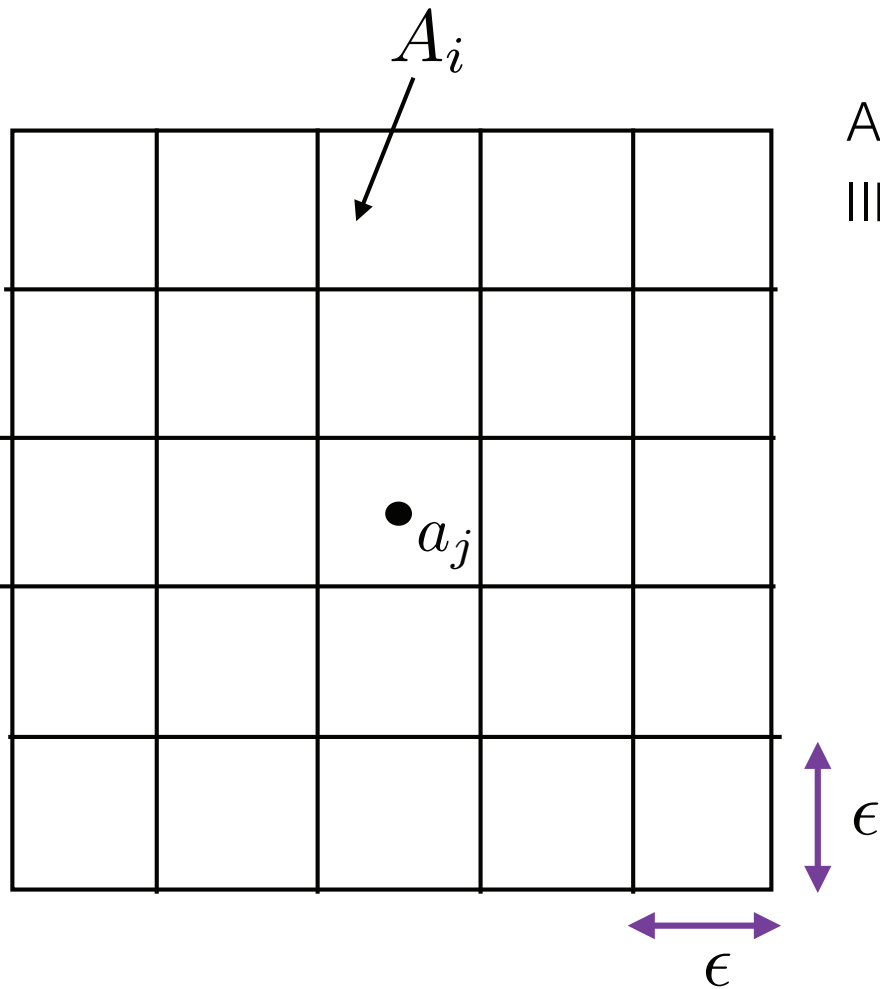
Want $\mathbf{X}(t)$ as a Markov Chain on $\mathbb{N}^{N_\epsilon}$ and want to work out a natural coupling with ϕ_t

Proof Sketch - Sufficient Condition

$$X_i(t) = \phi_t^\epsilon(A_i)$$

*Want $\mathbf{X}(t)$ as a Markov Chain on $\mathbb{N}^{N_\epsilon}$
and want to work out a natural coupling with ϕ_t*

