

## Axiomatic Set Theory Exercises: Selected Solutions

**Exercise 1.1** In ST (= MATH32000 Set Theory) TC was given an alternative (but equivalent) definition, and was shown to satisfy the definition of TC above. Rework this by showing, using the above definition, that:  
 (i) If  $\text{Trans}(t) \wedge x \subseteq t \longrightarrow \text{TC}(x) \subseteq t$ . Hence  $\text{TC}(x)$  is the  $\subseteq$ -smallest transitive set  $t$  satisfying  $x \subseteq t$ . Moreover  $\text{Trans}(x) \iff \text{TC}(x) = x$ . (ii)  $x \in y \longrightarrow \text{TC}(x) \subseteq \text{TC}(y)$ . (iii) Define by recursion on  $\omega$ :

$$\bigcup^0 x = x; \bigcup^{n+1} x = \bigcup (\bigcup^n x); \text{tc}(x) = \bigcup \{\bigcup^n(x) \mid n < \omega\}. \text{ Show that } \text{tc}(x) = \text{TC}(x)$$

Solution: For (i) By  $\in$ -induction assume for all  $y \in x$  we have proven:

“ $\text{TC}(y)$  is the  $\subseteq$ -smallest transitive set  $t$  satisfying  $y \subseteq t$ ”.

We do the same for  $x$ . We first show  $\text{Trans}(\text{TC}(x))$ . Suppose  $z \in y \in \text{TC}(x) =_{\text{df}} x \cup \bigcup \{\text{TC}(w) \mid w \in x\}$ . If  $y \in x$  then  $\text{TC}(y) \subseteq \text{TC}(x)$  and  $z \in y \subseteq \text{TC}(y)$ . Hence  $z \in \text{TC}(x)$ . If  $y \notin x$  then for some  $w \in x$  we must have  $y \in \text{TC}(w)$ . By the Inductive Hypothesis,  $\text{Trans}(\text{TC}(w))$ . Hence  $z \in y \subseteq \text{TC}(w) \subseteq \text{TC}(x)$ . Hence  $\text{Trans}(\text{TC}(x))$ . Now let  $t \supseteq x$  be any other transitive set containing  $x$ . Let  $y \in x$ . As  $x \subseteq t \wedge \text{Trans}(t)$ , we have  $y \subseteq t$ . By the Ind. Hyp.  $\text{TC}(y)$  is the smallest transitive set containing  $y$ . Hence  $\text{TC}(y) \subseteq t$ . But then, as  $y$  was arbitrary,  $\text{TC}(x) = x \cup \bigcup \{\text{TC}(y) \mid y \in x\} \subseteq t$  as required.

That “ $\text{Trans}(x) \iff \text{TC}(x) = x$ ” is immediate from this.

(ii) is immediate:  $\text{TC}(y)$  by (the notes) definition is  $y \cup \bigcup \{\text{TC}(x) \mid x \in y\}$ . Hence  $x \in y \longrightarrow \text{TC}(x) \subseteq \text{TC}(y)$ .

For (iii):

First note as well that: (a)  $y \in x \longrightarrow \text{tc}(y) \subseteq \text{tc}(x)$  as it is easy to see that  $\bigcup^n y \subseteq \bigcup^{n+1} x$ ; and (b)  $z \in \bigcup^{n+1} x$  for some  $n$ , implies that  $z \in \bigcup^n y$  for some  $y \in x$ . (Both these statements could be proven rigorously by a (metatheoretic) induction on  $n$ , but they are reasonably obvious from the chain-like nature of the definition of  $\bigcup^{n+1} x$ :  $z \in \bigcup^{n+1} x \iff \exists y \exists z_0 \dots \exists z_n (z = z_n \in z_{n-1} \dots z_0 \in y \in x)$ .)

Now we show:  $\text{tc}(x) = \text{TC}(x)$ : Proof by  $\in$ -induction. Suppose  $\text{tc}(y) = \text{TC}(y)$  for all  $y \in x$ .

( $\supseteq$ ) Then  $\text{TC}(x) = x \cup \bigcup \{\text{tc}(y) \mid y \in x\} \subseteq \text{tc}(x)$  using (a).

( $\subseteq$ ) However if  $z \in \text{tc}(x)$ , then either  $z \in \bigcup^0 x = x$  in which case immediately  $z \in \text{TC}(x)$  or else  $z \in \bigcup^{n+1} x$  for some  $n$ , and hence  $z \in \bigcup^n y$  for some  $y \in x$  by (b). But then  $z \in \text{tc}(y) = \text{TC}(y)$  (by the Inductive Hypothesis). But  $\text{TC}(y) \subseteq \text{TC}(x)$  by definition of  $\text{TC}(x)$ . Hence  $z \in \text{TC}(x)$ .

**Exercise 1.2** (i) Show that the relation  $xRy \iff x \in \text{TC}(y)$  is wellfounded. (ii) Show for any  $x$  that  $\rho(x) = \rho(\text{TC}(x))$ ; (iii) Let  $X$  be any set. Show that  $\text{Trans}(X) \rightarrow \rho$  “ $X \in \text{On}$ ”.

Solution: (i) Let  $t$  be a non-empty term, and pick  $y \in t$ . If  $y$  is not  $R$ -minimal in  $t$ , let  $y_0$  be  $\rho$ -minimal in  $\text{TC}(y) \cap t$ . Then  $y_0$  is  $R$ -minimal in  $t$ . [Check:  $zRy_0 \longrightarrow z \in \text{TC}(y_0)$ . But the latter implies that  $\rho(z) < \rho(y_0)$ . (Check!) So  $z \notin t$ .]

(ii) As  $x \subseteq \text{TC}(x)$ ,  $\rho(x) \leq \rho(\text{TC}(x))$  is trivial. But note that if  $x \subseteq V_\alpha$  then  $\bigcup^k x \subseteq V_\alpha$  for any  $k \in \omega$  also. Hence  $\text{TC}(x) \subseteq V_\alpha$  too, and so  $\rho(x) \geq \rho(\text{TC}(x))$ .

(iii) Suppose there is  $X$  a transitive set, so that  $\rho$  “ $X \notin \text{On}$ ”; by  $\in$ -induction we may choose such an  $X$  with  $\rho(X)$  least. Then there is some  $\gamma < \rho(X)$  with no  $z \in X$  with  $\rho(z) = \gamma$ , whilst at the same time there are  $u \in X$  with  $\gamma < \rho(u)$ . By part (iii)  $\rho(u) = \rho(\text{TC}(u))$  and by hypothesis as  $\rho(u) < \rho(X)$ ,  $\rho$  “ $\text{TC}(u) \in \text{On}$ ”. So there is  $z \in \text{TC}(u)$  with  $\rho(z) = \gamma$ . But  $\text{TC}(u) \subseteq X$ , so  $z \in X$ . This is a contradiction.

**Exercise 1.4** Convince yourself of the truth of Lemma 1.23. [Hint: At least set out the base cases of the induction: suppose  $\varphi$  is  $v_0 \in v_1$  and let  $t_0 = x$ ,  $t_1 = \{z \mid \psi\}$ . Then  $(x \in t_1)^W \leftrightarrow (x \in \{z \mid \psi\})^W \leftrightarrow \psi(x/z)^W \leftrightarrow x \in \{z \mid z \in W \wedge \psi^W\} \leftrightarrow x^W \in (t_1)^W$ . The other base cases are relatively straightforward, but a little lengthy to write out. The inductive step for non-atomic formulae is easy by comparison.]

Solution: Assume that  $(x \in t_1)^W$  has been proven as shown. We do the other cases in the following order: (a)  $t_1 = t_2$ , (b)  $x = t_1$ , (c)  $t_1 \in t_2$ , then (d)  $t_1 \in x$ . Suppose in the following  $t_1 = \{z \mid \psi\}$ , as above, and  $t_2 = \{y \mid \varphi\}$ .

$$\begin{aligned} \text{(a)} \quad & (t_1 = t_2)^W \leftrightarrow (\{z \mid \psi\} = \{y \mid \varphi\})^W \\ & \leftrightarrow (\forall w (w \in \{z \mid \psi(z)\} \leftrightarrow w \in \{y \mid \varphi(y)\}))^W \\ & \leftrightarrow \forall w \in W ((w \in \{z \mid \psi(z)\} \leftrightarrow w \in \{y \mid \varphi(y)\}))^W \\ & \leftrightarrow \forall w (w \in \{z \mid \psi(z)\}^W \leftrightarrow w \in \{y \mid \varphi(y)\}^W) \end{aligned}$$

(using that  $\text{Trans}(W)$  in this last equivalence, and  $\{z|\psi(z)\}^W, \{y|\varphi(y)\}^W \subseteq W$ )

$$\leftrightarrow \{z|\psi(z)\}^W = \{y|\varphi(y)\}^W \leftrightarrow t_1^W = t_2^W.$$

(b)  $(x = \{z|\psi\})^W \leftrightarrow (\forall y(y \in x \leftrightarrow y \in \{z|\psi\}))^W$

$$\leftrightarrow \forall y \in W(y \in x \leftrightarrow y \in \{z|\psi\})^W$$

$$\leftrightarrow \forall y(y \in x \leftrightarrow (y \in \{z|\psi\})^W)$$

$$\leftrightarrow x^W = \{z|\psi\}^W.$$

(c)  $(\{z|\psi\} \in \{y|\varphi\})^W \leftrightarrow (\exists x(x = \{z|\psi\} \wedge x \in \{y|\varphi\}))^W$

$$\leftrightarrow (\exists x \in W(x = \{z|\psi\} \wedge x \in \{y|\varphi\}))^W$$

$$\leftrightarrow (\exists x \in W(x = \{z|\psi\} \wedge \varphi(x/y))^W)$$

$$\leftrightarrow \exists x(x \in W \wedge x = \{z|\psi\}^W \wedge \varphi(x/y)^W)$$

$$\leftrightarrow \{z|\psi\}^W \in \{x \in W | \varphi^W\} = \{y|\varphi\}^W.$$

(d)  $(\{z|\psi\} \in x)^W \leftrightarrow (\exists y(y = \{z|\psi\} \wedge y \in x))^W$

$$\leftrightarrow \exists y \in W((y = \{z|\psi\})^W \wedge (y \in x)^W) \quad (\text{but } x \subseteq W \text{ and so:})$$

$$\leftrightarrow ((\exists y(y = \{z|\psi\}))^W \wedge y \in x) \leftrightarrow \{z|\psi\}^W \in x^W.$$

Now the lemma is established for atomic formula with terms substituted, the inductive steps for the more complex formulae are straightforward. We only consider the case  $\varphi$  is  $\exists x\psi(x, y_1, \dots, y_n)$  with terms  $t_1, \dots, t_n$ . Then:

$$\varphi(t_1, \dots, t_n)^W \leftrightarrow \exists x \in W(\psi(x, t_1, \dots, t_n))^W \leftrightarrow \exists x \in W\psi(x, t_1^W, \dots, t_n^W)^W \leftrightarrow (\exists x\psi(x, t_1^W, \dots, t_n^W))^W$$

with the inductive hypothesis used at second equivalence.

**Exercise 1.5** Show (ii) and (iv) of Lemma 1.25: Let  $W$  be a transitive term.

(ii) If  $\omega \in W$  then  $(\text{Ax.Infinity})^W$ .

(iv) If for any  $x \in W$ , and term  $f$  with  $f^W$  is a function, then  $f^W \ulcorner x \in W$ , then  $(\text{Ax.Replacement})_W$ .

Solution: (ii)  $(\exists x(\emptyset \in x \wedge \forall y(y \in x \longrightarrow y \cup \{y\} \in x))^W \leftrightarrow$

$$\exists x \in W((\emptyset \in x \wedge \forall y(y \in x \longrightarrow y \cup \{y\} \in x))^W \leftrightarrow$$

$$\exists x \in W(((\emptyset \in x)^W \wedge \forall y \in W(y \in x \longrightarrow y \cup \{y\} \in x)^W).$$

But  $\emptyset^W = \emptyset$ , and  $(y \cup \{y\} \in x)^W \leftrightarrow y \cup \{y\} \in x$  since  $x \subseteq W$  as  $\text{Trans}(W)$ . Taking  $x$  as  $\omega$  which is in  $W$  we see the above last equivalence holds.

