

# Connectivity and rate with physical layer security over boundaries

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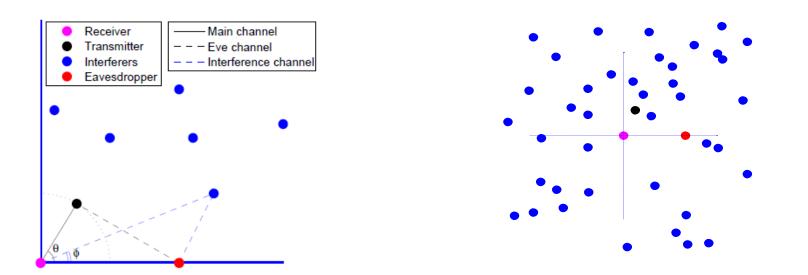


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## **Motivation**

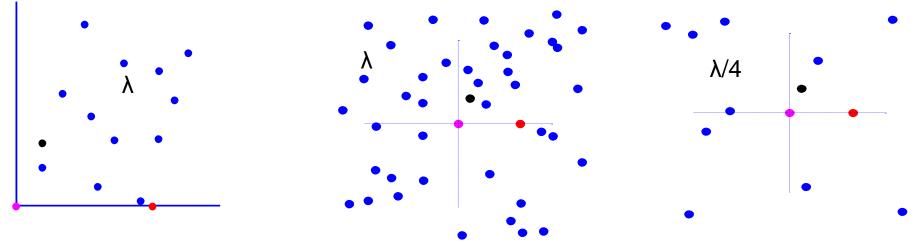
- Two deployment scenarios; Corner vs. Bulk
- Single eavesdropper fixed and known location
- When it becomes beneficial to hide the receiver at the corner?





## **Motivation**

- At the corner the mean and variance are scaled by <sup>1</sup>⁄<sub>4</sub>
- Impact of interferer's intensity in the bulk



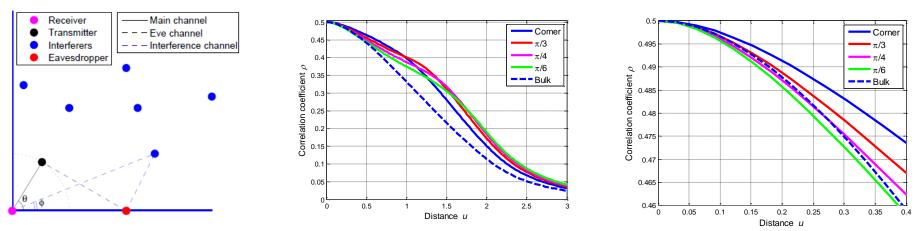
- How the performance at the corner with intensity of interferers  $\lambda$  differs from the performance in the bulk with scaled intensity ( $\lambda/4$ ) of interferers?
  - Interference at the receiver and the eavesdropper
  - Spatial correlation of interference



# Spatial correlation of interference

• Spatial correlation of interference is independent of the user density

$$\rho_{\mathbf{x}}(u) = \frac{\lambda \xi \int_{0}^{\infty} \int_{0}^{\phi_{\mathbf{x}}} g(r) g\left(\|re^{j\phi} - u\|\right) r \mathrm{d}\phi \,\mathrm{d}r}{\sqrt{\mathbb{V}\mathrm{ar}\left\{\mathcal{I}_{\mathbf{x}}(u)\right\}} \sqrt{\mathbb{V}\mathrm{ar}\left\{\mathcal{I}_{\mathbf{x}}\right\}}}$$



 Spatial correlation of interference is higher when the receiver is located at the corner – trade off



# Physical layer security

- Wyner encoding scheme
  - *R*<sub>t</sub> is the rate of the transmitted codewords
  - *R<sub>s</sub>* is the rate of confidential messages
  - $R_e = R_t R_s$  is the rate cost for securing the message against eavesdropping
- Unknown CSI
  - The rates R<sub>t</sub> and R<sub>s</sub> are kept fixed
  - SIR associated with the rates  $\mu = 2^{R_t} 1$   $\sigma = 2^{R_e} 1$
  - Probability of secure connectivity  $\mathbb{P}_{\mathbf{x}}^{\mathrm{sc}}(u) = \mathbb{P}(\gamma_{\mathbf{x},\mathbf{r}} > \mu, \gamma_{\mathbf{x},\mathbf{e}}(u) < \sigma)$
- Known CSI
  - AMC based on the instantaneous SIR
  - Average secrecy rate describes the performance