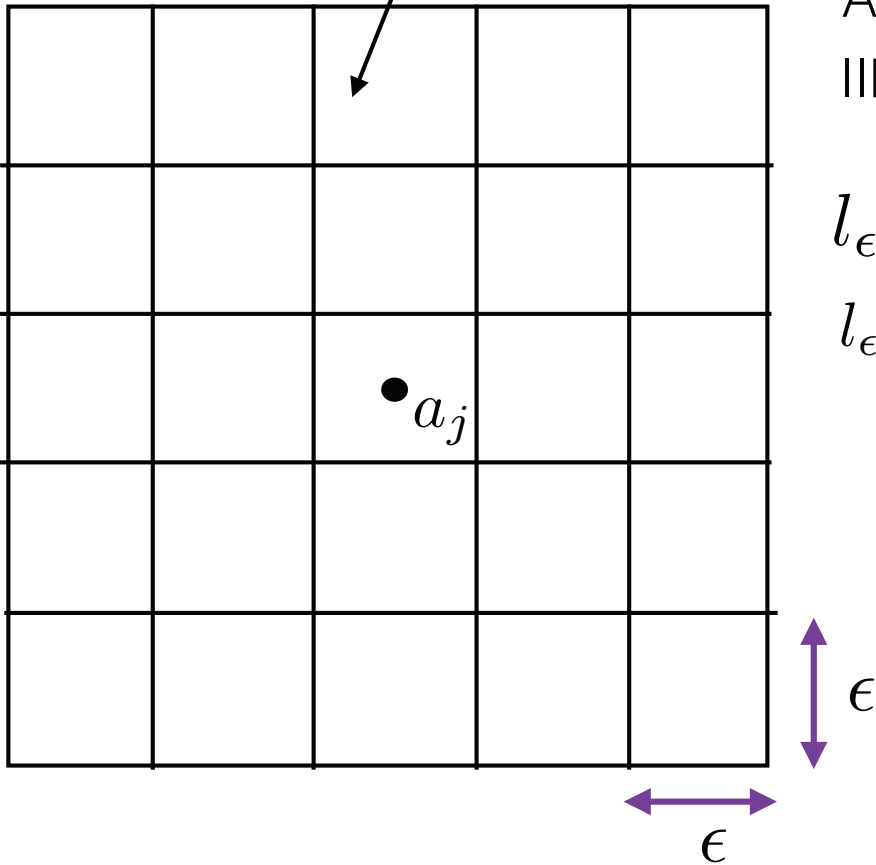
Arrivals - PPP on $\mathbb{R} \times \mathbf{S}$ with intensity λ
IID exponential File Sizes of mean L