(iv) Suppose  $f$  given as hypothesized. We require  $(\text{Fun}(f) \longrightarrow f \ulcorner x \in V)^W$ . We are given that  $\text{Fun}(f)^W$  and that  $f^W \ulcorner x \in W = V^W$ . hence (with the help of Lemma 1.23 - meaning the Exercise ? 1.4 above on the atomic sentence of type  $t_1 \in t_2$ ) we have  $(f \ulcorner x \in V)^W$ .

**Exercise 1.7** Let  $W$  be class term  $\{\emptyset\}$ . Which axioms of ZFC hold in  $\langle W, \in \rangle$ ? Let  $U$  be the class term  $\text{On}$ . Which axioms of ZFC hold in  $\langle U, \in \rangle$ ? (NB For the latter,  $\langle \text{On}, \in \rangle$  just is  $\langle \text{On}, < \rangle$  of course.)

Solution:  $\langle W, \in \rangle$ : Note that  $\text{Trans}(W)$  (hence  $\text{Ax.Ext.}$  holds) and  $W$  has one member, namely  $\emptyset$  (hence  $\text{Ax.Empty}$  holds).  $\text{Ax.Pair}$ ,  $\text{Ax.Infinity}$ ,  $\text{Ax. Power}$  all fail. (For example, for  $\text{Ax.Pairing}$ :  $(\{\emptyset, \emptyset\} = \{\emptyset\}) \in V^W$  is false.) The others hold more or less vacuously:  $\text{Union}$ ,  $\text{Foundation}$ ,  $\text{Replacement}$ ,  $\text{Separation}$ . (For  $\text{Foundation}$ , note that the only non-empty term  $a$  is  $V^W = W = \{\emptyset\}$ . And  $\exists x(x \in a \wedge x \cap a = \emptyset)$ . For  $\text{Replacement}$ , note that there are no terms  $f$  with  $f^W$  a function - since  $W$  has no ordered pairs.) Note that  $\emptyset$  is a function: it is the empty function with empty domain! However technically,  $\prod \emptyset = \{\emptyset\} \neq \emptyset$ , (as the empty function  $\emptyset$  itself satisfies the clauses for  $h$  to be in  $\prod \emptyset$ !) so  $\text{AC}$  (in our version) is vacuously true (but very silly).

$\langle U, \in \rangle$ :  $\text{Trans}(\text{On})$  so  $\text{Ax.Ext.}$  holds, as easily does  $\text{Axs. Empty}$ ,  $\text{Infinity}$ , and  $\text{Foundation}$ . Similarly if  $x \in \text{On}$  then  $\bigcup x \in \text{On}$ . Hence  $\text{Ax.Union}$  holds.  $\text{Ax.Pair}$  fails, in particular, as the only unordered pair that exists is  $\{\emptyset, \emptyset\} = 1!$   $\text{Ax.Replacement}$  thus holds vacuously in  $U$  (as the only term defining a function is  $\emptyset$ ). For any  $\alpha \in \text{On}$ , we have  $(\mathcal{P}(\alpha) = \alpha + 1)^U$  (Check!) so  $\text{Ax.Power}$  holds. However  $\text{Ax.Separation}$ , fails. Does the  $\text{Ax.Choice}$  hold? This is  $\text{Fun}(f) \wedge \text{dom}(f) \in V \wedge \emptyset \notin \text{ran}(f) \longrightarrow \prod f \neq \emptyset$ . We've remarked that  $\text{Fun}(f) \longrightarrow f = \emptyset$  is the only function term for  $U$ . As above  $\prod \emptyset = \{\emptyset\} \neq \emptyset$  As  $\{\emptyset\} \in \text{On}$  it again looks like that  $\text{Ax9}$  holds in  $U!$

**Exercise 1.6** Let  $\varphi(v_0, \dots, v_n)$  be any formula. Let  $g_\varphi(\vec{y}) \approx$  the least  $\beta$  such that  $\exists x \varphi(x, \vec{y}) \longrightarrow \exists x \in V_\beta \varphi(x, \vec{y})$  if such an  $x$  and  $\beta$  exist; let it be 0 otherwise. Show that  $\forall \xi g_\varphi \text{ `` } V_\xi \in V$ . Deduce that  $f_\varphi(\xi) =_{\text{df}} \sup(g_\varphi \text{ `` } V_\xi)$  is a welldefined function.

Sol: This is just a use of the AxRep.:  $g_\varphi$  is a function and its action on  $V_\xi$ ,  $g_\varphi \text{ `` } V_\xi$ , yields a set for any  $\xi$ . Then the supremum of this set is an ordinal, and so  $f_\varphi$  is properly defined.

**Exercise 1.8** Which axioms of ZFC hold in  $V_\omega$ ?

Sol: All axioms other than Ax. Infinity. (See Lemmata 1.24-1.26.)

**Exercise 1.9** Check, or recheck, the following basic properties of the  $V_\alpha$  using the Definitions 1.17, 1.18 of  $\rho$  and  $V_\alpha$ :

- (i)  $\text{Trns}(V_\alpha)$  and if  $x \in V_\alpha$  then  $\forall y \in x (y \in V_\alpha \wedge \rho(y) < \rho(x))$ ;
- (ii)  $\alpha < \beta \longrightarrow V_\alpha \subseteq V_\beta$ ;
- (iii)  $V_{\alpha+1} = \mathcal{P}(V_\alpha)$ ;
- (iv) If  $x \in V$ , then  $\rho(x) =$  least  $\alpha$  so that  $x \subseteq V_\alpha =$  least  $\alpha$  such that  $x \in V_{\alpha+1}$ .
- (v)  $\rho(\alpha) = \alpha$ ;
- (vi)  $\text{On} \cap V_\alpha = \alpha$ .

Solution: (i),(ii) are immediate from the definitions 1.17, 1.18.

For (iii): If  $x \subseteq V_\alpha =_{\text{df}} \{z \mid \rho(z) < \alpha\}$  then  $\forall y \in x (\rho(y) < \alpha)$ . Thus  $\rho(x) =_{\text{df}} \sup \{\rho(y) + 1 \mid y \in x\} \leq \alpha$ . Hence  $x \in V_{\alpha+1}$ . Thus  $V_{\alpha+1} \supseteq \mathcal{P}(V_\alpha)$ ; But  $z \in V_{\alpha+1} \iff \rho(z) < \alpha + 1$ , which implies that any  $u \in z \in V_{\alpha+1}$  satisfies  $\rho(u) + 1 \leq \rho(z) < \alpha + 1$ . Hence  $\rho(u) < \alpha$  and  $u \in V_\alpha$ . Thus  $z \subseteq V_\alpha$ . Hence  $V_{\alpha+1} \subseteq \mathcal{P}(V_\alpha)$ .

For (iv): By  $\in$ -induction. Assume as Ind. Hyp. that for any  $y \in x$  the conclusions hold. Then by definition  $\rho(x) = \sup \{\rho(y) + 1 \mid y \in x\} = \sup \{\alpha + 1 \mid y \in x \wedge y \in V_{\alpha+1}\}$  which is the least  $\beta$  with  $x \subseteq V_\beta$ . That this  $\beta$  is the least  $\beta$  for which  $x \in V_{\beta+1}$  follows from (iii).

For (v): By induction on  $\alpha$ : Trivial for  $\alpha = 0$ . Suppose true for  $\beta < \alpha$ : then  $\rho(\alpha) =_{\text{df}} \sup \{\rho(\beta) + 1 \mid \beta \in \alpha\} = \sup \{\beta + 1 \mid \beta \in \alpha\} = \alpha$ .

(vi) is then immediate from (v) and Defs. 1.17, 1.18.

**Exercise 1.10** Sol: That this is a total order is easy, even trivial, to see. We check that it is a wellorder. Let  $X \neq \emptyset$  be a set of  $n$ -tuples of ordinals. Let  $Z \subseteq X$  be the subset of  $X$  with least max. Now just pick the lexicographically least  $\vec{\alpha}$  from  $Z$ . (That is, let  $\vec{\alpha}_0$  be least first element of some  $\vec{\alpha} \in Z$ . Then let  $\vec{\alpha}_1$  be least second element amongst those sequences starting out:  $\vec{\alpha}_0, \vec{\alpha}_1$  etc for  $n$ -steps.)

**Exercise 1.11\*E** Sol: first note that  $<^*$  is the lexicographic ordering when restricted to *descending* finite sequences of ordinals  $p: p_0 > p_1 > \dots > p_k, q: q_0 > q_1 > \dots > q_l$ . So, given two sets  $p, q$  as above note that  $p <^* q$  or  $q <^* p$  depending on which is lexicographically first when written out in this descending fashion. So the ordering is a strict total ordering. Let  $\emptyset \neq X$  be a set of finite sets of ordinals. Assume each  $p \in X$  is enumerated in descending order. Let

$$d_0 = \text{least } \{\delta \mid \delta = p_0 \text{ for some } p_0 \in p \in X\};$$

$d_1 = \text{least } \{\delta \mid \delta = p_1 \text{ for some } d_0 = p_0 > p_1 \in p \in X\}$ . Then  $d_1 < d_0$ . Proceeding in this way, after a finite number of steps (as the ordinals are wellordered) we arrive at some  $d_k < d_{k-1} \dots < d_1 < d_0$  the minimal element of  $X$ .

(What is the order type of  $\langle F(\omega + 1), <^* \rangle, \langle F(\omega^\omega + 1), <^* \rangle$  (where  $F(\delta)$  is the set of all finite subsets of  $\delta$ )?)

**Exercise 2.1** Prove (iii) of Lemma 2.7 : If  $f: \alpha \longrightarrow \beta$  is cofinal and strictly increasing, then  $\text{cf}(\alpha) = \text{cf}(\beta)$  and the corollary following: If  $\text{Lim}(\lambda)$  then  $\text{cf}(\omega_\lambda) = \text{cf}(\lambda)$ .

Sol: Suppose  $f, \alpha, \beta$  are as hypothesised. By definition of  $\text{cf}$  let  $g: \text{cf}(\alpha) \longrightarrow \alpha$  and  $h: \text{cf}(\beta) \longrightarrow \beta$  be (1-1) strictly increasing maps unbounded in  $\alpha, \beta$  respectively. (We can assume this by part (i) of Lem.2.7.) Then  $f \circ g$  is cofinal into  $\beta$  and shows  $\text{cf}(\beta) \leq \text{cf}(\alpha)$ . To show  $\text{cf}(\alpha) \leq \text{cf}(\beta)$  define  $k: \text{cf}(\beta) \longrightarrow \alpha$  by  $k(\tau) = \text{least } \xi$  with  $f(\xi) > h(\tau)$ . As  $f, h$  are strictly increasing unbounded maps,  $k$  is cofinal into  $\alpha$ .

Corollary: Let  $\text{Lim}(\lambda)$ : since  $F_\aleph \upharpoonright \lambda: \lambda \longrightarrow \omega_\lambda$  is cofinal and strictly increasing this is immediate by what we have just shown.

**Exercise 2.5** Let  $f: \Omega \longrightarrow \Omega$  be strictly increasing. Then  $f$  is normal iff  $\text{ran}(f)$  is c.u.b. in  $\Omega$ .

Sol: ( $\Rightarrow$ ) Let  $\xi < \Omega$ . As  $f$  is increasing  $f(\xi) \geq \xi$ . Hence  $\text{ran}(f)$  is trivially unbounded in  $\Omega$ . Let  $\zeta$  be a limit ordinal with  $\text{ran}(f)$  unbounded in  $\zeta$ . We require that  $\zeta \in \text{ran}(f)$ . Using the continuity of  $f$  we have that  $f(\xi) = \zeta$ , where  $\xi = \sup \{\eta \mid f(\eta) < \zeta\}$ .

( $\Leftarrow$ ) Let  $\text{ran}(f)$  be cub in  $\Omega$ . As we are told  $f$  is increasing we only have to check (ii) in the definition of normality. Let  $\text{Lim}(\lambda)$ . Let  $\tau = \sup f \ll \lambda$ . But  $\text{ran}(f)$  is closed. Thus  $\tau \in \text{ran}(f)$ . Thus  $\tau$  must be  $f(\lambda)$ , just as required by (ii) of normality.

