

On Field's  
*"Saving Truth  
from Paradox"*

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(I) In current semantic theories

(a) there is a lack of a viable conditional  $\supset$

(b) there may be occurrences of failure of the  $T$ -scheme:

(that for some sentence(s)  $A$  we do not have ' $T(\ulcorner A \urcorner) \leftrightarrow A$ ' )

(c) there may be failure of *the intersubstitutivity* of  $T(\ulcorner A \urcorner)$  for  $A$  where the latter is a subformulae of some  $B$ .

So Field:

- (II) introduces a *binary operator*  $\longrightarrow$  to function as a form of generalised conditional;
- (III) provides for set models (one or more) semantics that remedies (a)-(c) in a non-classical logic;
- (IV) provides an analysis of the 'defectiveness' of *e.g.* the liar sentence through a hierarchy of *determinateness relations*.

## $\mathcal{M}^+$ Expansions: *G-solutions*

- $\mathcal{M}$  an  $\mathcal{L} = \mathcal{L}_{\mathcal{M}}$ -structure : expanded to  $\mathcal{M}^+$  in  $\mathcal{L}^+ = \mathcal{L}_{\mathcal{M}}^+$  (containing additionally  $T, \longrightarrow$ )
- $\mathcal{L}^+$  evaluated in a 3-valued  $\{0, \frac{1}{2}, 1\}$  logic, (or in a De Morgan function algebra  $V_c$ ).

- He takes issue with the idea that we can *define* ‘real truth’ by using notions of designated semantic values obtained by a variety of methods (Kripkean minimal fixed points, strong Kleene, supervaluational ... or *via* revision theory. In particular Tarski like constructions using *set-sized* models for which we can give mathematical or inductive definitions of ‘designated truth value’ cannot deliver for us a theory of ‘real truth.’
- But in order to give some description of the theory he is aiming for, he has given a number of *G*-solutions or *G*-models, in particular a ‘principal’ one deriving the consistency of the naive theory of truth (The *T*-scheme), the Intersubstitutivity Principle, and with  $\longrightarrow$ .
- We thus have *G*-models  $\mathcal{M}^+$  over *e.g.* ground models  $\mathcal{M}$  :

$$\langle \mathbb{N}, +, \times, 0, S \rangle, \quad \langle V_\alpha, \in \rangle, \dots$$

## Real Validity

*“it might be better to adopt the view that what is validated by a given version of the formal semantics [i.e. a G-model] outruns “real validity”: that the genuine logical validities are some effectively generable subset of those inferences that preserve value 1”*

(emphasis now mine). He continues:

*“... there would doubtless be some arbitrariness in which effectively generable subset to choose, but that is perfectly acceptable unless one wants to put high (and I think unreasonable) demands on the significance of the distinction between those inferences that are valid and those that are not.”*

## The complexity of the principal model: a concern

- The principal model over  $\mathbb{N}$  is supposed to deliver a *first order* theory of truth with  $\longrightarrow$ ;
- however it requires (as a piece of applied mathematics) a stronger subsystem of second order number theory (  $\Pi_3^1$ -CA ) than any other piece of ‘ordinary’ mathematics
- Well beyond the reach of any current proof-theoretical ordinal analysis
- Can there be a simpler ‘consistency proof’? A simpler ‘principal model’?

# Internal Structure of the Principal Model

- Recall that we have a determinateness operator  $D(A) \equiv A \wedge \neg(A \rightarrow \neg A)$
- $D^{n+1}(A) \equiv D(D^n A)$  ;  $D^\omega(A) \equiv \forall n \forall y (y = \ulcorner D^n(A) \urcorner \rightarrow T(y))$ .
- Comes with ‘determinate liars’:  $Q^\alpha \leftrightarrow \neg D^\alpha(\ulcorner Q^\alpha \urcorner)$ .
- How far can these hierarchies go?