Proof Sketch - Sufficient Condition

$$X_i(t) = \phi_t^\epsilon(A_i)$$

A_i



Want $\mathbf{X}(t)$ as a Markov Chain on $\mathbb{N}^{N_\epsilon}$
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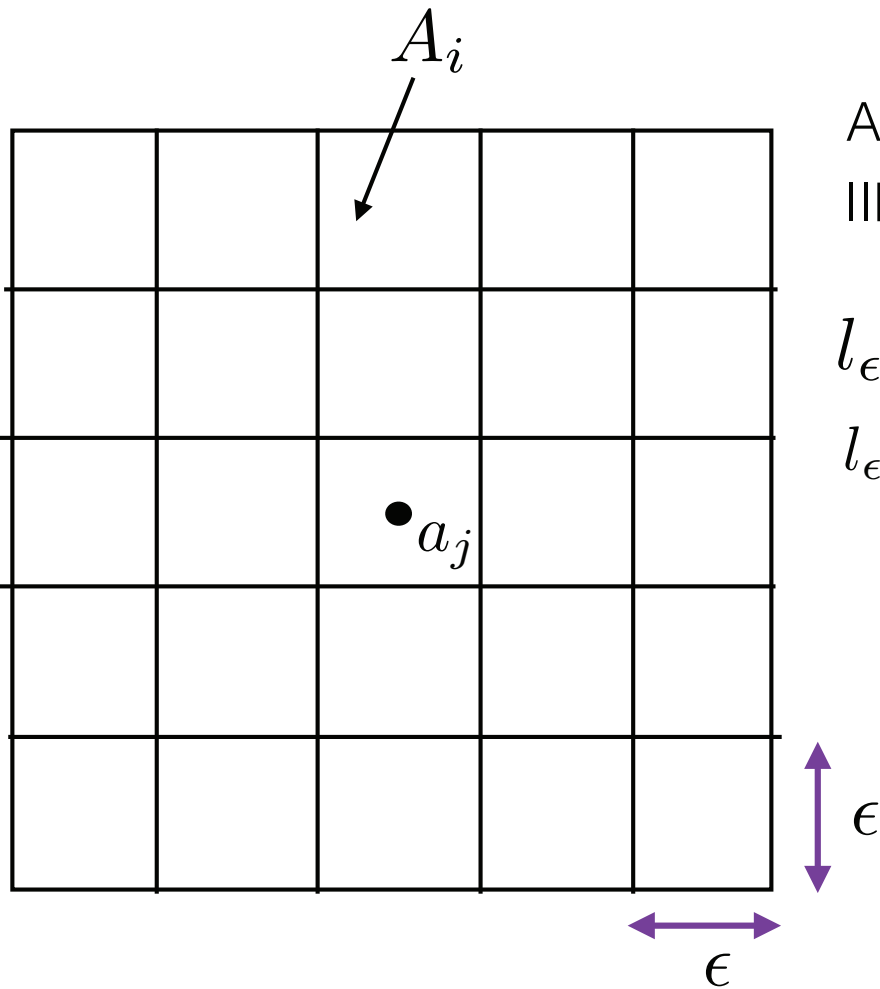