**Exercise 2.6** For any  $E \subseteq \text{On}$  define  $E^*$  to be the class of limit points of  $E$ : namely those limit ordinals  $\beta$  such that  $E \cap \beta$  is unbounded in  $\beta$ . Show that  $E^*$  is a closed class, and if  $E \in V$ , with  $\text{cf}(\sup(E)) > \omega$ , then  $E^*$  is c.u.b. below  $\sup(E)$ .

Sol: That  $E^*$  is closed is immediate: if  $E^* \cap \beta$  is unbounded in  $\beta$ , then for any  $\delta < \beta$  there is  $\tau \in E^* \cap \beta$  with  $\tau > \delta$ . But as  $\tau \in E^*$  then there is a  $\tau' \in E \cap \tau$  with  $\tau > \tau' > \delta$  too. Hence  $E \cap \beta$  is also unbounded in  $\beta$  and thus  $\beta \in E^*$ . Suppose  $E$  is a set. Let  $\eta_0 < \sup(E)$  be arbitrary. Any  $\omega$ -sequence  $\eta_0 < \eta_1 < \dots < \eta_n < \dots$  of elements from  $E$  has supremum in  $E^*$  which is below  $\sup(E)$  as the latter has uncountable cofinality. Hence  $E^*$  is unbounded in  $\sup(E)$ . Closure was the first part.

**Exercise 2.7** suppose  $\Omega$  is a regular cardinal. (i) Let  $C, D \subseteq \Omega$  be c.u.b.in  $\Omega$ . Show that  $C \cap D$  is c.u.b.in  $\Omega$ . (ii) Now generalise this argument: let  $\gamma < \Omega$ . Let  $\langle C_\xi \mid \xi < \gamma \rangle$  be a sequence of c.u.b.in  $\Omega$  classes. Show that  $\bigcap_{\xi < \gamma} C_\xi$  is c.u.b. in  $\Omega$ .

Sol: (i) Let  $\alpha < \Omega$  be arbitrary. Pick alternately:  $\alpha \leq \alpha_0 < \beta_0 < \alpha_1 < \dots < \alpha_n < \beta_n < \dots$  for  $n < \omega$  with  $\alpha_n \in C$ ,  $\beta_n \in D$ . Let  $\gamma = \sup \{\alpha_n\} = \sup \{\beta_n\}$ . By closure of  $C$  and  $D$   $\gamma \in C \cap D$  provided that  $\gamma < \Omega$ . However this is true: if  $\Omega$  is a regular cardinal it is because  $\text{cf}(\Omega) = \Omega$  (and not  $\omega$ !). Hence  $C \cap D$  is unbounded in  $\Omega$ . But it is also closed: if  $\text{Lim}(\mu)$  and  $(C \cap D) \cap \mu$  is unbounded in  $\mu$ , then both  $C \cap \mu$  and  $D \cap \mu$  are unbounded in  $\mu$  and thus  $\mu \in C \cap D$ . Hence  $C \cap D$  is closed.

(ii): Let  $\alpha < \Omega$  be arbitrary. Using repeatedly the fact that  $\Omega$  is regular and  $\gamma < \Omega$ , we can successfully pick alternately:  $\alpha \leq \alpha_0^0 < \alpha_1^0 < \alpha_2^0 < \dots < \alpha_i^0 < \dots$  for  $i < \gamma$  with  $\alpha_i^0 \in C_i$ . Let  $\alpha^1 = \sup_i \alpha_i^0$ . Then  $\alpha^1 < \Omega$  again, as the latter is regular and  $\gamma < \Omega$ . Now repeat this again with  $\alpha^1 \leq \alpha_0^1 < \alpha_1^1 < \alpha_2^1 < \dots < \alpha_i^1 < \dots$  for  $i < \gamma$ , and define  $\alpha^2 = \sup \{\alpha_i^1 \mid i < \gamma\}$ . Then  $\alpha^2 < \Omega$ . Etc for  $n < \omega$ . Let  $\alpha^\omega = \sup_n \alpha^n < \Omega$  (once more by regularity of  $\Omega$ ). As  $\alpha^\omega = \sup \{\alpha_i^n \mid n < \omega\}$  for any  $i < \gamma$ , then by closure of  $C_i$  we have that  $\alpha^\omega \in C_i$ . Hence  $\alpha^\omega \in \bigcap_{\xi < \gamma} C_\xi$ . This shows the latter intersection is unbounded in  $\Omega$ . Closure is immediate: if  $\text{Lim}(\mu)$  then  $((\bigcap_{\xi < \gamma} C_\xi) \cap \mu)$  is unbounded in  $\mu$  implies  $\forall \xi < \gamma$  ( $C_\xi \cap \mu$  is unbounded in  $\mu$ ) which implies  $\forall \xi < \gamma$  ( $\mu \in C_\xi$ ).

**Exercise 2.8 (Diagonal Intersections)** Let  $\Omega \in \text{Reg}$ . Let  $\langle E_\xi \mid \xi < \Omega \rangle$  be a sequence of subsets of  $\Omega$ . Define the diagonal intersection of the sequence to be the set  $D =_{\text{df}} \Delta_{\xi < \Omega} \langle E_\xi \mid \xi < \Omega \rangle = \{\alpha < \Omega \mid \forall \beta < \alpha (\alpha \in E_\beta)\}$ . Now suppose that the  $E_\xi$  are all c.u.b.in  $\Omega$ . Show that (i) the diagonal intersection  $D$  is c.u.b.in  $\Omega$  (ii)  $D = \bigcap_{\alpha < \Omega} (E_\alpha \cup (\alpha + 1))$ .

Sol: (i) Closure is immediate again as in the last Exercise. ( $x \subseteq D \longrightarrow \sup x \in E_\xi \cup \{\Omega\}$  for any  $\xi < \sup x$ .) Unboundedness: let  $\alpha^0 < \Omega$  be arbitrary; by the unboundedness of  $\bigcap_{\xi < \gamma} E_\xi$  for any  $\gamma < \Omega$  (by the last Exercise) we may for any  $\gamma < \Omega$ , define  $f(\gamma)$  to be the least element of  $\bigcap_{\xi < \gamma} E_\xi$  above  $\gamma$ . Now define  $\alpha^0 < \alpha^1 = f(\alpha^0) < \dots < \alpha^{n+1} = f(\alpha^n)$ , and  $\alpha^\omega = \sup_n \alpha^n$ ; then check  $\alpha^\omega \in D$ . For (ii): Let  $\beta \in D$ . Then  $\forall \gamma < \beta (\beta \in E_\gamma) \wedge \forall \gamma \geq \beta (\beta \in \gamma + 1)$ . Conversely  $\beta \in \bigcap_{\alpha < \Omega} (E_\alpha \cup (\alpha + 1))$  implies *a fortiori* that  $\forall \gamma < \beta (\beta \in E_\gamma)$  and thus  $\beta \in D$ .

**Exercise 2.9** Show that  $\forall \alpha (|V_{\omega+\alpha}| = \beth_\alpha)$ .

Sol: By induction on  $\alpha$ :  $|V_\omega| = \omega = \beth_0$ ; if  $|V_{\omega+\alpha}| = \beth_\alpha$  is proven, then  $|V_{\omega+\alpha+1}| = |\mathcal{P}(V_{\omega+\alpha})| = 2^{|V_{\omega+\alpha}|} = \beth_{\alpha+1}$ ; if  $|V_{\omega+\alpha}| = \beth_\alpha$  for  $\alpha < \lambda$ , then  $|V_{\omega+\lambda}| = |\bigcup_{\alpha < \lambda} V_{\omega+\alpha}| = \bigcup_{\alpha < \lambda} \beth_\alpha = \beth_\lambda$ . (Remark: once  $\alpha \geq \omega^2$  then we have  $\omega + \alpha = \alpha$  in any case, so the annoying “ $\omega +$ ” in the subscript can be dropped.)

**Exercise 2.10** (i) Show that the GCH (Generalised Continuum Hypothesis, that  $\forall \alpha (2^{\aleph_\alpha} = \aleph_{\alpha+1})$ ) implies that  $\forall \alpha (\aleph_\alpha = \beth_\alpha)$ . (ii) Show that the first fixed point of the  $\beth$  function has cofinality  $\omega$ . (iii) Show that for any regular cardinal  $\kappa$  there is  $\alpha$ , a fixed point of the  $\beth$  function, with  $\text{cf}(\alpha) = \kappa$ . [Hint: this is simple, just consider an enumeration of the fixed points.]

Sol: (i): This is almost by definition, GCH implies that  $\forall \alpha (2^{\aleph_\alpha} = \aleph_{\alpha+1})$ , if, by induction we have proven  $\aleph_\alpha = \aleph_\alpha$ , then  $\aleph_{\alpha+1} = 2^{\aleph_\alpha} = 2^{\aleph_\alpha} = \aleph_{\alpha+1}$ . Then  $\aleph_\lambda = \aleph_\lambda$  as both are defined continuously at limit  $\lambda$ .

(ii) Let  $f(0) = \aleph_0; f(n+1) = \aleph_{f(n)}$ ; let  $f(\omega) = \sup \{f(n) \mid n < \omega\}$ . Then  $\aleph_{f(\omega)} = f(\omega)$  (For, if  $\gamma < \aleph_{f(\omega)}$  then  $\gamma < \aleph_{f(n)} = f(n+1)$  for some  $n$ , hence  $\aleph_\gamma < f(\omega)$ , hence  $\aleph_{f(\omega)} \leq f(\omega)$ ). As the  $\aleph$ -function is increasing,  $\gamma \leq \aleph_\gamma$  for all  $\gamma$ .)  $f(\omega)$  obviously has cofinality  $\omega$ . By construction we could not have any  $\kappa < \aleph_{f(\omega)}$  a smaller fixed point.

(iii) Just continue the construction of (ii) for  $\kappa$  many steps, then  $\aleph_{f(\kappa)} = f(\kappa)$ , and  $\text{cf}(\aleph_{f(\kappa)}) = \kappa$ .

**Exercise 2.12** Let  $S \subseteq \Omega$  be stationary and  $C \subseteq \Omega$  be c.u.b. Then  $S \cap C$  is stationary.

Sol: Let  $D \subseteq \Omega$  be an arbitrary c.u.b. class. Then  $D \cap C$  is c.u.b. Hence  $(S \cap C) \cap D \neq \emptyset$  as required.

**Exercise 2.13** Let  $S \subseteq \Omega$  be stationary. Show that  $S \cap S^*$  is stationary.

Sol: Immediate from the last Exercise, since  $S^*$  is c.u.b.

**Exercise 2.14** Can you generalise Example 1 to larger cardinals?

Sol: For any  $n \geq 2$  we could have  $S_{\omega_k} = \{\alpha < \omega_n \mid \text{cf}(\alpha) = \omega_k\}$  and these would be disjoint stationary sets for  $k < n$ . More generally for  $\kappa > \lambda \geq \omega_0, \kappa, \lambda \in \text{Reg}$ , define  $S_\lambda = \{\alpha < \kappa \mid \text{cf}(\alpha) = \lambda\}$  and then for different such  $\lambda, \lambda'$   $S_\lambda, S_{\lambda'}$  are disjoint and stationary.

**Exercise 2.15** Find  $S_n \subseteq \aleph_{\omega+1}$  stationary, for  $n < \omega$ , with  $S_{n+1} \subseteq S_n$ , but with  $\bigcap_n S_n = \emptyset$ .

Sol: Let  $S_n = \{\alpha < \aleph_{\omega+1} \mid \text{cf}(\alpha) \geq \aleph_n\}$ . Then each  $S_n$  is stationary, but for every  $\alpha < \aleph_{\omega+1}$ , for just one  $n$  is  $\text{cf}(\alpha) = \aleph_n$ . Hence the intersection of the  $S_n$  must be empty.

**Exercise 2.26** Let  $\langle H, R \rangle \in \text{WO}$ . Apply the Collapsing Lemma. What is the outcome?

Sol:  $\pi: \langle H, R \rangle \cong \langle \alpha, \in \rangle$  where  $\alpha$  is the order type of  $\langle H, R \rangle$ .