## Taking $\mathcal{M}$ as $\mathbb{N}$

- To go beyond recursive ordinals let *sentences of  $\mathcal{L}^+$  themselves* stand in for ordinal notations:

### Definition

$\rho(A) \simeq$  least  $\rho$  such that semantic value of  $\rho$  is constant from  $\rho$  onwards.



We abbreviate  $A \prec B$  for  $P_{\prec}(\ulcorner A \urcorner, \ulcorner B \urcorner) = 1$  etc.

- If  $\|A\| = 1$  (or 0) say, then  $\{B : B \prec A\} = \{B : \|P_{\prec}(\ulcorner A \urcorner, \ulcorner B \urcorner)\| = 1\}$  is a prewellordering of order type some ordinal  $\xi < \Delta_0$ .
- We let  $\text{Field}(\prec)$  denote the set of sentences stabilizing on 0 or 1.

Moreover:

### Lemma

There is a formula  $P_{\prec}(v_0, v_1)$  in  $\mathcal{L}^+$  so that for any sentences  $A, B \in \mathcal{L}^+$ , we have  $\|P_{\prec}(\ulcorner A \urcorner, \ulcorner B \urcorner)\| = 1$  iff  $\rho(A) \downarrow, \rho(B) \downarrow$  and  $\rho(A) < \rho(B)$ ;

$$\begin{aligned} &= 0 \text{ iff } \rho(A) \downarrow, \rho(B) \downarrow \text{ and } \rho(A) \geq \rho(B); \\ &= \frac{1}{2} \text{ otherwise.} \end{aligned}$$

### Lemma

For any  $\xi < \Delta_0$  there is a sentence  $A = A_{\xi}$  in  $\text{Field}(\prec)$  with the order type of  $\{B \mid B \prec A\}$  equalling  $\xi$ .

- We may define for *any* sentence  $C$

$$D^C(A) \equiv \forall B \prec C \forall y (y = \ulcorner D^B(A) \urcorner \rightarrow T(y)).$$

- For  $C \in \text{Field}(\prec)$  this defines a bivalent determinateness hierarchy of length  $\rho(C)$ .
- However it is not a bivalent matter as to whether a general  $C$  is or is not in  $\text{Field}(\prec)$ . (In other words  $\text{Field}(\prec)$  is not a crisp subclass of  $\mathbb{N}$ .) However if  $C \in \text{Field}(\prec)$  then it can be shown that it is a bivalent matter whether a general  $B$  is  $\prec$ -below  $C$  or not.
- Consequently the expression

“ $\langle D^B(v_0) \mid B \prec C \text{ forms a determinateness hierarchy} \rangle$ ”

is not in the classical part of the language  $\mathcal{L}^+$  to which the Law of Excluded Middle holds.

- Thus the internally defined determinateness hierarchy over  $\mathbb{N}$  breaks down, not fuzzily, but precisely, at  $\Delta_0$ . There is no internally definable maximal hierarchy.

## Axiomatising $F =_{\text{df}} \{A : \|A\| = 1\}$

- (Martin) We have an *open game representation* in  $\mathcal{L}$  of the least Strong Kleene fixed point over  $\mathbb{N}$  as an *open game*.
- A game for  $F$  can be formulated but is an  $\exists\forall\exists$  game.
- One can have an open game representation of  $F$  over  $\mathbb{N}$  in a language with a generalised quantifier  $\mathcal{L}^+(Q)$ , where
- ‘ $Qx\varphi x$ ’ iff for ‘path-many’  $x \varphi(x)$ ; that is

$$Qx\varphi x \Leftrightarrow \exists A \in \text{Field}(\prec)(\{n \in \mathbb{N} \mid \varphi(n)\} \supseteq \{B \mid B \preceq A\}).$$

- Speculatively this suggests a possible axiomatisation of a theory of truth (with a  $\longrightarrow$ ) *together with* determinateness satisfying the laws or properties Field has already given for  $T$  and  $D$ .