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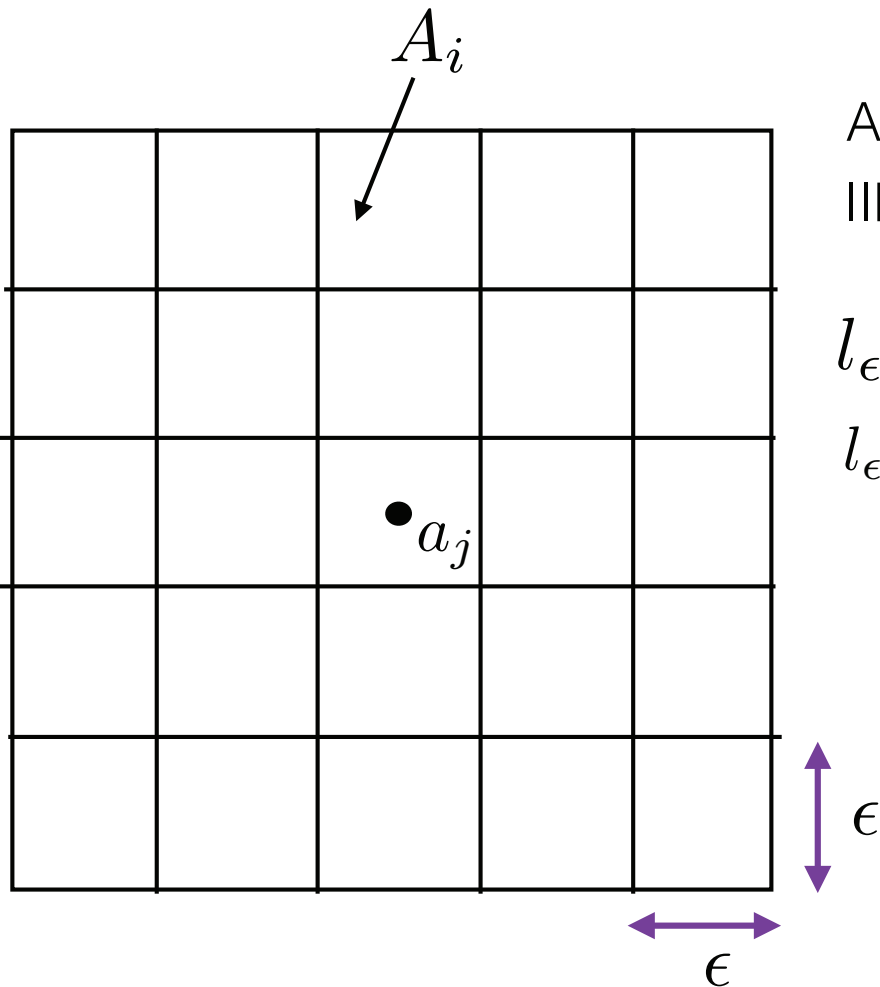
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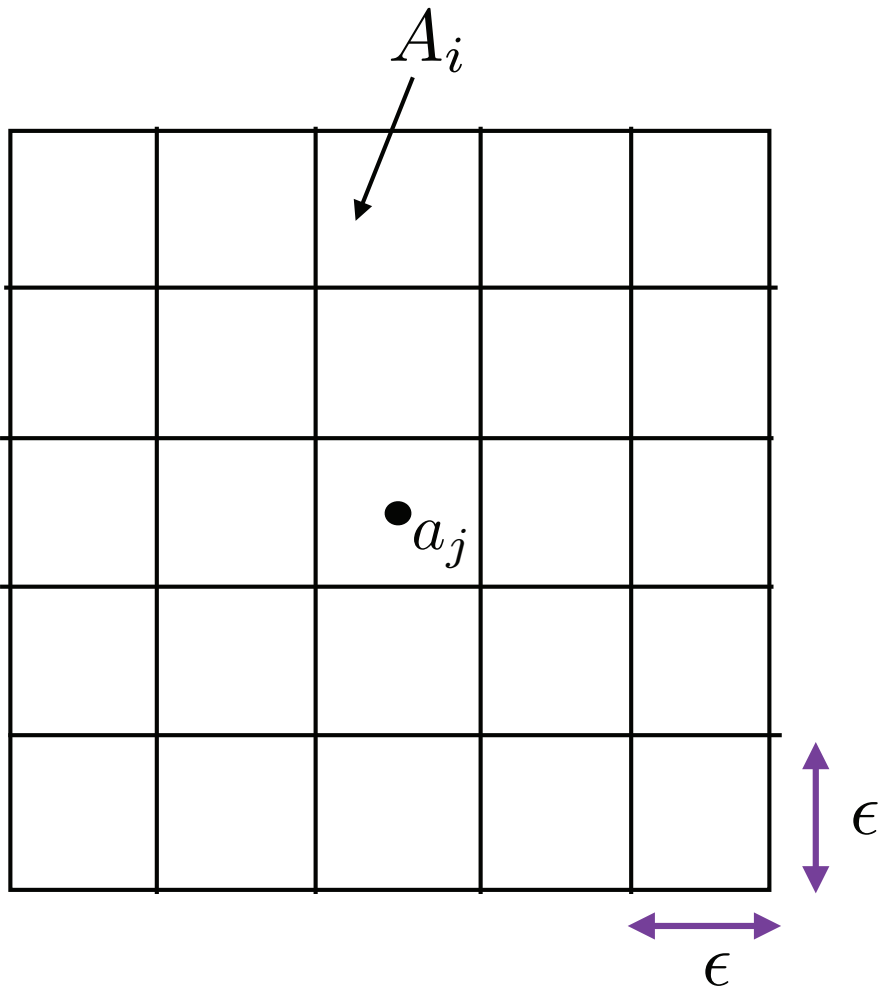
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Need monotonicity of $l(r)$!!