**Exercise 2.27** Show that  $V_\omega$  can be 'coded' as a subset of  $\omega$ : that is there is  $E \subseteq \omega$  so that  $\langle \omega, E \rangle \cong \langle V_\omega, \in \rangle$ . [Hint: Define  $nEm \leftrightarrow_{\text{df}}$  the "2<sup>n</sup>" column in the binary expansion of  $m$  contains a 1; (thus  $\{n \mid nE11\} = \{0, 1, 3\}$ ); check there is  $u$  satisfying  $\langle \omega, E \rangle \cong \langle u, \in \rangle$  with  $\text{Trans}(u)$ . Show  $u = V_\omega$ .]

Sol: We use the given  $E$  as our coding method. We check that  $E$  is (i) wellfounded (this is because  $\langle \mathbb{N}, < \rangle$  is wellfounded) (ii)  $E$  is extensional (this is because two different  $n, n'$  have different binary expansions!). Hence by the Mostowski-Shepherdson theorem there is a  $\langle u, \in \rangle$  with  $\pi: \langle \omega, E \rangle \cong \langle u, \in \rangle$ . ( $u \subseteq V_\omega$ ): suppose  $v \in u$  and all members of  $v$  are in  $V_\omega$ . As  $v$  has only finitely many members  $v \subseteq V_n$  for some  $n$ . Hence  $v \in V_{n+1} \subseteq V_\omega$ . The result follows by  $\in$ -induction.

( $u \supseteq V_\omega$ ): conversely by induction on  $n$  suppose  $V_n \subseteq u$ . If  $X \subseteq V_n$  is an arbitrary subset,  $X$  is of course finite, and its members are coded by  $\pi^{-1} \text{``} X \subseteq \omega$ . If  $\pi^{-1} \text{``} X$  is  $\{n_1, n_2, \dots, n_k\}$  it is a simple matter to find an  $n$  so that  $\{m \mid mEn\} = \{n_1, n_2, \dots, n_k\}$ . Then  $\pi(n) = X \in u$ . Hence  $\mathcal{P}(V_n) = V_{n+1} \subseteq u$ .

**Exercise 2.28** Show that if  $A, B$  are transitive sets, and  $f: \langle A, \in \rangle \rightarrow \langle B, \in \rangle$  is an isomorphism, that  $f = \text{id} \upharpoonright A$ .

Sol: By  $\in$ -induction: suppose for any  $x \in A$ , that we have shown for all  $y \in x$  that  $f(y) = y$ . (Note that as  $\text{Trans}(A), x \subseteq A$ .) As  $f$  is an isomorphism for any  $y \in A: y \in x \leftrightarrow f(y) \in f(x)$ ; Thus  $f(x) = \{f(y) \mid y \in x\}$ . But then  $f(x) = \{f(y) \mid y \in x\} = \{y \mid y \in x\} = x$ . As this is true for any  $x \in A$  we have the desired conclusion by " $\in$ -induction", that is we have used the following principle:

$$\forall x [\forall y \in x (\Phi(y) \rightarrow \Phi(x))] \rightarrow \forall x \Phi(x).$$

**Exercise 2.30** Find an example of an  $\langle x, \in \rangle$  which is not extensional. If we nevertheless apply the Mostowski-Shepherdson Collapse function  $\pi$  to it, what happens?

Sol: e.g. take  $x = \{0, \{1\}\}$ . Then  $\pi(0) = 0 = \pi(\{1\})$ . That is  $\pi: \langle x, \in \rangle \rightarrow \langle \{0\}, \in \rangle$ , and the map is not (1-1).

**Exercise 2.31** Prove (i),(iii)-(iv) here.

(i)  $\text{On} \cap H_\kappa = \kappa; \text{Trans}(H_\kappa)$ ;

(iii)  $y \in H_\kappa \wedge x \subseteq y \rightarrow x \in H_\kappa$ ;

- (iv)  $x, y \in H_\kappa \longrightarrow \bigcup x, \{x, y\} \in H_\kappa$ ;  
(v) Give an example to show that  $x \in H_\kappa \leftrightarrow x \subseteq H_\kappa \wedge |x| < \kappa$  fails if  $\kappa$  is a singular cardinal.

Sol: (i): All ordinals are already transitive; hence  $\alpha \in \text{On} \cap H_\kappa \leftrightarrow |\alpha| < \kappa \leftrightarrow \alpha < \kappa$ . Let  $y \in x \in H_\kappa$ , then  $\text{TC}(y) \subseteq \text{TC}(x) \longrightarrow |\text{TC}(y)| \leq |\text{TC}(x)| < \kappa$ . Hence  $y \in H_\kappa$  and so  $\text{Trans}(H_\kappa)$ .

(iii):  $|\text{TC}(y)| < \kappa \wedge x \subseteq y \longrightarrow \text{TC}(x) \subseteq \text{TC}(y) \wedge |\text{TC}(x)| < \kappa$ .

(iv) :  $|\text{TC}(\bigcup x)| \leq |\text{TC}(x)|$  (since  $\text{TC}(\bigcup x) \subseteq \text{TC}(x)$ ); hence  $x \in H_\kappa \longrightarrow \bigcup x \in H_\kappa$  is immediate. Similarly  $|\text{TC}(\{x, y\})| \leq \max\{|\text{TC}(x)|, |\text{TC}(y)|\}$ , from which  $x, y \in H_\kappa \longrightarrow \{x, y\} \in H_\kappa$  is also immediate.

(v) fails if  $\kappa = \aleph_\omega$  for example:  $A = \{\aleph_n | n < \omega\} \subseteq H_{\aleph_\omega} \wedge |A| = \aleph_0 < \aleph_\omega$  but  $A \notin H_{\aleph_\omega}$ .

**Exercise 2.32** Show that  $V_\omega = \text{HF}$ . [Hint: For  $(\subseteq)$  use induction on  $n$  to show  $V_n \in \text{HF}$ . For  $(\supseteq)$  use  $\in$ -induction].

Sol: Follow the Hint:  $(\subseteq)$  suppose as inductive hypothesis that  $V_n \in \text{HF}$ . Then  $V_n$  is finite, as is  $\mathcal{P}(V_n) = V_{n+1}$ . But  $V_{n+1} \subseteq \text{HF}$ , as any  $X \in V_{n+1}$  is a subset of  $V_n \subseteq \text{HF}$  and is finite, and so is in HF. (see Lemma 2.31(v)). For  $(\supseteq)$  suppose  $x \in \text{HF}$  and as an  $\in$ -inductive hypothesis that  $\forall y \in x (y \in V_\omega)$ . As  $x$  is finite,  $m = \max\{\rho(y) | y \in x\} < \omega$ . So  $x \subseteq V_m$ . Hence  $x \in V_{m+1} \subseteq V_\omega$ .

**Exercise 2.33** Check that HF is closed under all the assumptions of Lemmata 1.24, 1.25, (except 1.25(ii)) and even the power set operation. Hence  $(\text{ZF} - \text{Ax.Inf})^{\text{HF}}$ .

Sol: This is very much routine checking given the definition of HF. As  $\text{Trans}(\text{HF})$  we have  $\text{Ax.Extensionality}$ . For the basic operations of pairing and union (and Power set) we can appeal to the Ex. 2.31(iv), and the fact that we know that  $V_\omega$  is closed under such operations; likewise  $\text{Ax.Empty}$  is trivial. Since HF is closed under the power set operation  $\text{Ax.Subsets}$  becomes trivial.  $\text{Ax.Replacement}$  holds because the range of a function on a finite set is finite. (More detail: if  $u \in \text{HF}$  and has size  $k$  and  $F$  is a function  $F: \text{HF} \longrightarrow \text{HF}$  then  $F''u \subseteq \text{HF}$  and has size at most  $k$  also. Hence  $F''u \in \text{HF}$ .) The negation of the  $\text{Ax.Infinity}$  is true, since no set  $x \in \text{HF}$  is closed under the successor function  $S(u) =_{\text{df}} u \cup \{u\}$ .  $\text{Ax.Found.}$  holds in any transitive structure  $\langle M, \in \rangle$ .

**Exercise 2.35.** If  $x \in \text{HC}$  then we have  $|\text{TC}(x)| \leq \omega$ . Define a wellfounded extensional relation  $E$  on  $\omega$  so that  $\langle \omega, E \rangle \cong \langle \text{TC}(x), \in \rangle$ . [Hint: We have a bijection  $f: N \longleftrightarrow \text{TC}(x)$  for some  $N \leq \omega$ ; define  $nEm \leftrightarrow f(n) \in f(m)$ .] If we use a recursive pairing bijection  $p: \omega \longleftrightarrow \omega \times \omega$  (for example  $p^{-1}(\langle k, l \rangle) = 2^k \cdot (2l + 1) - 1$ ) we may further code  $E$  as a subset  $\bar{E} \subseteq \omega$ . We thus have effectively coded up  $\text{TC}(x)$  as a subset of  $\omega$ . [By using further such coding devices we may take any countable structure with domain in HC and code it up as a subset of  $\omega$ .]

Sol: The more interesting case here is with  $\text{TC}(x)$  infinite. The Hint here really says it all: as opposed to the coding we did earlier of  $V_\omega$  onto  $\langle \omega, E \rangle$  in Ex.2.27, we don't really need to fix on a particular fancy coding device: any bijection  $f: N \longleftrightarrow \text{TC}(x)$  will do, as will any pairing function  $p: \omega \longleftrightarrow \omega \times \omega$ . The point is just to code up the  $\in$ -information up by natural numbers. There are many  $\bar{E}$  with  $\langle \omega, \bar{E} \rangle \cong \langle \text{TC}(x), \in \rangle$  (in fact  $2^{\aleph_0}$  many).

**Exercise 2.36** Which axioms of ZF hold in  $V_\alpha$  if  $\text{Lim}(\alpha)$ ? Find a wellordering  $\langle A, R \rangle \in V_{\omega+\omega}$  but for which there is no ordinal  $\beta$  with  $\langle A, R \rangle \cong \langle \beta, < \rangle$ ; hence find an instance of the  $\text{Ax.Replacement}$  that fails in  $V_{\omega+\omega}$ . [The latter is a model of Z, the system Zermelo which is more or less ZF with Replacement removed. For almost all regions of mathematical discourse,  $V_{\omega+\omega}$  is a sufficiently large "universe" - mathematicians rarely need sets outside of this set.]

Sol: If  $\text{Lim}(\alpha)$  then all axioms of ZF relativise to  $V_\alpha$  except the  $\text{Ax.Replacement}$ . (One easily sees that such a  $V_\alpha$  is closed under pairing, union, subsets etc, by calculating ranks of sets as we did in ST). (1) Let  $\langle A, R \rangle$  be  $\langle \mathbb{N}, < \rangle$  where  $< \prime$  is the ordering that puts all the Evens before all the Odds, but otherwise retains the usual order. Then  $\text{ot}(\langle A, R \rangle) = \omega + \omega \notin V_{\omega+\omega}$ . Then  $G(2k + 1) = \text{ot}(\langle \mathbb{N}, < \prime \upharpoonright 2k + 1 \rangle)$  is a definable function that maps Odds unboundedly into  $\omega + \omega$ . So  $\text{ran}(G)$  is not a set in  $V_{\omega+\omega}$  although Odds is. Thus  $\text{AxRep.}$  fails.

(2) Alternative viewpoint: Consider  $V_{\omega+\omega}$  and the class term function  $F(\delta) = A_\omega(\delta) = \omega + \delta$ . However  $F''\omega = \{\omega, \omega + 1, \dots, \omega + k, \dots\} \notin V_{\omega+\omega}$ . ( $F''\omega \in V_{\omega+\omega} \longrightarrow \sup F''\omega = \bigcup F''\omega = \omega + \omega \in V_{\omega+\omega}$ . This is absurd as  $\text{On} \cap V_{\omega+\omega} = \omega + \omega$ .) Thus again  $\text{Ax.Replacement}$  fails.