$$\overline{C}_{\mathbf{x}}^{\mathbf{sc}}(u) = \int_{0}^{\infty} \int_{0}^{\gamma_{\mathbf{x},\mathbf{r}}} \log_{2}\left(\frac{1+\gamma_{\mathbf{x},\mathbf{r}}}{1+\gamma_{\mathbf{x},\mathbf{e}}}\right) f_{\mathbf{r},\mathbf{e}}(\gamma_{\mathbf{x},\mathbf{r}},\gamma_{\mathbf{x},\mathbf{e}}) \,\mathrm{d}\gamma_{\mathbf{x},\mathbf{e}} \mathrm{d}\gamma_{\mathbf{x},\mathbf{r}}$$

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## Probability of secure connectivity – unknown CSI

• Probability of secure connectivity with correlated interference

 $\mathbb{P}_{\mathbf{x}}^{\mathrm{sc}}(u) = \mathbb{P}\left(\gamma_{\mathrm{x},\mathrm{r}} > \mu, \gamma_{\mathrm{x},\mathrm{e}}(u) < \sigma\right)$  $= \mathbb{E}\left\{e^{-s\mathcal{I}_{\mathbf{X},\mathbf{r}}}\left(1 - e^{-s_e\mathcal{I}_{\mathbf{X},\mathbf{e}}(u)}\right)\right\} = \mathbb{P}_{\mathbf{X},\mathbf{r}}^{\mathsf{c}} - \mathcal{J}_{\mathbf{X}}(u)$  $\mathcal{J}_{\mathbf{X}}(u) = \mathbb{E}\left\{e^{-s\mathcal{I}_{\mathbf{X},\mathbf{f}}-s_e\mathcal{I}_{\mathbf{X},\mathbf{e}}(u)}\right\}$  $= \int_{0}^{\frac{\pi}{2}} \exp\left(-\lambda \int_{\Theta} \left(1 - \frac{1}{1 + sg\left(r\right)} \frac{1}{1 + s_{e}g\left(d\right)}\right) \mathrm{d}S\right) f_{\Theta} \mathrm{d}\theta$  $= \int \exp\left(-\lambda \int \left(1 - \frac{1}{1 + \mu g(r) + \sigma z^{-1} g(d)}\right) dS\right) f_Z dz$ 





# Bulk vs. Corner at low transmission rates $R_t$

- Expanding the probability of secure connectivity for low  $\mu$ ,  $\sigma$
- Connection probability

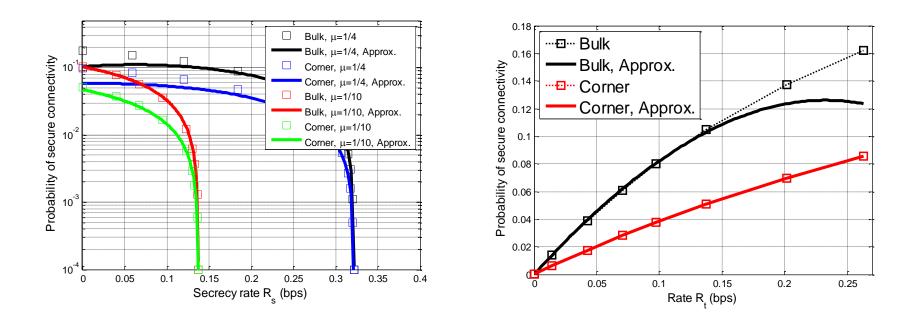
$$\begin{aligned} \mathbb{P}_{\mathbf{x},\mathbf{r}}^{\mathbf{c}} &= \exp\left(-\lambda \int_{0}^{\infty} \int_{0}^{\phi_{\mathbf{x}}} \frac{sg(r)}{1+sg(r)} r \mathrm{d}\phi \,\mathrm{d}r\right) \\ &\approx \exp\left(-\lambda \left(\int_{S_{\mathbf{x}}} \left(\mu g(r) - \mu^{2}g^{2}(r)\right) \,\mathrm{d}S\right)\right) \\ &= \exp\left(-\mu \mathbb{E}\{\mathcal{I}_{\mathbf{x},\mathbf{r}}\} + \frac{\mu^{2}}{2} \mathbb{V}\mathrm{ar}(\mathcal{I}_{\mathbf{x},\mathbf{r}})\right) \\ &\approx 1 - \mu \mathbb{E}\{\mathcal{I}_{\mathbf{x},\mathbf{r}}\} + \frac{\mu^{2}}{2} \left(\mathbb{V}\mathrm{ar}(\mathcal{I}_{\mathbf{x},\mathbf{r}}) + \mathbb{E}\{\mathcal{I}_{\mathbf{x},\mathbf{r}}\}^{2}\right) \end{aligned}$$