Proof Sketch - Sufficient Condition

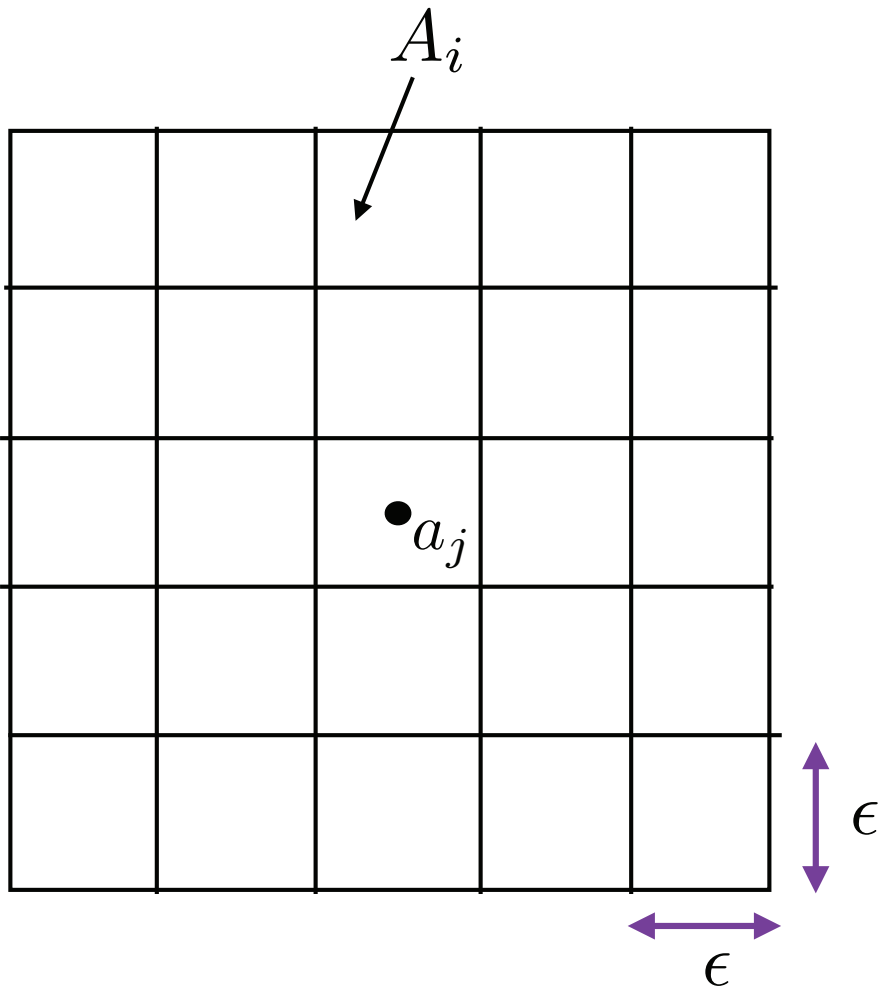


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Hence for a given λ if ϕ_t^{ϵ} is stable, then ϕ_t is stable for that λ

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