**Exercise 2.38** Let  $\kappa \in \text{Card}$ . Show that  $|H_\kappa| = 2^{<\kappa}$ .

Sol: The Hint was to use the last Exercise: that concluded that  $|H_{\kappa^+}| = |\mathcal{P}(\kappa)|$ . This equals  $2^\kappa = 2^{<\kappa^+}$ . So this is proven for successor cardinals  $\kappa^+$ . Let now  $\kappa$  be a limit cardinal, and as an inductive hypothesis that  $|H_\mu| = 2^{<\mu}$  for  $\mu \in \text{Card} \wedge \mu < \kappa$ . Then  $H_\kappa = \bigcup \{H_\mu \mid \mu \in \text{Card} \cap \kappa\}$ . As the cardinality of  $H_\mu$  is increasing with  $\mu$ , we have:  $|H_\kappa| = \sup \{|H_\mu| \mid \mu \in \text{Card} \cap \kappa\} = \sup \{2^{<\mu} \mid \mu \in \text{Card} \cap \kappa\} = 2^{<\kappa}$ .

**Exercise 2.39** (Levy) Let  $h(\kappa)$  be the class of sets  $x$  with (i)  $|x| < \kappa$  (ii)  $\forall y \in \text{TC}(x) (|y| < \kappa)$ . Show that if  $\kappa \in \text{Reg}$ , then  $H_\kappa = h(\kappa)$ ; find an example where this fails if  $\kappa$  is singular.

Sol: Suppose  $\kappa$  is  $\aleph_\omega$ . Then the set  $x = \{\aleph_n \mid n < \omega\}$  is a counterexample: it is in  $h(\kappa)$  but not  $H_\kappa$ . Suppose now  $\kappa \in \text{Reg}$ . Let  $x \in H_\kappa$ . Then  $|\text{TC}(x)| < \kappa$ . This implies both that  $|x|$  and  $|\text{TC}(y)| < \kappa$  for any  $y \in x$ , and so  $x \in h(\kappa)$ . Conversely let  $x \in h(\kappa)$ . Then  $\text{TC}(x)$  is the union of less than  $\kappa$  many sets of cardinality  $< \kappa$ . Hence, as  $\kappa \in \text{Reg}$  it has cardinality  $< \kappa$  and thus is in  $H_\kappa$ .

**Exercise 2.41** Let  $W$  be a transitive class term. Then (i) any  $\Sigma_1$ -formula  $\varphi$  is upwards absolute for  $W$ ; (ii) any  $\Pi_1$ -formula  $\varphi$  is downwards absolute for  $W$ .

Sol: (i) Suppose  $\varphi$  is the formula  $\exists x \psi(x, y)$  with the single free variable  $y$  displayed (for simplicity) and with  $\psi(x, y) \Delta_0$ . Let  $W$  be a transitive class. Then for any  $x, y \in W$  we have  $\psi(x, y) \leftrightarrow \psi(x, y)^W$  since  $\Delta_0$  formulae are absolute by the last Exercise). By the definition of upward absoluteness we must show:  $\forall y \in W (\varphi^W \rightarrow \varphi)$ . Let  $y \in W$  be arbitrary and assume  $(\exists x \psi(x, y))^W$ . Let  $x_0 \in W$  with  $(\psi(x_0, y))^W$ . By the comment above  $\psi(x_0, y)$  holds (in  $V$ ). Hence so does  $\exists x \psi(x, y)$  i.e.  $\varphi$ .

(ii) This follows from (i) where now  $\varphi$  is  $\Pi_1$  that is (for simplicity) of the form  $\forall x \psi(x, y)$  with  $\psi$  in  $\Delta_0$ . But then if  $y \in W$  we have that  $\forall x \psi(x, y) \rightarrow \forall x \in W \psi(x, y) \leftrightarrow \forall x \in W (\psi(x, y))^W \leftrightarrow \varphi^W$ .

**Exercise 2.42** Verify the very last part of Lemma 2.40.

Sol: [This just asks that you understand that the Lemma follows from the *Claim*.] The  $\beta$  just constructed above is a closure point of the functions  $F_i$ . By appealing to Lemma 2.40, we then see that any formula on the list of the  $\vec{\varphi}$  is absolute between  $Z_\beta$  and  $Z$ . There are arbitrarily large such  $\beta$  by the construction, which we performed within the ZF axiom system. So the Lemma 2.40 follows.

**Exercise 2.43** Show that for every formula  $\varphi$  of  $\mathcal{L}$ :

ZF  $\vdash$  "There is a c.u.b. class  $C \subseteq \text{On}$  so that  $\forall \alpha \in C \forall \vec{x} \in V_\alpha (\varphi(\vec{x}) \leftrightarrow (\varphi(\vec{x}))^{V_\alpha})$ "

[Hint: The reasoning of Lemma 2.40 pretty much gives the relevant cub class as the closure points of the  $F_i$ .] Remark: One might think that one could enumerate all the axioms of ZF  $\varphi_0, \varphi_1, \dots$ , find the appropriate classes  $C_{\varphi_n}$  and take  $D = \bigcap_n C_{\varphi_n}$ . This appears then only to be an intersection of only countable many c.u.b. classes and so must be c.u.b. in  $\text{On}$ ? But for any element  $\alpha \in D$  we'd have  $(\text{ZF})^{V_\alpha}$ , and we appear to have proven the existence of models of ZF - contradicting Goedel. What is wrong with this reasoning?

Sol: This is really only an observation on the proof of the Montague-Levy Reflection Theorem. Let  $\vec{\varphi} = \varphi_0, \dots, \varphi_n$  be a subformula closed list with  $\varphi$  equal to  $\varphi_n$  (without loss of generality). Just note that the function  $F_n$  defined in that proof is not just monotone but *normal*, hence  $\text{ran}(F)$  is just the desired c.u.b. class  $C_\varphi$ . [It is more interesting to ask what is wrong with the given reasoning. The point is just as we cannot prove the Reflection theorem for *all* formulae of the language at once, so we cannot prove more than the existence of *finitely many* such  $C_\varphi$  at any one time in ZF. Hence the given reasoning cannot be given a formal version within ZF. (A platonist who believes absolutely in the universe of all sets  $V$  might also believe in the existence of all such classes and (s)he could then take such an intersection without any qualms!)]

**Exercise 2.44** Find a sentence  $\sigma$  so that if  $\sigma$  is absolute for  $V_\alpha$  then  $\alpha$  is a limit ordinal. Repeat the exercise and find  $\tau$  so that if  $\tau$  is absolute for  $V_\beta$  then  $\beta = \omega_\beta$  (the  $\beta$ 'th infinite cardinal). [Hint: consider the statement: "For every  $\beta$   $\omega_\beta$  exists".]

Sol:  $\sigma \equiv \forall \gamma \exists \beta (\beta \in \text{On} \wedge \beta > \gamma)$ . Then if  $\sigma$  is absolute for  $V_\alpha$  there can be no largest ordinal in  $V_\alpha$ . In other words  $\text{Lim}(\alpha)$  since  $V_\alpha \cap \text{On} = \alpha$  holds for any  $\alpha \in \text{On}$ .

$\tau \equiv \forall \gamma \exists f (\text{ran}(f) \subseteq \text{Card} \wedge \text{dom}(f) = \gamma \wedge f \text{ is a monotone increasing function})$

Then if  $\tau$  is absolute for  $V_\alpha$  it must be that for every  $\gamma < \alpha$  ("there are  $\gamma$  many cardinals below  $\alpha$ ") $^{V_\alpha}$ , in other words (" $\omega_\gamma$  exists for any  $\gamma$ ") $^{V_\alpha}$ . By absoluteness if  $(\kappa \text{ is a cardinal})^{V_\alpha}$  then  $\kappa$  really is a cardinal:  ${}^\kappa \kappa \subseteq V_\alpha$ , so  $V_\alpha$  knows all functions  $f: \gamma \rightarrow \kappa$  for any  $f$ , any  $\gamma \leq \kappa$ . hence  $(\omega_\gamma)^{V_\alpha} = \omega_\gamma$ ; hence  $\alpha$  is a limit point of the cardinal enumeration function  $\gamma \rightarrow \omega_\gamma$ , and we have just reasoned that  $\alpha$  must be  $\omega_\alpha$  itself.

**Exercise 2.45** Let  $x$  be any set, and  $f_i: {}^n V \rightarrow V$  for  $i < \omega$  be any collection of finitary functions (meaning that  $n_i < \omega$ ); show that there is a  $y \supseteq x$  which is closed under each of the  $f_i$  (thus  $f_i \upharpoonright y \subseteq y$  for each  $i$ ) and  $|y| \leq \max\{\omega, |x|\}$ . [Hint: no need for a formal argument here: build up a  $y$  in  $\omega$  many stages  $y_k \subseteq y_{k+1}$  at each step applying all the  $f_i$ .]

Sol: Let  $y_0 = x$ ; let  $y_{k+1} = y_k \cup \{f_i \upharpoonright y_k \mid i < \omega\}$ . Let  $y = \bigcup_{i < \omega} y_i$ . Suppose  $|x| \geq \omega$ . Then each  $|y_{k+1}| = |y_k|$  as  $y_{k+1}$  is the result of adding countably many functions applied to finite tuples from  $y_k$ . But  ${}^{<\omega} y_k$  (the set of finite sequences from  $y_k$ ) has cardinality that of  $y_k$ . So altogether  $|y_{k+1}| = |y_k| + \aleph_0 \cdot |{}^{<\omega} y_k| = |y_k| + |y_k| = |y_k|$ . Then  $y$  is the union of countably many sets of cardinality that of  $x$ . So  $|y| = |x| \cdot \aleph_0 = |x|$ . If  $x$  is finite then possibly  $y_1 = \aleph_0$  (as there are countably many  $f_i$ ), but thereafter all the  $y_k$  and  $y$  will have cardinality  $\aleph_0$ .

**Exercise 2.46** Verify that  $\kappa$  is weakly inaccessible iff  $\kappa$  is regular and  $\kappa = \aleph_\kappa$

Sol: ( $\Rightarrow$ ) By definition  $\kappa$  is regular. If  $\kappa = \aleph_\alpha$  for an  $\alpha < \kappa$  then  $\text{cf}(\kappa) = \text{cf}(\aleph_\alpha) < \kappa$  which contradicts  $\kappa$ 's regularity. Hence, as  $\beta \leq \aleph_\beta$  for any  $\beta$ ,  $\kappa$  must be  $\aleph_\kappa$ .

( $\Leftarrow$ )  $\text{Lim}(\kappa) \rightarrow \aleph_\kappa$  is a limit cardinal. As it is regular, it is weakly inaccessible.

**Exercise 2.47** Does  $\kappa > \omega \wedge V_\kappa = H_\kappa$  imply that  $\kappa$  is strongly inaccessible?

Sol: No. Consider  $\kappa$  the least fixed point of the beth function:  $\kappa = \beth_\kappa$ . This has cofinality  $\omega$  and so is singular. Check that it satisfies  $V_\kappa = H_\kappa$ . [ $H_\kappa \subseteq V_\kappa$  is Lemma 2.31 (ii). Conversely  $|V_\alpha| = 2^{2^\alpha}$  (for  $\omega^2 \leq \alpha$ ; see Ex.2.7). So if  $\alpha < \kappa$ ,  $V_\alpha$  is a transitive set of cardinality  $< \kappa$ . So  $V_\alpha \in H_\kappa$ .]

**Exercise 2.48** Let  $\lambda$  be the least weakly inaccessible cardinal which is itself a limit of weakly inaccessible cardinals (meaning the weakly inaccessible below  $\lambda$  are unbounded in  $\lambda$ ). Show that  $\lambda$  is not weakly Mahlo. Repeat the exercise replacing "weakly" by "strongly" throughout.

Sol: Suppose  $\lambda$  were weakly Mahlo. Then by Lemma 2.50,  $E = \{\alpha < \lambda \mid \alpha \text{ is weakly inaccessible}\}$  is stationary below  $\lambda$ . As usual let  $E^*$  be the set of limit points of  $E$ . Let  $f_E: \lambda \rightarrow E^*$  be the enumerating function of  $E^*$ . The set of  $\alpha < \lambda$  which are fixed points of  $f_E$ , so for which  $f_E(\alpha) = \alpha$ , is a cub set below  $\lambda$  (see Lemma 2.13),  $\bar{E}$  say, and consists of cardinals  $\alpha$  with order type  $\alpha$ -many weakly inaccessible below it. But then  $E \cap \bar{E}$  is stationary, and any  $\alpha \in E \cap \bar{E}$  is then the  $\alpha$ 'th weakly inaccessible cardinal. But  $\lambda$  was supposed to be the least such. The argument is exactly the same for "strongly" in place of "weakly".