- Similarly we can expand  $\mathcal{J}_{\mathbf{x}}(u)$
- Probability of secure connectivity at low-rate transmissions is proportional to the mean interference at the eavesdropper



# Bulk vs. Corner at low transmission rates $R_t$

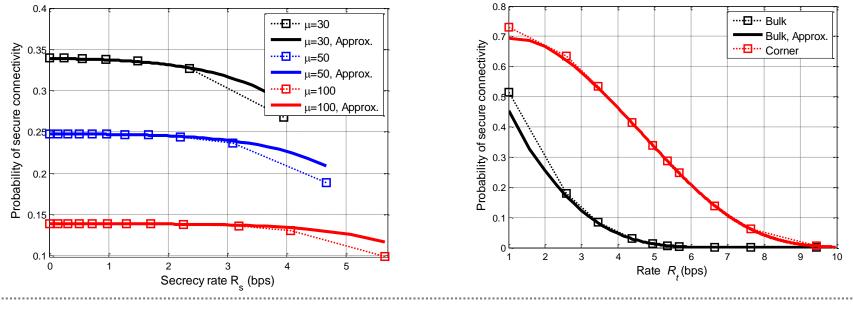
- 1. Mean interference along the boundary is always less than mean interference in the bulk  $\rightarrow$  Bulk is preferable for low rate transmissions
- 2. The impact of interference correlation vanishes at low rates





# Bulk vs. Corner at high transmission rates $R_t$

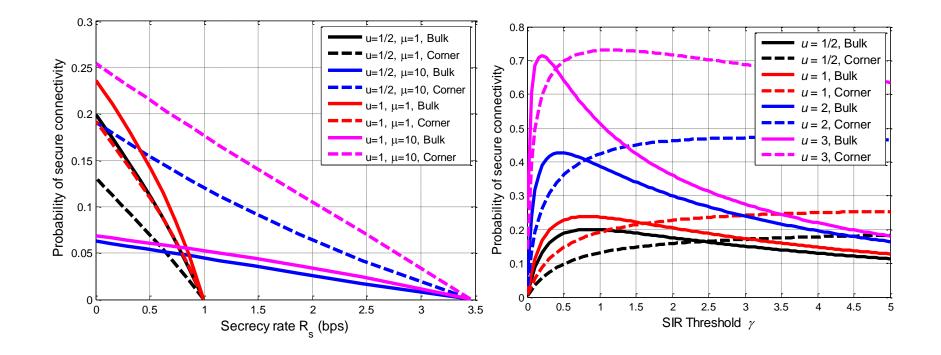
- 3. For large distance separation u between the receiver and the eavesdropper, placing the receiver at the corner is preferable at high transmission rates  $R_t$ 
  - Assume independent interference and expand for high μ, σ





# Bulk vs. Corner at high transmission rates $R_t$

4. For small distance separation *u* between the receiver and the eavesdropper, placing the receiver at the corner is still preferable



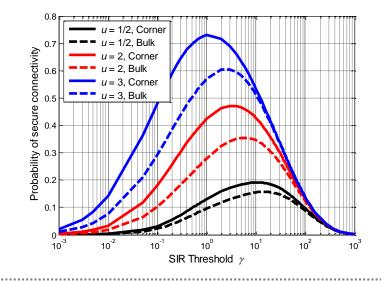


## Average secrecy rate – known CSI

• Physical layer security reduces the average rate by a quantity that depends on the joint connection probability of receiver and eavesdropper for  $\mu = \sigma = \gamma$ 

$$\overline{C}_{\mathbf{x}}^{\mathrm{sc}}(u) = \overline{C}_{\mathbf{x}} - \frac{1}{\log(2)} \int_{0}^{\infty} \frac{\mathcal{J}_{\mathbf{x}}(u,\gamma)}{1+\gamma} \mathrm{d}\gamma$$

5. The average secrecy rate at the corner is higher than in the bulk even if the density of interferers over there is 4 times higher than in the bulk







#### Future work

- Point Process for the eavesdroppers
- Interference correlation in more complex geometries
- Performance of secrecy enhancement techniques