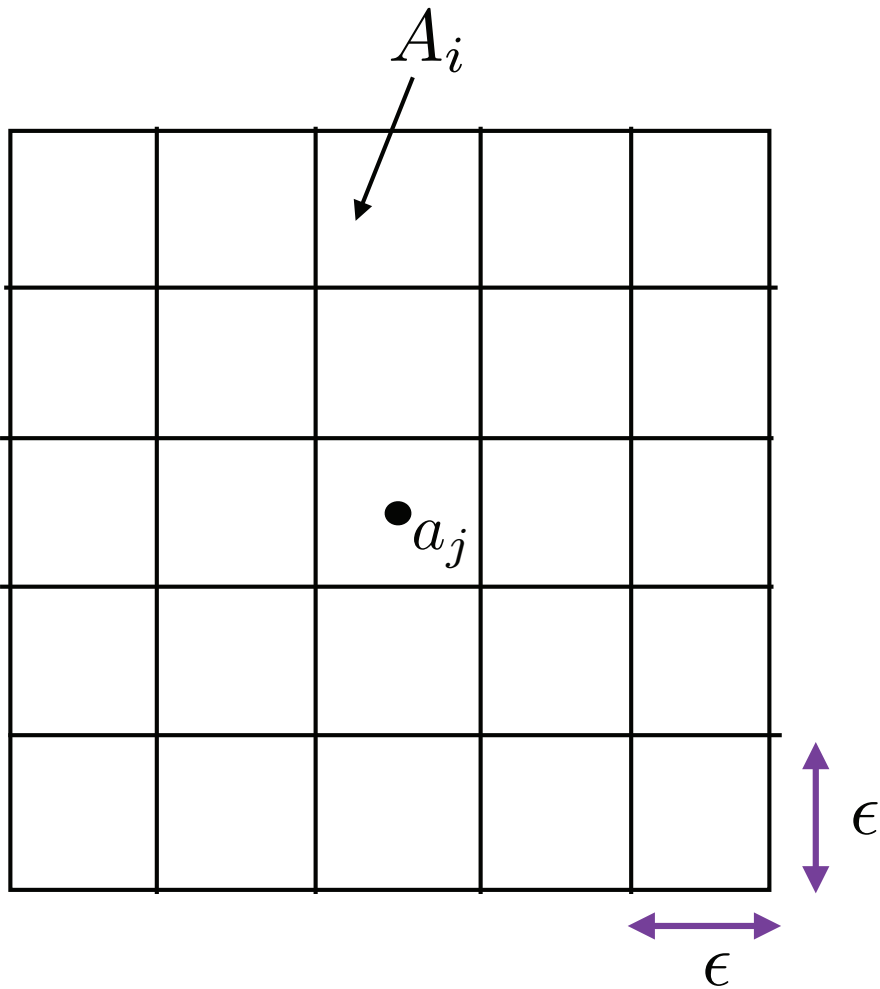
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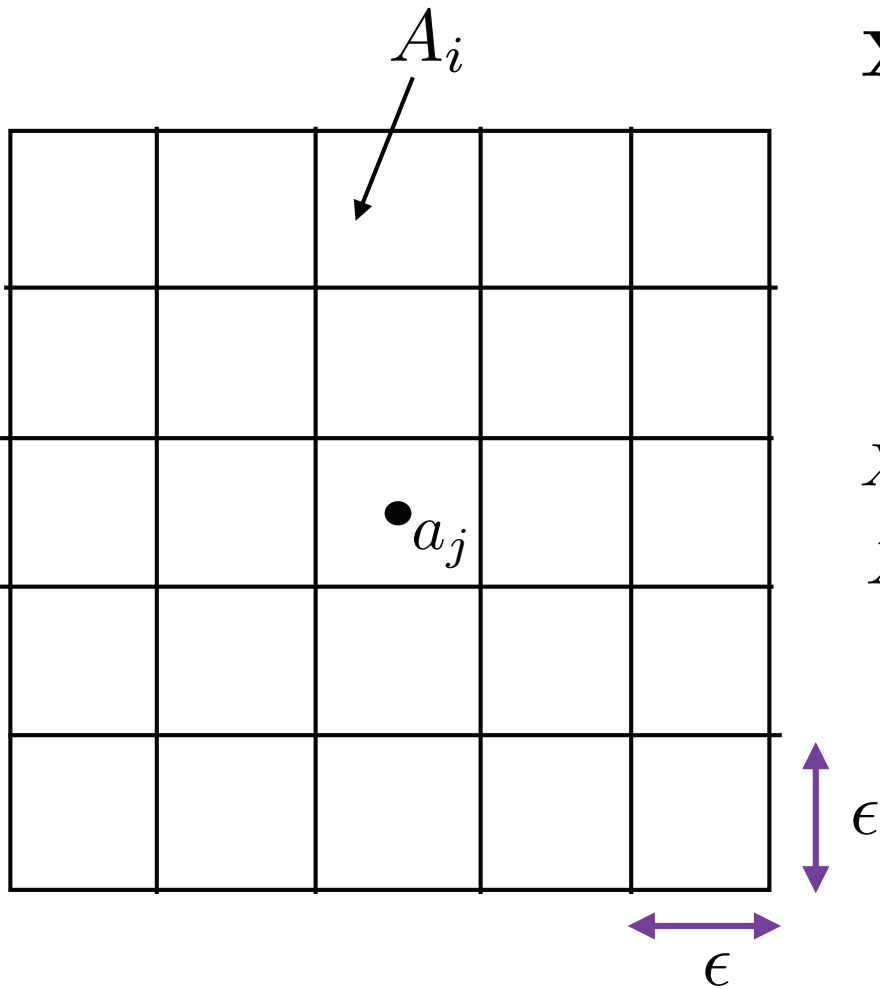
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Obtaining the best possible bound (by optimizing over ϵ) gives that