**Exercise 3.1** (i) Show that "z is a total order of y" can be expressed in a  $\Delta_0$  fashion.

(ii) Complete (ix), (x), (xvi) and find a correct  $t$  to finish (xvii) of Lemma 3.7.

Sol: (i)  $\forall u, v, w \in \text{dom}(x) \cup \text{ran}(x) [(u = v \vee \langle u, v \rangle \in x \vee \langle v, u \rangle \in x) \wedge (\langle u, v \rangle \in x \wedge \langle v, w \rangle \in x \rightarrow \langle u, w \rangle \in x) \wedge \langle u, u \rangle \notin x]$  (The last clause if the ordering is strict.)

(ii): (ix):  $z \upharpoonright x = \{u \in z \mid (u)_0 \in x\}$ ; (x)  $z^{-1} = \{\langle (u)_1, (u)_0 \rangle \mid u \in z\}$

(xvi): "x is unbounded in  $\beta$ ";  $\forall u \in \beta \exists v \in x (u \in v)$ .

"z:  $\alpha \rightarrow \beta$  is a cofinal function":  $\text{Fun}(z) \wedge \text{ran}(z)$  is unbounded in  $\beta$ .

"x  $\subseteq \beta$  is a closed and unbounded set":  $\forall u \in \beta (\sup(x \cap u) = u \rightarrow u \in \beta) \wedge x$  is unbounded in  $\beta$ .

(xvii) Let  $t(u, v) = u \cup \bigcup v$ . Then define  $f(x) = t(x, \{f(w) \mid w \in x\})$ . Hence  $f$  is definite. Check that  $\text{TC}(x) = f(x)$ .

**Exercise 3.2.** Suppose  $\text{Trans}(W) \wedge (\text{ZF}^-)^W$ . Show  $(V_\alpha)^W = V_\alpha \cap W$ . {Hint: use that the rank function is absolute.}

Sol: Two ways: (i) Use that the rank function  $\rho(x)$  is a.d. and that we know that  $V_\alpha = \{y \mid \rho(y) < \alpha\}$ . Then  $(\rho(y) < \alpha)^W \leftrightarrow \rho(y) < \alpha$ . Thence  $(V_\alpha)^W = \{y \in W \mid (\rho(y) < \alpha)^W\} = \{y \in W \mid \rho(y) < \alpha\} = V_\alpha \cap W$ .

(ii) Or directly by induction on  $\alpha \in W$  (without using the rank function). For  $\alpha = 0$  this is trivial. If true for  $\beta < \alpha$  then (if  $\alpha = \beta + 1$ ) we have that  $(V_\alpha)^W = (\mathcal{P}(V_\beta))^W = \{X \mid X \in W \wedge (X \subseteq V_\beta)^W\}$  by Ind. Hyp.:

$= \{X \mid X \in W \wedge (X \subseteq V_\beta \cap W)\}$  (as " $u \subseteq v$ " is absolute for transitive models)

$= \{X \mid X \in V_\alpha \cap W\}$ .

If  $\text{Lim}(\alpha)$  then  $(V_\alpha)^W = (\bigcup_{\beta < \alpha} V_\beta)^W = \bigcup_{\beta < \alpha} V_\beta \cap W = V_\alpha \cap W$  (using that  $\bigcup$  is absolute, and the Ind.Hyp. for the second equality).



**Exercise 3.3** Let  $\lambda$  be a limit; show that the following are absolute for  $V_\lambda$ : (i)  $\mathcal{P}(x)$  (ii) “ $\alpha$  is a cardinal” (and hence  $(\text{Card})^{V_\lambda} = \text{Card} \cap \lambda$ ); (iii)  $\text{cf}(\alpha)$  (iv) “ $\alpha$  is (strongly) inaccessible” (v)  $V_\alpha$  (vi)  $\aleph_\alpha$  (vii)  $\beth_\alpha$ .

Sol: These are really rather immediate from the definition of the  $V_\alpha$ -hierarchy and the assumption that  $\text{Lim}(\lambda)$ : (i) if  $x \in V_\lambda$  then for some  $\alpha < \lambda$   $x \in V_\alpha \longrightarrow \mathcal{P}(x) \in V_{\alpha+1} \subseteq V_\lambda$ ; (ii) if  $\alpha \in V_\lambda$  then (as  $\text{On} \cap V_\lambda = \lambda$ ) we have that  $\alpha < \lambda$ . But  ${}^\alpha\alpha \subseteq V_{\alpha+5}$ . So  $V_\lambda$  is correct about whether  $\alpha$  is a cardinal or not; likewise  $V_\lambda$  knows if  $\text{cf}(\alpha) = \gamma \leq \alpha$  or not, as all the relevant functions are in  $V_\lambda$  - which is (iii). It is then clear that “ $\alpha$  is inaccessible” is also absolute for  $V_\lambda$ . For *strongly inaccessible* note this comes from (i) and (iii) (as  $\alpha \in \text{Reg} \leftrightarrow \text{cf}(\alpha) = \alpha$ ):

$$(\alpha \text{ is strongly inaccessible})^{V_\lambda} \leftrightarrow (\forall x \in V_\alpha (\mathcal{P}(x) \in V_\alpha) \wedge \alpha \text{ is regular})^{V_\lambda}.$$

Then for any  $\alpha < \lambda$  we shall have  $V_\alpha = (V_\alpha)^{V_\lambda}$ ; if  $(y \text{ is the } \alpha\text{'th cardinal})^{V_\lambda}$  then  $y \text{ is the } \alpha\text{'th cardinal}$ . (More fully expanded, one should use that “ $y \text{ is the } \alpha\text{'th cardinal}$ ” iff  $\exists f(\text{Func}(f) \wedge \text{dom}(f) = \alpha + 1 \wedge \forall \xi < \zeta < \alpha + 1 ((f(\xi) < f(\zeta)) \wedge \text{ran}(f) = \text{Card} \cap y + 1 \wedge f(\alpha) = y)$ ). Then, if  $(y \text{ is the } \alpha\text{'th cardinal})^{V_\lambda}$  there will be such a function  $f \in V_{\alpha+3}$ , and as by (ii)  $(\text{Card})^{V_\lambda} = \text{Card} \cap \lambda$ , all parts of the definition of  $f$  are absolute for  $V_\lambda$ , and so  $y$  really is the  $\alpha$ 'th cardinal, that is  $y \text{ is } \aleph_\alpha$ ). Similarly if  $y = \beth_\alpha < \lambda \leftrightarrow (y = \beth_\alpha)^{V_\lambda}$ . (v) That  $V_\alpha = (V_\alpha)^{V_\lambda}$  may formally be proven by induction on  $\alpha < \lambda$ , but really we have already proven this in Ex 3.2.

**Exercise 3.4** Finish (ii)  $\text{Trans}(x) \longrightarrow x \subseteq \text{Def}(x)$ ; and (iii)  $\forall z \subseteq x (|z| < \omega \rightarrow z \in \text{Def}(x))$ ; of Lemma 3.19.

Sol: (ii) Let  $y \in x$ . Then  $y = \{z \in x \mid \langle x, \in \rangle \vDash \ulcorner v_0 \in v_1 \urcorner [z, y]\}$ .

(iii) Suppose  $z = \{w_1, \dots, w_n\}$ . Then

$$z = \{w \in x \mid \langle x, \in \rangle \vDash \ulcorner v_0 = v_1 \vee \dots \vee v_0 = v_n \urcorner [w, w_1, \dots, w_n]\}$$

**Exercise 3.5** Let  $\langle x, \in \rangle$  be a transitive  $\in$ -model. Show that  $\text{Trans}(\text{Def}(x))$ . If  $y, z \in x$  then is  $\langle y, z \rangle \in \text{Def}(x)$ ? Is  $\{x\}$ ? [Hint (for the last question): If  $\rho(x) = \alpha$ , compute  $\rho(\text{Def}(x))$  and compare this with the given sets.]

Sol: Let  $z \in \text{Def}(x)$ . Then for some  $\varphi$  some  $\vec{y} \in Q_x$ ,  $z = \{u \in x \mid \langle x, \in \rangle \vDash \ulcorner \varphi \urcorner [u, \vec{y}]\}$  so obviously  $z \subseteq x \subseteq \text{Def}(x)$  by the last Exercise. Thus  $\text{Trans}(\text{Def}(x))$ .

However in general  $\langle y, z \rangle \notin \text{Def}(x)$ . Nor is  $\{x\}$ . For the latter this is for the simple reason that  $\{x\} \in \text{Def}(x)$  would imply that  $x \in x$ ! Note that  $z \in \text{Def}(x) \longrightarrow z \subseteq x \longrightarrow \rho(z) \leq \rho(x)$ . But  $x \in \text{Def}(x)$ . Hence  $\rho(\text{Def}(x)) = \rho(x) + 1$ . But if, e.g.,  $\rho(x) = \gamma + 1$  and  $y \in x$  with  $\rho(y) = \gamma$ , then  $\rho(\langle y, y \rangle) = \gamma + 2$ . Hence  $\langle y, y \rangle$  cannot be in  $\text{Def}(x)$  as this has rank only  $\gamma + 2$  itself.

**Exercise 3.6** Let us say that  $w$  is *outright definable in the set*  $\langle x, \in \rangle$  if for some formula  $\varphi$  with only free variable  $v_0$  then  $w$  is the unique element in  $x$  so that  $\langle x, \in \rangle \vDash \ulcorner \varphi \urcorner [w]$ . We may thus define a variant on the Def function by:

$$\text{Def}_0(x) = \{z \mid \{z\} = \{w \mid \langle x, \in \rangle \vDash \ulcorner \varphi \urcorner [w]\}, \text{Fmla}(\varphi) = 1, \text{FVbl}(\varphi) = \{v_0\} \wedge w \in x\}$$

of the sets outright definable in  $\langle x, \in \rangle$ , definable without use of parameters. Show that  $|\text{Def}_0(x)| \leq \omega$  for any  $x$ .

Sol: Just note that irrespective of the size of  $x$ , there are only countably many possible definitions using formulae  $\varphi$  from  $\mathcal{L}$ .

**Exercise 3.7** (i) Show that: (a)  $\text{On} \subseteq \text{OD}^*$ ; (b)  $\forall \beta V_\beta \in \text{OD}^*$ ; (c)  $\forall x (x \in \text{OD}^* \longrightarrow \{x\} \in \text{OD}^*)$ .

(ii) (\*) Show that there is a (countable) set  $X$  so that for unboundedly many ordinals  $\beta$   $X \in \text{Def}_0(V_\beta)$ . Hint: consider the theory of each  $V_\beta$ : the set of all codes of sentences  $\sigma$  so that  $V_\beta \vDash \ulcorner \sigma \urcorner$ . This is a subset of  $V_\omega$ .

Sol: (i) (a)  $\{\alpha\} = \{w \mid V_{\alpha+1} \vDash \ulcorner w \text{ is the largest ordinal} \urcorner\}$ . Similarly (b)  $\{V_\beta\} = \{w \mid V_{\beta+1} \vDash \ulcorner w \text{ is the set of all sets of rank less than the largest ordinal} \urcorner\}$ . (c) Suppose  $x \in \text{Def}_0(V_\beta)$  for some  $\beta$ . Then  $\{x\} = \{w \mid V_\beta \vDash \ulcorner \varphi \urcorner [w]\}$  for some suitable  $\varphi(v_0)$ . But then

$$\{\{x\}\} = \{w \mid V_{\beta+1} \vDash \ulcorner w \text{ is the set of } t \text{ so that } V_\beta \vDash \ulcorner \varphi \urcorner [t] \text{ where } y \text{ is the largest ordinal} \urcorner\}.$$

(ii) For  $\beta \in \text{On}$  let  $\text{Th}(\beta)$  be the ‘theory’ of  $\langle V_\beta, \in \rangle$ , that is

$$\text{Th}(\beta) = \{\ulcorner \sigma \urcorner \mid \ulcorner \sigma \urcorner \text{ is a code of a sentence } \wedge V_\beta \vDash \ulcorner \sigma \urcorner\}.$$

Then note that  $\text{Th}(\beta) \in \text{Def}_0(V_{\beta+1})$ . But each  $\text{Th}(\beta)$  is a, necessarily, countable subset of HF, and then it is easy to see that for unboundedly many  $\beta$  the  $\text{Th}(\beta)$  must be the same  $T$  (as the countable subsets of HF form a set not a proper class!) So this  $T \in \text{Def}_0(V_{\beta+1})$  for unboundedly many  $\beta$ .

**Exercise 3.8** Suppose  $\kappa$  is strongly inaccessible. Verify that  $\langle V_\kappa, \in \rangle \vDash \ulcorner \text{ZFC} \urcorner$ .

Sol: Just note that we have shown (Lemma 2.48 (iii)) for every axiom  $\sigma$  of ZFC that we have  $(\sigma)^{V_\kappa}$ , so by Correctness, Thm 3.21,  $\langle V_\kappa, \in \rangle \vDash \ulcorner \sigma \urcorner$ .

**Exercise 3.9** (\*) (E) Let  $\mathcal{A}, \mathcal{B}$  be structures. We write  $\mathcal{A} < \mathcal{B}$  if for every formula  $u$ , every  $h \in Q_{\mathcal{A}}$  if  $\mathcal{B} \models u[h]$  then  $\mathcal{A} \models u[h]$ . Suppose that  $\kappa < \lambda$  are such that  $\langle V_{\kappa}, \in \rangle < \langle V_{\lambda}, \in \rangle$ . Show that  $\kappa$  is a strong limit cardinal and that both  $\langle V_{\kappa}, \in \rangle, \langle V_{\lambda}, \in \rangle$  are models of ZFC.

Sol: Note that both  $\kappa, \lambda$  are limit ordinals: if  $V_{\kappa} \models \text{“}\alpha \text{ is the largest ordinal”}$ , then  $V_{\lambda} \models \text{“}\alpha \text{ is the largest ordinal”}$ , but  $\alpha + 1$  cannot be both  $\kappa$  and  $\lambda$ . So for any limit  $\tau$  every axiom of ZFC holds in  $V_{\tau}$  except possibly for the Ax Replacement. Let  $f^{V_{\kappa}}$  be a term which defines a function over  $V_{\kappa}$ . Let  $u = \text{dom}(f^{V_{\kappa}})$  and suppose that  $\text{ran}(f^{V_{\kappa}})$  is not a set of  $V_{\kappa}$ . (In other words  $\text{ran}(f^{V_{\kappa}}) \not\subseteq V_{\tau}$  for any  $\tau < \kappa$ , because otherwise  $\text{ran}(f^{V_{\kappa}}) \in V_{\tau+1}$ .) By the definition of  $<$  we have that if we define  $f^{V_{\lambda}}$  we have  $f^{V_{\kappa}}(v) = f^{V_{\lambda}}(v)$  for every  $v \in u$ . But  $V_{\lambda} \models \text{“}\text{there exists } \kappa' \text{ with } \text{ran}(f^{V_{\lambda}}) \subseteq V_{\kappa'}\text{”}$ . However this is not true of  $f^{V_{\kappa}}$  in  $V_{\kappa}$  - a contradiction. So our supposition is false, and  $\text{ran}(f^{V_{\kappa}})$  is a set of  $V_{\kappa}$ ; so AxRep holds in both  $V_{\kappa}$  (and by definition of  $<$ , in  $V_{\lambda}$  too). So  $\langle V_{\kappa}, \in \rangle, \langle V_{\lambda}, \in \rangle$  are models of ZFC. And thus in both models the statement “ $\forall \gamma (\gamma \in \text{Card} \rightarrow \exists \eta (\eta = 2^{\gamma}))$ ” holds (as this follows from ZFC). So both  $\kappa$  (and  $\lambda$ ) are strong limit cardinals. NB we have not shown, and it does not follow from our assumptions, that either of  $\kappa, \lambda$  is regular.

**Exercise 3.12** Suppose  $\langle X, \in \rangle \models T$  for some set of sentences  $T$ . Show that there is a countable transitive  $x$  with  $\langle x, \in \rangle \models T$ . [Hint: The Downward-Loewenheim Skolem Theorem says for any cardinal  $\lambda$  with  $\omega \leq \lambda \leq |X|$  there is a  $Y$  with  $\langle Y, \in \rangle < \langle X, \in \rangle$  and  $|Y| = \lambda$ . Then use the Mostowski-Shepherdson Collapsing Lemma.] In particular if there is an  $\in$ -structure which is a model of ZFC then there is a countable transitive one.

Sol: The Hint really says it all here: by the quoted theorem there is  $\langle Y, \in \rangle < \langle X, \in \rangle$  with  $|Y| = \omega$ . Hence  $\langle Y, \in \rangle \models T$  and  $Y$  has the correct cardinality. It may simply lack transitivity. But this is what the Collapsing Lemma provides: an  $\langle x, \in \rangle \cong \langle Y, \in \rangle$  with  $\text{Trans}(x)$ . As the latter two structures are isomorphic,  $\langle x, \in \rangle \models T$  also. The final sentence applies when  $T$  is ZFC itself.

**Exercise 4.1** (i) Verify that for all  $\alpha \in \text{On}$ ,  $\rho_L(\alpha) = \rho(\alpha) = \alpha$ ; (ii) Prove that for  $n \leq \omega$   $L_n = V_n$ .

Sol: (i) The second equality is already proven. By (iv) of Lemma 4.4  $\alpha = \text{On} \cap L_{\alpha}$ . Hence  $\alpha \in \text{Def}(\langle L_{\alpha}, \in \rangle)$ . ( $\alpha = \{x \in L_{\alpha} \mid \langle L_{\alpha}, \in \rangle \vDash \exists v_0 \in \text{On} \neg [x]\}$ .) Hence  $\alpha \in L_{\alpha+1}$  and  $\rho_L(\alpha) \leq \alpha$ . But  $\alpha$  cannot be a member of  $L_{\alpha}$ , hence we have equality here.

(ii) By induction on  $n$ : for  $n=0$  this is trivial. Suppose true for  $n=k$ . Since  $\text{Def}(\langle L_k, \in \rangle) \subseteq \mathcal{P}(L_k)$ , then by the induction hypothesis we have  $\text{Def}(\langle L_k, \in \rangle) \subseteq \mathcal{P}(V_k) = V_{k+1}$ . Let  $x = \{x_1, \dots, x_m\} \in \mathcal{P}(V_k)$ . Then  $x \subseteq L_k$ , and  $x \in \text{Def}(\langle L_k, \in \rangle) = L_{k+1}$ , by (iii) of Exercise 3.1.

**Exercise 4.3** Show that “ $x$  is a cardinal” and “ $x$  is regular” are downward absolute from  $V$  to  $L$ . Deduce that if  $\kappa$  is a (regular) limit cardinal then  $(\kappa \text{ is a (regular) limit cardinal})^L$

Sol: Being a cardinal, and being regular are downward absolute (so this is nothing special about  $L$ ) as follows.

$x \text{ is a cardinal} \leftrightarrow x \in \text{On} \wedge \forall f (\text{Func}(f) \wedge \text{dom}(f) < x \rightarrow \text{ran}(f) \neq x)$ .

Being an ordinal is definite hence absolutely definite, and the second conjunct is a universal quantifier in front of a definite,  $\Sigma_0$ , matrix. Hence if this holds for all functions  $f$  it certainly holds for all functions in any subclass of  $V$  (such as  $L$ ). Similarly

$x \text{ is a regular} \leftrightarrow x \in \text{Card} \wedge \forall f (\text{Func}(f) \wedge \text{dom}(f) < x \wedge \text{ran}(f) \subseteq x \rightarrow \sup \text{ran}(f) \neq x)$

is also (equivalent to) a universally quantified statement, and both are downward absolute to  $L$ . Finally then, if  $x$  is a limit of cardinals  $\lambda$ , then those  $\lambda$  are cardinals in  $L$ , hence  $(x \text{ is a limit of cardinals})^L$ . If  $x$  is regular in addition, then  $x$  will be a regular limit cardinal (and hence inaccessible) in  $L$ .

**Exercise 4.5** Show that  $\text{ot}(L_{\kappa}, <_{\kappa}) = \kappa$  for  $\kappa$  an infinite cardinal; Show that  $\text{ot}(L, <_L) = \text{On}$ .

Sol: If  $\alpha < \kappa$  we saw (Lemma 4.5 (iii)) that  $|L_{\alpha}| = |\alpha| < \kappa$ . Hence  $\text{ot}(L_{\alpha}, <_{\alpha})$  must be strictly less than the cardinal  $\kappa$ . As  $<_{\kappa} = \bigcup_{\alpha < \kappa} <_{\alpha}$  and  $\alpha < \beta < \kappa \rightarrow <_{\beta}$  is an end extension of  $<_{\alpha}$ , we must have that  $\text{ot}(L_{\kappa}, <_{\kappa}) \leq \kappa$  (for if  $x \in L_{\kappa}$  is the  $\kappa$ 'th set in  $L_{\kappa} = \bigcup_{\alpha < \kappa} L_{\alpha}$  then for some  $\beta < \kappa$  we should have that  $x$  is the  $\kappa$ 'th set in  $L_{\beta}$  which is a contradiction). As  $\alpha < \kappa \rightarrow \alpha \in L_{\alpha+1}$  we must have that  $\kappa \subseteq \text{Field}(<_{\kappa})$  and so  $\text{ot}(L_{\kappa}, <_{\kappa}) \geq \kappa$ . Entirely the same argument works to show  $\text{ot}(L, <_L) = \text{On}$  thinking now of  $\text{On}$  in place of  $\kappa$ .

**Exercise 4.6** Suppose there is a transitive set model of ZFC. Show that there is a *minimal model* of ZFC, that is for some countable ordinal  $\beta_0$ ,  $(ZFC)^{L_{\beta_0}}$  and that  $L_{\beta_0}$  is a subclass of any other transitive set model of ZF.

Sol: Suppose there is a transitive set model of ZFC. Then there is such of minimal rank, say  $M_0$ , with  $\rho(M_0) = \beta_0$ . Then  $\langle M_0, \in \rangle \models \ulcorner ZFC \urcorner$  implies  $\langle M_0, \in \rangle \models \ulcorner \sigma_1 \urcorner$  where  $\sigma_1$  is the conjunction of finitely many axioms of ZF needed to prove absoluteness and results about the  $L$  construction (as in Lemma 4.12). By the Correctness Theorem we have  $(\sigma_1)^{M_0}$  and thus by Lemma 4.12,  $L^{M_0} = L_{\beta_0}$ ; moreover  $\langle M_0, \in \rangle \models \ulcorner L_{\beta_0} \models \ulcorner ZFC \urcorner \urcorner$ .  $L_{\beta_0}$  must then be the minimal model of ZFC as required.

**Exercise 4.7** (Shepherdson) Show that there is no class term  $W$  so that  $ZFC \vdash \text{IM}(W)$  and  $ZFC \vdash (\neg\text{CH})^W$ . [This Exercise shows that Goedel's argument was essentially a "one-off": there is no way one can define in ZFC alone an inner model and hope that it is a model of all of ZF plus, e.g.,  $\neg\text{CH}$ .]

Sol: Suppose there were such a term  $W$ . Then let  $U$  be the term  $W$  relativised to  $L$ :  $U = (W)^L$ . Thus  $U \subseteq L$ . However then By Lemma 4.12  $(L)^U = L = U$ . But  $(\neg\text{CH})^U$  since by assumption  $ZFC \vdash (\neg\text{CH})^W$  and  $(ZFC)^L$ . Contradiction.

**Exercise 4.8** Show that if there is a weakly inaccessible cardinal  $\kappa$  then  $(ZFC)^{L_\kappa}$ . Hence  $ZFC \not\vdash \exists \kappa (\kappa \text{ a weakly inaccessible cardinal.})$  [Hint: Use the fact that  $(\text{GCH})^L$ .]