$$\liminf_{\epsilon \downarrow 0} \lambda L \int_{x \in \mathbf{S}} l_\epsilon(x, 0) dx < \log_2(e) \quad \text{as the stability region of } \phi_t$$

Proof Sketch - Sufficient Condition



$$\mathbf{X}(t) \in \mathbb{N}^{N_\epsilon} \quad \text{with} \quad X_i(t) = \phi_t^\epsilon(A_i)$$

ϕ_t^ϵ is ergodic if and only if $\mathbf{X}(t)$ is ergodic

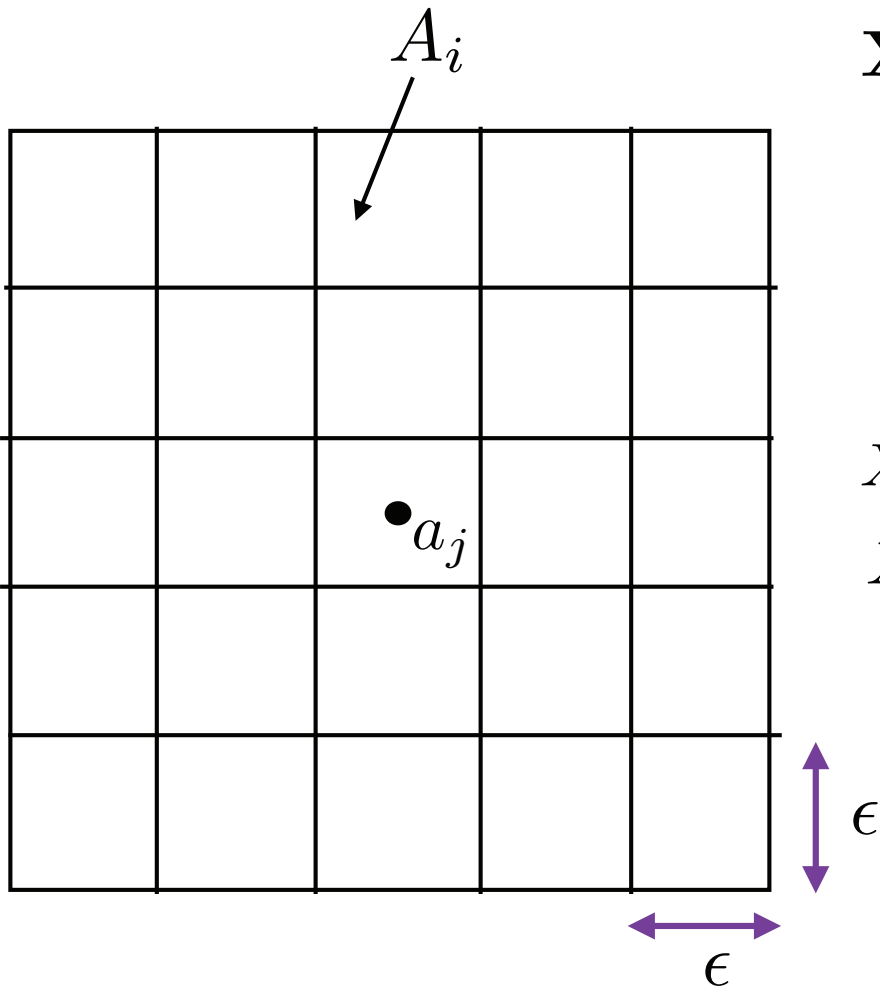
The Evolution

$$X_i \rightarrow X_i + 1 \quad \text{at rate} \quad \lambda \epsilon^2$$

$$X_i \rightarrow X_i - 1 \quad \text{at rate} \quad \frac{1}{L} C X_i \log_2 \left(1 + \frac{1}{N_0 + I_i^\epsilon(X)} \right)$$

$$I_i^\epsilon(X) = \sum_{j=1}^{N_\epsilon} (X_j - \mathbf{1}(j = i)) l_\epsilon(a_i, a_j)$$

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Analyze this evolution through Fluid Limit techniques of [Dai 95], [Massoulié 07].

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The Fluid model of the above evolution

$$x_i(t) := x_i(0) + \lambda\epsilon^2 t - D_i(t)$$

with the derivative of the interference satisfying $\dot{D}_i(t) = \frac{x_i(t)}{I_i^{\epsilon,f}(t)}$

Interference in the fluid scale $I_i^{\epsilon,f}(t) = \sum_k x_k(t) l_\epsilon(a_k, a_i)$

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Or equivalently

$$\frac{d}{dt} x_i(t) = \lambda\epsilon^2 - \frac{x_i(t) C \log_2(e)}{L \sum_k x_k(t) l_\epsilon(a_k, a_i)}$$

The fluid model arises as a result of appropriate space-time scaling

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More precisely, one can show that for a sequence of initial conditions $\{X^{(k)}(0)\}_{k \geq 1}$ and sequence of numbers $\{z_k\}_{k \geq 1}$ such that $z_k \rightarrow \infty$ and $\lim_{k \rightarrow \infty} \frac{X^{(k)}(0)}{z_k} = x(0)$, one has $\frac{X^{(k)}(z_k t)}{z_k} \xrightarrow{\mathbb{P}} x(t)$ u.o.c i.e. $\forall \epsilon > 0$ and $\forall T \in (0, \infty)$

$$\lim_{k \rightarrow \infty} \mathbb{P} \left(\inf_{f \in S(x(0))} \sup_{t \in [0, T]} |z_k^{-1} X^{(k)}(z_k t) - f(t)| > \epsilon \right) = 0$$

Need $l(r)$ to be bounded

Idea borrowed from [Massoulié, 07]

Proof Sketch - Sufficient Condition

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