Sol: If  $\kappa$  is a regular cardinal then  $H_\kappa$  is a model of  $ZFC^-$  by Lemma 2.29 (although Power Set may fail). As  $(H_\kappa)^L = L_\kappa$  we have all instances of ZFC axioms with the possible exception of Power Set holding in  $L_\kappa$ . But if  $\alpha < \kappa$  is any cardinal, then  $(|\mathcal{P}(\alpha)| = \alpha^+ < \kappa)^L$  as  $(\text{GCH})^L$ . (Note that  $(\alpha^+)^L \leq \alpha^+ < \kappa$  as  $\kappa$  is a limit cardinal.) Thus Power Set does hold in  $L_\kappa$  after all. Thus if  $\kappa$  is weakly inaccessible cardinal, we have deduced that  $\kappa$  is a strongly inaccessible cardinal of  $L$ . But then (Lemma 2.48)  $(ZFC)^{L_\kappa}$ , and we have a model of the ZFC axioms. The result follows by the Second Incompleteness theorem.

**Exercise 4.10** Assume  $V = L$ . When does  $V_\alpha = L_\alpha$ ?

**Exercise 4.12** Show that (i) if  $\kappa$  is a weakly inaccessible cardinal, then  $(\kappa \text{ is strongly inaccessible})^L$ ; (ii) if  $\kappa$  is a weakly Mahlo cardinal, then  $(\kappa \text{ is strongly Mahlo})^L$ . [Hint: See Exercises 4.3 & 4.8. For (ii) show that the property of being cub in  $\kappa$  is preserved upwards from  $L$  to  $V$ .]

Sol: (i) Similar to Ex 4.8, if  $\alpha < \kappa$  is a cardinal then  $(\alpha \text{ is a cardinal})^L$  (We check this one more time: for if  $\alpha$  was not a cardinal in  $L$ , then in  $L$  for some  $\gamma < \alpha$  there is a function  $f \in L$  with  $f: \gamma \rightarrow \alpha$  onto; but this is absolute: so  $f$  really does show in  $V$  that  $\alpha$  is not a cardinal!) Hence  $(\kappa \text{ is a limit cardinal})^L$ . But  $\kappa$  is not a singular cardinal in  $L$  for the same reason that there is no constructible function demonstrating its singularity. So  $\kappa$  is a regular limit cardinal of  $L$ . Hence it is (strongly) inaccessible, because - as in Exercise 2.8 - no  $\alpha < \kappa$  has its constructible power set of cardinality  $\geq \kappa$ .

(ii) This is slightly different. Suppose  $C \in L$  and  $(C \text{ is closed and unbounded in } \kappa)^L$ . Note how being "closed and unbounded in  $\kappa$ " is in fact an absolute notion. So  $C$  really is such. As  $\kappa$  is weakly Mahlo in  $V$  there is a regular cardinal  $\alpha \in C$ . As being regular is downward absolute to  $L$  (Lemma 4.3 again)  $(\alpha \text{ is regular})^L$ . As  $C$  was arbitrary,  $L$  sees that  $\kappa$  is weakly Mahlo. As weakly Mahlo implies that  $\kappa$  is weakly inaccessible, we have that  $\kappa$  is strongly inaccessible in  $L$  by part (i) and therefore *strongly* Mahlo there.

**Exercise 4.13** (i) Let  $\langle x, \in \rangle \prec L_{\omega_1}$  where  $\omega_1 = (\omega_1)^L$ . Show that already  $\text{Trans}(x)$  and so  $x = L_\gamma$  for some  $\gamma \leq \omega_1$ . [Hint: For  $\delta < \omega_1$  note that  $(|\delta| = |L_\delta| = \omega)^{L_{\omega_1}}$ . Hence for  $\delta \in x$ , in  $L_{\omega_1}$ , and thus in  $x$ , there is an onto map  $f: \omega \rightarrow L_\delta$ . Thus, as  $\omega \subseteq x \wedge f \in x$  we deduce that  $\text{ran}(f) = L_\delta \subseteq x$ . Deduce that  $\text{Trans}(x)$ .]

Sol: (i) Notice that it suffices to show that  $\text{Trans}(x)$ : for then  $(ZF^-)^x$  since  $(ZF^-)^{L_{\omega_1}}$  (as  $(ZF^-)^{H_{\omega_1}}$  by Lemma 2.29, but in  $L$ ,  $H_{\omega_1} = L_{\omega_1}$ ). By Lemma 4.12 this means  $L^x = L_\gamma$  where  $\gamma = \text{On} \cap x$ . However  $x$  is itself a model of " $V = L$ " (i.e.  $\forall z \exists \beta (z \in L_\beta)$ ) because the same is true of  $L_{\omega_1}$ . Thus  $x$  can only be  $L_\gamma$ ! We now show transitivity: first note that  $L_{\omega_1}$  is a model of the statement that everything is countable. Following the Hint, we see that for every  $\delta$  in  $x$ , there is, also in  $x$ , an onto function  $f_\delta: \omega \rightarrow L_\delta$ . But  $\text{ran}(f_\delta) \subseteq x$  since for every  $n \in \omega$ ,  $n$  has a first order definition and so  $n \in x$ ; as  $f_\delta \in x$ , we conclude that  $f_\delta(n) \in x$ . We deduce that  $L_\delta \subseteq x$ . As  $x$  is a model of  $V = L$ , for every  $z \in x$  there is some  $\beta \in x$  with  $z \in L_\beta$ . But we have just seen that  $L_\beta \subseteq x$ . Hence  $z \subseteq L_\beta \subseteq x$  (the first inclusion coming from  $\text{Trans}(L_\beta)$ ).

**Exercise 4.18** Show (i)  $z \in \text{HOD} \leftrightarrow z \in \text{OD} \wedge \forall y \in z (y \in \text{HOD})$ . (ii)  $\mathcal{P}(\omega) \cap \text{OD} = \mathcal{P}(\omega) \cap \text{HOD}$ .

Sol: (i)  $z \in \text{HOD} \leftrightarrow_{\text{df}} z \in \text{OD} \wedge \text{TC}(z) \subseteq \text{OD} \leftrightarrow z \in \text{OD} \wedge (z \cup \bigcup \{ \text{TC}(y) \mid y \in z \}) \subseteq \text{OD}$   
 $\leftrightarrow z \in \text{OD} \wedge \forall y \in z (\{y\} \cup \text{TC}(y)) \subseteq \text{OD} \leftrightarrow z \in \text{OD} \wedge \forall y \in z (y \in \text{HOD})$ .

(ii): ( $\supseteq$ ): trivial. ( $\subseteq$ ): Every integer is in HOD. hence if  $x \in \mathcal{P}(\omega) \cap \text{OD}$ , then  $x \in \text{HOD}$  by part(i).

**Exercise 4.19** Show that for any  $\beta$ ,  $V_\beta \cap \text{HOD} \in \text{HOD}$ .

Sol: If  $z = V_\beta \cap \text{HOD}$ , then  $z \subseteq \text{HOD}$  so by the last exercise it suffices to show  $z \in \text{OD}$ . But  $z$  is uniquely defined from the ordinal  $\alpha$  as “the set of  $y \in V_\beta$  so that  $\forall u \in \text{TC}(\{y\}) \exists \alpha (u \in \text{Def}_0(V_\alpha))$ ”.

**Exercise 4.20** Show that the following are equivalent: (i)  $V = \text{OD}$ , (ii)  $V = \text{HOD}$ , (iii)  $\text{Trans}(\text{OD})$ , (iv)  $(\text{AxExt})^{\text{OD}}$ . [Hint: Use that for any  $\alpha$   $V_\alpha \in \text{OD} \wedge V_\alpha \cap \text{OD} \in \text{OD}$ .]

Sol: (i)  $\Rightarrow$  (ii) is straightforward from the definition of HOD. (ii)  $\Rightarrow$  (iii):  $\text{HOD} \subseteq \text{OD} \subseteq V \wedge \text{Trans}(V)$ . (iii)  $\Rightarrow$  (iv):  $\text{Trans}(W) \longrightarrow (\text{AxExt})^W$  for any class term  $W$ . (iv)  $\Rightarrow$  (i): As  $V_\alpha \cap \text{OD} \in \text{OD}$  (almost identical to Ex. 4.18) and also  $V_\alpha \in \text{OD}$ , then by  $(\text{AxExt})^{\text{OD}}$  if these are different, then for some  $x \in \text{OD}$ ,  $x \in V_\alpha \cap \text{OD} \Delta V_\alpha$ . But these facts are contradictory.

**Exercise 4.21** Show that  $\text{HOD} \cap \mathcal{P}(\omega)$  is the largest subset of  $\mathcal{P}(\omega)$  with a definable wellorder.

Sol: If  $B = \text{HOD} \cap \mathcal{P}(\omega)$ , then  $B \subseteq \text{OD}$  and so by restricting the definable wellorder of OD to  $B$ , it has a definable wellorder. By Lemma 4.23 if  $A \subseteq \mathcal{P}(\omega)$  has a definable wellorder then  $A \subseteq \text{OD}$ . But by Ex 4.17 then  $A \subseteq \text{HOD} \cap \mathcal{P}(\omega)$ .

**Exercise 4.22** Suppose that  $W$  is a term defining an inner model of ZF and there is a definable global wellorder of  $W$  (that, as in  $L$ , there is a formula defining a wellorder  $<_W$  of the whole of  $W$  in order type  $\text{On}$ ). Show that  $W \subseteq \text{HOD}$ . (Consequently HOD is the largest inner model  $W$  with a definable bijection  $\text{On} \leftrightarrow W$ .)

Sol: Again HOD has a definable wellorder by restricting the definable wellorder of OD to it. By Lemma 4.23  $W \subseteq \text{OD}$ . But  $\text{IM}(W) \Rightarrow \text{Trans}(W)$ , so  $W \subseteq \text{HOD}$ .

**Exercise 4.23** Define “ $\Pi_2$ -OD” (and  $\Pi_2$ -HOD) just as we did for OD and HOD but now restrict the formulae allowed in definitions to be  $\Pi_2$  only. Show that  $\Pi_2\text{-OD} = \text{OD}$  and  $\Pi_2\text{-HOD} = \text{HOD}$ . Now do the same for  $\Sigma_2\text{-OD}$  and  $\Sigma_2\text{-HOD}$ .

Sol: First note that “ $\rho(z) < \alpha$ ” can be expressed both in  $\Sigma_1$  and  $\Pi_1$ -form: this is  $\exists f [\Phi(f, z) \wedge f(z) < \alpha]$  and  $\forall f [\Phi(f, z) \longrightarrow f(z) < \alpha]$  where

$$\Phi(f, z) \equiv \text{Fun}(f) \wedge \text{TC}(\{z\}) \subseteq \text{dom}(f) \wedge \forall t \in z (f(t) = \sup \{f(s) + 1 \mid s \in t\}).$$

Then note that “ $y = V_\alpha$ ” is  $\Pi_1$ . This is  $\forall z (z \in y \leftrightarrow \rho(z) < \alpha)$  (using each form in turn for  $\rho(z) < \alpha$ ).

Let  $x \in \text{OD}$ . Then there is some  $\varphi(v_0)$  and some  $\beta$  so that

$$\{x\} = \{w \mid V_\beta \vDash \varphi(w)\} = \{w \mid \forall y (y = V_\beta \rightarrow \langle y, \in \rangle \vDash \varphi[w])\}.$$

Now use that “ $\langle u, \in \rangle \vDash \varphi[w]$ ” is  $\Delta_1$  expressible and we have a  $\Pi_2$  definition of  $x$ . But we can also write:

$$\{x\} = \{w \mid \exists y \exists \beta (y = V_\beta \wedge \langle y, \in \rangle \vDash \varphi[w])\} \text{ in a } \Sigma_2 \text{ form. Thus } \Sigma_2\text{-OD will be OD.}$$

The equalities for  $\Pi_2\text{-HOD} = \text{HOD}$  and similarly  $\Sigma_2\text{-HOD} = \text{HOD}$  now just follow from the definition of HOD.

**Exercise 4.24\*** Show that there is a *single* formula  $\varphi_0(v_0)$  with just the free variable shown, so that  $\text{OD}$  is the class of all those  $x$  so that  $x \in \text{Def}_0(V_\beta)$  for some  $\beta$ , that is for some  $\beta$ ,  $\{x\} = \{z \mid \varphi_0(z)^{V_\beta}\}$ .